

Formalizing MLTL in Isabelle/HOL

Zili Wang and Katherine Kosaian

March 17, 2025

Abstract

Building on the formalization of Mission-time Linear Temporal Logic (MLTL) in Isabelle/HOL, we formalize the correctness of the algorithms for the WEST tool [1, 2], which converts MLTL formulas to regular expressions. We use Isabelle/HOL’s code export to generate Haskell code to validate the existing (unverified) implementation of the WEST tool.

Contents

1	Key algorithms for WEST	2
1.1	Custom Types	2
1.2	Trace Regular Expressions	2
1.3	WEST Operations	4
1.3.1	AND	4
1.3.2	Simp	5
1.3.3	AND and OR operations with WEST-simp	7
1.3.4	Useful Helper Functions	7
1.3.5	WEST Temporal Operations	8
1.3.6	WEST recursive reg Function	9
1.3.7	Adding padding	11
2	Some examples and Code Export	12
3	WEST Proofs	12
3.1	Useful Definitions	12
3.2	Proofs about Traces Matching Regular Expressions	13
3.3	Facts about the WEST and operator	13
3.3.1	Commutative	13
3.3.2	Identity and Zero	16
3.3.3	WEST-and-state	16
3.3.4	WEST-and-trace	18
3.3.5	WEST-and correct	21

3.4	Facts about the WEST or operator	22
3.5	Pad and Match Facts	22
3.6	Facts about WEST num vars	23
3.6.1	Facts about num vars for different WEST operators	23
3.7	Correctness of WEST-simp	27
3.7.1	WEST-count-diff facts	27
3.7.2	Orsimp-trace Facts	29
3.7.3	WEST-orsimp-trace-correct	30
3.7.4	Simp-helper Correct	31
3.7.5	WEST-simp Correct	32
3.8	Correctness of WEST-and-simp/WEST-or-simp	33
3.9	Facts about the WEST future operator	33
3.10	Facts about the WEST global operator	34
3.11	Facts about the WEST until operator	35
3.12	Facts about the WEST release Operator	36
3.13	Top level result: Shows that WEST reg is correct	37
3.14	Top level result for padded version	37
4	Key algorithms for WEST	38
5	Regex Equivalence Correctness	39

1 Key algorithms for WEST

theory *WEST-Algorithms*

imports *Mission-Time-LTL.MLTL-Properties*

begin

1.1 Custom Types

datatype *WEST-bit* = *Zero* | *One* | *S*

type-synonym *state* = *nat set*

type-synonym *trace* = *nat set list*

type-synonym *state-regex* = *WEST-bit list*

type-synonym *trace-regex* = *WEST-bit list list*

type-synonym *WEST-regex* = *WEST-bit list list list*

1.2 Trace Regular Expressions

fun *WEST-get-bit:: trace-regex $\Rightarrow$ nat $\Rightarrow$ nat $\Rightarrow$ WEST-bit*

where *WEST-get-bit regex timestep var = (*

if timestep $\geq$ length regex then S

else let regex-index = regex ! timestep in

if var $\geq$ length regex-index then S

else regex-index ! var

)

Returns the state at time i, list of variable states

```
fun WEST-get-state:: trace-regex  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$  state-regex
where WEST-get-state regex time num-vars = (
  if time  $\geq$  length regex then (map ( $\lambda$  k. S) [0 ..< num-vars])
  else regex ! time
)
```

Checks if one state of a trace matches one timeslice of a WEST regex

```
definition match-timestep:: nat set  $\Rightarrow$  state-regex  $\Rightarrow$  bool
where match-timestep state regex-state = ( $\forall$  x::nat. x < length regex-state  $\longrightarrow$ 
(
  ((regex-state ! x = One)  $\longrightarrow$  x  $\in$  state)  $\wedge$ 
  ((regex-state ! x = Zero)  $\longrightarrow$  x  $\notin$  state)))
```

```
fun trim-reversed-regex:: trace-regex  $\Rightarrow$  trace-regex
where trim-reversed-regex [] = []
| trim-reversed-regex (h#t) = (if ( $\forall$  i<length h. (h!i) = S)
  then (trim-reversed-regex t) else (h#t))
```

```
fun trim-regex:: trace-regex  $\Rightarrow$  trace-regex
where trim-regex regex = rev (trim-reversed-regex (rev regex))
```

```
definition match-regex:: nat set list  $\Rightarrow$  trace-regex  $\Rightarrow$  bool
where match-regex trace regex = (( $\forall$  time<length regex.
  (match-timestep (trace ! time) (regex ! time)))
 $\wedge$  (length trace  $\geq$  length regex))
```

```
definition match:: nat set list  $\Rightarrow$  WEST-regex  $\Rightarrow$  bool
where match trace regex-list = ( $\exists$  i. i < length regex-list  $\wedge$ 
  (match-regex trace (regex-list ! i)))
```

```
lemma match-example:
shows match [{0::nat,1}, {1}, {0}]
[
  [[Zero,Zero]],
  [[S,S], [S,One]]
] = True
<proof>
```

```
definition regex-equiv:: WEST-regex  $\Rightarrow$  WEST-regex  $\Rightarrow$  bool
where regex-equiv rl1 rl2 = (
 $\forall$   $\pi$ ::nat set list. (match  $\pi$  rl1)  $\longleftrightarrow$  (match  $\pi$  rl2))
```

```
lemma (regex-equiv [[[S,S]]]
  [[[S,One]]],
```

$$[[One, S],$$

$$[[Zero, Zero]]]) = True$$

$$\langle proof \rangle$$

1.3 WEST Operations

1.3.1 AND

fun *WEST-and-bitwise*:: *WEST-bit* $\Rightarrow$
 $WEST-bit \Rightarrow$
 $WEST-bit\ option$
where *WEST-and-bitwise* *b* *One* = (if *b* = *Zero* then *None* else *Some One*)
| *WEST-and-bitwise* *b* *Zero* = (if *b* = *One* then *None* else *Some Zero*)
| *WEST-and-bitwise* *b* *S* = *Some b*

fun *WEST-and-state*:: *state-regex* $\Rightarrow$ *state-regex* $\Rightarrow$ *state-regex option*
where *WEST-and-state* [] [] = *Some* []
| *WEST-and-state* (*h1* # *t1*) (*h2* # *t2*) =
(case *WEST-and-bitwise* *h1* *h2* of
None $\Rightarrow$ *None*
| *Some b* $\Rightarrow$ (case *WEST-and-state* *t1* *t2* of
None $\Rightarrow$ *None*
| *Some L* $\Rightarrow$ *Some (b* # *L*)))
| *WEST-and-state* - - = *None*

fun *WEST-and-trace*:: *trace-regex* $\Rightarrow$ *trace-regex* $\Rightarrow$ *trace-regex option*
where *WEST-and-trace* *trace* [] = *Some trace*
| *WEST-and-trace* [] *trace* = *Some trace*
| *WEST-and-trace* (*h1* # *t1*) (*h2* # *t2*) =
(case *WEST-and-state* *h1* *h2* of
None $\Rightarrow$ *None*
| *Some state* $\Rightarrow$ (case *WEST-and-trace* *t1* *t2* of
None $\Rightarrow$ *None*
| *Some trace* $\Rightarrow$ *Some (state* # *trace*)))

fun *WEST-and-helper*:: *trace-regex* $\Rightarrow$ *WEST-regex* $\Rightarrow$ *WEST-regex*
where *WEST-and-helper* *trace* [] = []
| *WEST-and-helper* *trace* (*t* # *traces*) =
(case *WEST-and-trace* *trace* *t* of
None $\Rightarrow$ *WEST-and-helper* *trace* *traces*
| *Some res* $\Rightarrow$ *res* # (*WEST-and-helper* *trace* *traces*)

fun *WEST-and*:: *WEST-regex* $\Rightarrow$ *WEST-regex* $\Rightarrow$ *WEST-regex*
where *WEST-and* *traceList* [] = []

```

| WEST-and [] traceList = []
| WEST-and (trace#traceList1) traceList2 =
(case WEST-and-helper trace traceList2 of
[] => WEST-and traceList1 traceList2
| traceList => traceList@(WEST-and traceList1 traceList2))

```

1.3.2 Simp

Bitwise simplification operation **fun** WEST-simp-bitwise:: WEST-bit => WEST-bit => WEST-bit

```

where WEST-simp-bitwise b S = S
| WEST-simp-bitwise b Zero = (if b = Zero then Zero else S)
| WEST-simp-bitwise b One = (if b = One then One else S)

```

fun WEST-simp-state:: state-regex => state-regex => state-regex

```

where WEST-simp-state s1 s2 = (
map (λ k. WEST-simp-bitwise (s1 ! k) (s2 ! k)) [0 ..< (length s1)])

```

fun WEST-simp-trace:: trace-regex => trace-regex => nat => trace-regex

```

where WEST-simp-trace trace1 trace2 num-vars = (
map (λ k. (WEST-simp-state (WEST-get-state trace1 k num-vars) (WEST-get-state
trace2 k num-vars)))
[0 ..< (Max {(length trace1), (length trace2)}]))

```

Helper functions for defining WEST-simp **fun** count-nonS-trace:: state-regex => nat

```

where count-nonS-trace [] = 0
| count-nonS-trace (h#t) = (if (h ≠ S) then (1 + (count-nonS-trace t)) else
(count-nonS-trace t))

```

fun count-diff-state:: state-regex => state-regex => nat

```

where count-diff-state [] [] = 0
| count-diff-state trace [] = count-nonS-trace trace
| count-diff-state [] trace = count-nonS-trace trace
| count-diff-state (h1#t1) (h2#t2) = (if (h1 = h2) then (count-diff-state t1 t2)
else (1 + (count-diff-state t1 t2)))

```

fun count-diff:: trace-regex => trace-regex => nat

```

where count-diff [] [] = 0
| count-diff [] (h#t) = (count-diff-state [] h) + (count-diff [] t)
| count-diff (h#t) [] = (count-diff-state [] h) + (count-diff [] t)
| count-diff (h1#t1) (h2#t2) = (count-diff-state h1 h2) + (count-diff t1 t2)

```

fun check-simp:: trace-regex => trace-regex => bool

```

where check-simp trace1 trace2 = ((count-diff trace1 trace2) ≤ 1 ∧ length trace1
= length trace2)

```

fun enumerate-pairs :: nat list => (nat * nat) list **where**

```

enumerate-pairs [] = [] |
enumerate-pairs (x#xs) = map (λy. (x, y)) xs @ enumerate-pairs xs

```

```

fun enum-pairs:: 'a list ⇒ (nat * nat) list
  where enum-pairs L = enumerate-pairs [0 ..< length L]

```

```

fun remove-element-at-index:: nat ⇒ 'a list ⇒ 'a list
  where remove-element-at-index n L = (take n L)@(drop (n+1) L)

```

This assumes (fst h) < (snd h)

```

fun update-L:: WEST-regex ⇒ (nat × nat) ⇒ nat ⇒ WEST-regex
  where update-L L h num-vars =
    (remove-element-at-index (fst h) (remove-element-at-index (snd h) L))@[WEST-simp-trace
    (L!(fst h)) (L!(snd h)) num-vars]

```

Defining and Proving Termination of WEST-simp **lemma** *length-enumerate-pairs*:
shows length (*enumerate-pairs* L) ≤ (length L) ^ 2
 ⟨proof⟩

lemma *length-enum-pairs*:
shows length (*enum-pairs* L) ≤ (length L) ^ 2
 ⟨proof⟩

lemma *enumerate-pairs-fact*:
assumes ∀ i j. (i < j ∧ i < length L ∧ j < length L) ⟶ (L!i) < (L!j)
shows ∀ pair ∈ set (*enumerate-pairs* L). (fst pair) < (snd pair)
 ⟨proof⟩

lemma *enum-pairs-fact*:
shows ∀ pair ∈ set (*enum-pairs* L). (fst pair) < (snd pair)
 ⟨proof⟩

lemma *enum-pairs-bound-snd*:
assumes pair ∈ set (*enumerate-pairs* L)
shows (snd pair) ∈ set L
 ⟨proof⟩

lemma *enum-pairs-bound*:
shows ∀ pair ∈ set (*enum-pairs* L). (snd pair) < length L
 ⟨proof⟩

lemma *WEST-simp-termination1-bound*:
fixes a::nat
shows a ^ 3 + a ^ 2 < (a+1) ^ 3
 ⟨proof⟩

lemma *WEST-simp-termination1*:
fixes L::WEST-regex

```

assumes  $\neg (idx\text{-}pairs \neq enum\text{-}pairs\ L \vee length\ idx\text{-}pairs \leq i)$ 
assumes  $check\text{-}simp\ (L\ !\ fst\ (idx\text{-}pairs\ !\ i))\ (L\ !\ snd\ (idx\text{-}pairs\ !\ i))$ 
assumes  $x = update\text{-}L\ L\ (idx\text{-}pairs\ !\ i)\ num\text{-}vars$ 
shows  $((x, enum\text{-}pairs\ x, 0, num\text{-}vars), L, idx\text{-}pairs, i, num\text{-}vars)$ 
 $\in measure\ (\lambda(L, idx\text{-}list, i, num\text{-}vars). length\ L \wedge 3 + length\ idx\text{-}list - i)$ 
<proof>

```

```

function WEST-simp-helper::  $WEST\text{-}regex \Rightarrow (nat \times nat)\ list \Rightarrow nat \Rightarrow nat \Rightarrow$ 
 $WEST\text{-}regex$ 
where WEST-simp-helper  $L\ idx\text{-}pairs\ i\ num\text{-}vars =$ 
 $(if\ (idx\text{-}pairs \neq enum\text{-}pairs\ L \vee i \geq length\ idx\text{-}pairs)\ then\ L\ else$ 
 $(if\ (check\text{-}simp\ (L!(fst\ (idx\text{-}pairs!i)))\ (L!(snd\ (idx\text{-}pairs!i))))\ then$ 
 $(let\ newL = update\text{-}L\ L\ (idx\text{-}pairs!i)\ num\text{-}vars\ in$ 
 $WEST\text{-}simp\text{-}helper\ newL\ (enum\text{-}pairs\ newL)\ 0\ num\text{-}vars)$ 
 $else\ WEST\text{-}simp\text{-}helper\ L\ idx\text{-}pairs\ (i+1)\ num\text{-}vars))$ 
<proof>
termination
<proof>

```

```

declare WEST-simp-helper.simps[simp del]

```

```

fun WEST-simp::  $WEST\text{-}regex \Rightarrow nat \Rightarrow WEST\text{-}regex$ 
where WEST-simp  $L\ num\text{-}vars =$ 
 $WEST\text{-}simp\text{-}helper\ L\ (enum\text{-}pairs\ L)\ 0\ num\text{-}vars$ 

```

```

value WEST-simp [[[S, S, One]], [[S, One, S]], [[S, S, Zero]]] 3
value WEST-simp [[[S, One]], [[One, S]], [[Zero, Zero]]] 2
value WEST-simp [[[One, One]], [[Zero, Zero]], [[One, Zero]], [[Zero, One]]] 2

```

1.3.3 AND and OR operations with WEST-simp

```

fun WEST-and-simp::  $WEST\text{-}regex \Rightarrow WEST\text{-}regex \Rightarrow nat \Rightarrow WEST\text{-}regex$ 
where WEST-and-simp  $L1\ L2\ num\text{-}vars = WEST\text{-}simp\ (WEST\text{-}and\ L1\ L2)$ 
 $num\text{-}vars$ 

```

```

fun WEST-or-simp::  $WEST\text{-}regex \Rightarrow WEST\text{-}regex \Rightarrow nat \Rightarrow WEST\text{-}regex$ 
where WEST-or-simp  $L1\ L2\ num\text{-}vars = WEST\text{-}simp\ (L1 @ L2)\ num\text{-}vars$ 

```

1.3.4 Useful Helper Functions

```

fun arbitrary-state::  $nat \Rightarrow state\text{-}regex$ 
where arbitrary-state  $num\text{-}vars = map\ (\lambda\ k. S)\ [0 ..< num\text{-}vars]$ 

fun arbitrary-trace::  $nat \Rightarrow nat \Rightarrow trace\text{-}regex$ 
where arbitrary-trace  $num\text{-}vars\ num\text{-}pad = map\ (\lambda\ k. (arbitrary\text{-}state\ num\text{-}vars))$ 
 $[0 ..< num\text{-}pad]$ 

fun shift::  $WEST\text{-}regex \Rightarrow nat \Rightarrow nat \Rightarrow WEST\text{-}regex$ 

```

where $\text{shift traceList num-vars num-pad} = \text{map } (\lambda \text{ trace. } (\text{arbitrary-trace num-vars num-pad}) @ \text{trace}) \text{ traceList}$

fun $\text{pad}:: \text{trace-regex} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{trace-regex}$
where $\text{pad trace num-vars num-pad} = \text{trace} @ (\text{arbitrary-trace num-vars num-pad})$

1.3.5 WEST Temporal Operations

fun $\text{WEST-global}:: \text{WEST-regex} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{WEST-regex}$
where $\text{WEST-global } L \ a \ b \ \text{num-vars} = ($
 $\text{if } (a = b) \text{ then } (\text{shift } L \ \text{num-vars } a)$
 $\text{else } (\text{if } (a < b) \text{ then } (\text{WEST-and-simp } (\text{shift } L \ \text{num-vars } b)$
 $\quad (\text{WEST-global } L \ a \ (b-1) \ \text{num-vars}) \ \text{num-vars})$
 $\text{else } []))$

fun $\text{WEST-future}:: \text{WEST-regex} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{WEST-regex}$
where $\text{WEST-future } L \ a \ b \ \text{num-vars} = ($
 $\text{if } (a = b)$
 $\text{then } (\text{shift } L \ \text{num-vars } a)$
 $\text{else } ($
 $\quad \text{if } (a < b)$
 $\quad \text{then } \text{WEST-or-simp } (\text{shift } L \ \text{num-vars } b) \ (\text{WEST-future } L \ a \ (b-1) \ \text{num-vars})$
 num-vars
 $\text{else } []))$

fun $\text{WEST-until}:: \text{WEST-regex} \Rightarrow \text{WEST-regex} \Rightarrow \text{nat} \Rightarrow$
 $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{WEST-regex}$
where $\text{WEST-until } L-\varphi \ L-\psi \ a \ b \ \text{num-vars} = ($
 $\text{if } (a=b)$
 $\text{then } (\text{WEST-global } L-\psi \ a \ a \ \text{num-vars})$
 $\text{else } ($
 $\quad \text{if } (a < b)$
 $\quad \text{then } \text{WEST-or-simp } (\text{WEST-until } L-\varphi \ L-\psi \ a \ (b-1) \ \text{num-vars})$
 $\quad (\text{WEST-and-simp } (\text{WEST-global } L-\varphi \ a \ (b-1) \ \text{num-vars})$
 $\quad \quad (\text{WEST-global } L-\psi \ b \ b \ \text{num-vars}) \ \text{num-vars}) \ \text{num-vars}$
 $\text{else } []))$

fun $\text{WEST-release-helper}:: \text{WEST-regex} \Rightarrow \text{WEST-regex} \Rightarrow$
 $\text{nat} \Rightarrow \text{nat} \Rightarrow \text{nat} \Rightarrow \text{WEST-regex}$
where $\text{WEST-release-helper } L-\varphi \ L-\psi \ a \ ub \ \text{num-vars} = ($
 $\text{if } (a=ub)$
 $\text{then } (\text{WEST-and-simp } (\text{WEST-global } L-\varphi \ a \ a \ \text{num-vars}) \ (\text{WEST-global } L-\psi \ a \ a$
 $\text{num-vars}) \ \text{num-vars})$
 $\text{else } ($
 $\quad \text{if } (a < ub)$
 $\quad \text{then } \text{WEST-or-simp } (\text{WEST-release-helper } L-\varphi \ L-\psi \ a \ (ub-1) \ \text{num-vars})$
 $\text{else } []))$

$(\text{WEST-and-simp } (\text{WEST-global } L\text{-}\psi \ a \ ub \ \text{num-vars})$
 $(\text{WEST-global } L\text{-}\varphi \ ub \ ub \ \text{num-vars}) \ \text{num-vars}) \ \text{num-vars}$
 $\text{else } []))$

fun *WEST-release*:: *WEST-regex* $\Rightarrow$ *WEST-regex* $\Rightarrow$ *nat*
 $\Rightarrow$ *nat* $\Rightarrow$ *nat* $\Rightarrow$ *WEST-regex*
where *WEST-release* *L-φ L-ψ a b num-vars* = (
if (*b* > *a*)
then (*WEST-or-simp* (*WEST-global* *L-ψ a b num-vars*) (*WEST-release-helper*
L-φ L-ψ a (b-1) num-vars) *num-vars*)
else (*WEST-global* *L-ψ a b num-vars*))

1.3.6 WEST recursive reg Function

lemma *exhaustive*:

shows $\bigwedge x::\text{nat mltl} \times \text{nat}. \bigwedge P::\text{bool}. (\bigwedge \text{num-vars}::\text{nat}. x = (\text{True-mltl}, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \text{num-vars}::\text{nat}. x = (\text{False-mltl}, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge p \text{ num-vars}::\text{nat}. x = (\text{Prop-mltl } p, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge p \text{ num-vars}::\text{nat}. x = (\text{Not-mltl } (\text{Prop-mltl } p), \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \psi \text{ num-vars}. x = (\text{Or-mltl } \varphi \psi, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \psi \text{ num-vars}. x = (\text{And-mltl } \varphi \psi, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \ a \ b \text{ num-vars}. x = (\text{Future-mltl } \varphi \ a \ b, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \ a \ b \text{ num-vars}. x = (\text{Global-mltl } \varphi \ a \ b, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \psi \ a \ b \text{ num-vars}. x = (\text{Until-mltl } \varphi \psi \ a \ b, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \psi \ a \ b \text{ num-vars}. x = (\text{Release-mltl } \varphi \psi \ a \ b, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \text{num-vars}. x = (\text{Not-mltl } \text{True-mltl}, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \text{num-vars}. x = (\text{Not-mltl } \text{False-mltl}, \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \psi \text{ num-vars}. x = (\text{Not-mltl } (\text{And-mltl } \varphi \psi), \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \psi \text{ num-vars}. x = (\text{Not-mltl } (\text{Or-mltl } \varphi \psi), \text{num-vars}) \Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \ a \ b \text{ num-vars}. x = (\text{Not-mltl } (\text{Future-mltl } \varphi \ a \ b), \text{num-vars}) \Rightarrow P)$
 $\Rightarrow$
 $(\bigwedge \varphi \ a \ b \text{ num-vars}. x = (\text{Not-mltl } (\text{Global-mltl } \varphi \ a \ b), \text{num-vars}) \Rightarrow P)$
 $\Rightarrow$
 $(\bigwedge \varphi \psi \ a \ b \text{ num-vars}. x = (\text{Not-mltl } (\text{Until-mltl } \varphi \psi \ a \ b), \text{num-vars}) \Rightarrow$
 $P) \Rightarrow$
 $(\bigwedge \varphi \psi \ a \ b \text{ num-vars}. x = (\text{Not-mltl } (\text{Release-mltl } \varphi \psi \ a \ b), \text{num-vars})$
 $\Rightarrow P) \Rightarrow$
 $(\bigwedge \varphi \text{ num-vars}. x = (\text{Not-mltl } (\text{Not-mltl } \varphi), \text{num-vars}) \Rightarrow P) \Rightarrow P$
 $\langle \text{proof} \rangle$

fun *WEST-termination-measure*:: (*nat*) *mltl* $\Rightarrow$ *nat*
where *WEST-termination-measure* *True_m* = 1
 $|$ *WEST-termination-measure* (*Not_m True_m*) = 4
 $|$ *WEST-termination-measure* *False_m* = 1
 $|$ *WEST-termination-measure* (*Not_m False_m*) = 4
 $|$ *WEST-termination-measure* (*Prop_m (p)*) = 1
 $|$ *WEST-termination-measure* (*Not_m (Prop_m (p))*) = 4

| *WEST-termination-measure* ($\varphi \text{ Or}_m \psi$) = 1 + (*WEST-termination-measure* φ) + (*WEST-termination-measure* ψ)
 | *WEST-termination-measure* ($\varphi \text{ And}_m \psi$) = 1 + (*WEST-termination-measure* φ) + (*WEST-termination-measure* ψ)
 | *WEST-termination-measure* ($F_m [a,b] \varphi$) = 1 + (*WEST-termination-measure* φ)
 | *WEST-termination-measure* ($G_m [a,b] \varphi$) = 1 + (*WEST-termination-measure* φ)
 | *WEST-termination-measure* ($\varphi \text{ U}_m[a,b] \psi$) = 1 + (*WEST-termination-measure* φ) + (*WEST-termination-measure* ψ)
 | *WEST-termination-measure* ($\varphi \text{ R}_m[a,b] \psi$) = 1 + (*WEST-termination-measure* φ) + (*WEST-termination-measure* ψ)
 | *WEST-termination-measure* ($\text{Not}_m (\varphi \text{ Or}_m \psi)$) = 1 + 3 * (*WEST-termination-measure* ($\varphi \text{ Or}_m \psi$))
 | *WEST-termination-measure* ($\text{Not}_m (\varphi \text{ And}_m \psi)$) = 1 + 3 * (*WEST-termination-measure* ($\varphi \text{ And}_m \psi$))
 | *WEST-termination-measure* ($\text{Not}_m (F_m[a,b] \varphi)$) = 1 + 3 * (*WEST-termination-measure* ($F_m[a,b] \varphi$))
 | *WEST-termination-measure* ($\text{Not}_m (G_m[a,b] \varphi)$) = 1 + 3 * (*WEST-termination-measure* ($G_m[a,b] \varphi$))
 | *WEST-termination-measure* ($\text{Not}_m (\varphi \text{ U}_m[a,b] \psi)$) = 1 + 3 * (*WEST-termination-measure* ($\varphi \text{ U}_m[a,b] \psi$))
 | *WEST-termination-measure* ($\text{Not}_m (\varphi \text{ R}_m[a,b] \psi)$) = 1 + 3 * (*WEST-termination-measure* ($\varphi \text{ R}_m[a,b] \psi$))
 | *WEST-termination-measure* ($\text{Not}_m (\text{Not}_m \varphi)$) = 1 + 3 * (*WEST-termination-measure* ($\text{Not}_m \varphi$))

lemma *WEST-termination-measure-not:*

fixes $\varphi::(\text{nat}) \text{ mltl}$
shows *WEST-termination-measure* ($\text{Not-mltl } \varphi$) = 1 + 3 * (*WEST-termination-measure* φ)
 <proof>

function *WEST-reg-aux*:: ($\text{nat}) \text{ mltl} \Rightarrow \text{nat} \Rightarrow \text{WEST-reg-expr}$

where *WEST-reg-aux* $\text{True}_m \text{ num-vars}$ = $[[(\text{map } (\lambda j. S) [0 ..< \text{num-vars}])]]$
 | *WEST-reg-aux* $\text{False}_m \text{ num-vars}$ = []
 | *WEST-reg-aux* ($\text{Prop}_m (p)$) num-vars = $[[(\text{map } (\lambda j. (\text{if } (p=j) \text{ then One else } S)) [0 ..< \text{num-vars}])]]$
 | *WEST-reg-aux* ($\text{Not}_m (\text{Prop}_m (p))$) num-vars = $[[(\text{map } (\lambda j. (\text{if } (p=j) \text{ then Zero else } S)) [0 ..< \text{num-vars}])]]$
 | *WEST-reg-aux* ($\varphi \text{ Or}_m \psi$) num-vars = *WEST-or-simp* (*WEST-reg-aux* $\varphi \text{ num-vars}$) (*WEST-reg-aux* $\psi \text{ num-vars}$) num-vars
 | *WEST-reg-aux* ($\varphi \text{ And}_m \psi$) num-vars = (*WEST-and-simp* (*WEST-reg-aux* $\varphi \text{ num-vars}$) (*WEST-reg-aux* $\psi \text{ num-vars}$) num-vars)
 | *WEST-reg-aux* ($F_m[a,b] \varphi$) num-vars = (*WEST-future* (*WEST-reg-aux* $\varphi \text{ num-vars}$) $a \ b \text{ num-vars}$)
 | *WEST-reg-aux* ($G_m[a,b] \varphi$) num-vars = (*WEST-global* (*WEST-reg-aux* $\varphi \text{ num-vars}$) $a \ b \text{ num-vars}$)

| $WEST\text{-}reg\text{-}aux (\varphi U_m[a,b] \psi) num\text{-}vars = (WEST\text{-}until (WEST\text{-}reg\text{-}aux \varphi num\text{-}vars) (WEST\text{-}reg\text{-}aux \psi num\text{-}vars) a b num\text{-}vars)$
 | $WEST\text{-}reg\text{-}aux (\varphi R_m[a,b] \psi) num\text{-}vars = WEST\text{-}release (WEST\text{-}reg\text{-}aux \varphi num\text{-}vars) (WEST\text{-}reg\text{-}aux \psi num\text{-}vars) a b num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m True_m) num\text{-}vars = WEST\text{-}reg\text{-}aux False_m num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m False_m) num\text{-}vars = WEST\text{-}reg\text{-}aux True_m num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m (\varphi And_m \psi)) num\text{-}vars = WEST\text{-}reg\text{-}aux ((Not_m \varphi) Or_m (Not_m \psi)) num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m (\varphi Or_m \psi)) num\text{-}vars = WEST\text{-}reg\text{-}aux ((Not_m \varphi) And_m (Not_m \psi)) num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m (F_m[a,b] \varphi)) num\text{-}vars = WEST\text{-}reg\text{-}aux (G_m[a,b] (Not_m \varphi)) num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m (G_m[a,b] \varphi)) num\text{-}vars = WEST\text{-}reg\text{-}aux (F_m[a,b] (Not_m \varphi)) num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m (\varphi U_m[a,b] \psi)) num\text{-}vars = WEST\text{-}reg\text{-}aux ((Not_m \varphi) R_m[a,b] (Not_m \psi)) num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m (\varphi R_m[a,b] \psi)) num\text{-}vars = WEST\text{-}reg\text{-}aux ((Not_m \varphi) U_m[a,b] (Not_m \psi)) num\text{-}vars$
 | $WEST\text{-}reg\text{-}aux (Not_m (Not_m \varphi)) num\text{-}vars = WEST\text{-}reg\text{-}aux \varphi num\text{-}vars$
 $\langle proof \rangle$
termination
 $\langle proof \rangle$

fun $WEST\text{-}num\text{-}vars:: (nat) mltl \Rightarrow nat$
where $WEST\text{-}num\text{-}vars True_m = 1$
 | $WEST\text{-}num\text{-}vars False_m = 1$
 | $WEST\text{-}num\text{-}vars (Prop_m (p)) = p+1$
 | $WEST\text{-}num\text{-}vars (Not_m \varphi) = WEST\text{-}num\text{-}vars \varphi$
 | $WEST\text{-}num\text{-}vars (\varphi And_m \psi) = Max \{(WEST\text{-}num\text{-}vars \varphi), (WEST\text{-}num\text{-}vars \psi)\}$
 | $WEST\text{-}num\text{-}vars (\varphi Or_m \psi) = Max \{(WEST\text{-}num\text{-}vars \varphi), (WEST\text{-}num\text{-}vars \psi)\}$
 | $WEST\text{-}num\text{-}vars (F_m[a,b] \varphi) = WEST\text{-}num\text{-}vars \varphi$
 | $WEST\text{-}num\text{-}vars (G_m[a,b] \varphi) = WEST\text{-}num\text{-}vars \varphi$
 | $WEST\text{-}num\text{-}vars (\varphi U_m[a,b] \psi) = Max \{(WEST\text{-}num\text{-}vars \varphi), (WEST\text{-}num\text{-}vars \psi)\}$
 | $WEST\text{-}num\text{-}vars (\varphi R_m[a,b] \psi) = Max \{(WEST\text{-}num\text{-}vars \varphi), (WEST\text{-}num\text{-}vars \psi)\}$

fun $WEST\text{-}reg:: (nat) mltl \Rightarrow WEST\text{-}regex$
where $WEST\text{-}reg F = (let nnf\text{-}F = convert\text{-}nnf F in WEST\text{-}reg\text{-}aux nnf\text{-}F (WEST\text{-}num\text{-}vars F))$

1.3.7 Adding padding

fun $pad\text{-}WEST\text{-}reg:: nat mltl \Rightarrow WEST\text{-}regex$
where $pad\text{-}WEST\text{-}reg \varphi = (let unpadding = WEST\text{-}reg \varphi in$

```

      (let complen = complen-mltl  $\varphi$  in
      (let num-vars = WEST-num-vars  $\varphi$  in
      (map ( $\lambda L$ . (if (length  $L$  < complen) then (pad  $L$  num-vars
      (complen-(length  $L$ )) else  $L$ )) unpadded)))

fun simp-pad-WEST-reg:: nat mltl  $\Rightarrow$  WEST-regex
  where simp-pad-WEST-reg  $\varphi$  = WEST-simp (pad-WEST-reg  $\varphi$ ) (WEST-num-vars
 $\varphi$ )

```

2 Some examples and Code Export

Base cases

```

value WEST-reg Truem
value WEST-reg Falsem
value WEST-reg (Propm (1))
value WEST-reg (Notm (Propm (0)))

```

Test cases for recursion

```

value WEST-reg ((Notm (Propm (0))) Andm (Propm (1)))
value WEST-reg (Fm[0,2] (Propm (1)))
value WEST-reg ((Notm (Propm (0))) Orm (Propm (0)))

value pad-WEST-reg ((Propm (0)) Um[0,2] (Propm (0)))
value simp-pad-WEST-reg ((Prop-mltl 0) Um[0,2] (Prop-mltl 0))

```

```

export-code WEST-reg in Haskell module-name WEST
export-code simp-pad-WEST-reg in Haskell module-name WEST-simp-pad

```

end

3 WEST Proofs

theory WEST-Proofs

imports WEST-Algorithms

begin

3.1 Useful Definitions

```

definition trace-of-vars::trace  $\Rightarrow$  nat  $\Rightarrow$  bool
  where trace-of-vars trace num-vars = (
   $\forall k$ . ( $k$  < (length trace)  $\longrightarrow$  ( $\forall p \in (\text{trace}!k)$ .  $p$  < num-vars)))

```

```

definition state-regex-of-vars::state-regex  $\Rightarrow$  nat  $\Rightarrow$  bool
  where state-regex-of-vars state num-vars = ((length state) = num-vars)

```

definition *trace-regex-of-vars*::*trace-regex* $\Rightarrow$ *nat* $\Rightarrow$ *bool*
where *trace-regex-of-vars* *trace* *num-vars* =
 $(\forall i < (\text{length } \text{trace}). \text{length } (\text{trace}!i) = \text{num-vars})$

definition *WEST-regex-of-vars*::*WEST-regex* $\Rightarrow$ *nat* $\Rightarrow$ *bool*
where *WEST-regex-of-vars* *traceList* *num-vars* =
 $(\forall k < \text{length } \text{traceList}. \text{trace-regex-of-vars } (\text{traceList}!k) \text{ num-vars})$

3.2 Proofs about Traces Matching Regular Expressions

value *match-regex* [{0::nat}, {0,1}, {}] []
lemma *arbitrary-regtrace-matches-any-trace*:
fixes *num-vars*::*nat*
fixes π ::*trace*
assumes π -of-*num-vars*: *trace-of-vars* π *num-vars*
shows *match-regex* π []
 $\langle \text{proof} \rangle$

lemma *WEST-and-state-diff lengths-is-none*:
assumes $\text{length } s1 \neq \text{length } s2$
shows *WEST-and-state* *s1* *s2* = *None*
 $\langle \text{proof} \rangle$

3.3 Facts about the WEST and operator

3.3.1 Commutative

lemma *WEST-and-bitwise-commutative*:
fixes *b1* *b2*::*WEST-bit*
shows *WEST-and-bitwise* *b1* *b2* = *WEST-and-bitwise* *b2* *b1*
 $\langle \text{proof} \rangle$

fun *remove-option-type-bit*:: *WEST-bit* *option* $\Rightarrow$ *WEST-bit*
where *remove-option-type-bit* (*Some* *i*) = *i*
| *remove-option-type-bit* - = *S*

lemma *WEST-and-state-commutative*:
fixes *s1* *s2*::*state-regex*
assumes *same-len*: $\text{length } s1 = \text{length } s2$
shows *WEST-and-state* *s1* *s2* = *WEST-and-state* *s2* *s1*
 $\langle \text{proof} \rangle$

lemma *WEST-and-trace-commutative*:
fixes *num-vars*::*nat*
fixes *regtrace1*::*trace-regex*
fixes *regtrace2*::*trace-regex*
assumes *regtrace1*-of-*num-vars*: *trace-regex-of-vars* *regtrace1* *num-vars*

assumes *regtrace2-of-num-vars: trace-regex-of-vars regtrace2 num-vars*
shows $(\text{WEST-and-trace } \text{regtrace1 } \text{regtrace2}) = (\text{WEST-and-trace } \text{regtrace2 } \text{regtrace1})$
 $\langle \text{proof} \rangle$

lemma *WEST-and-helper-subset:*
shows $\text{set } (\text{WEST-and-helper } h \ L) \subseteq \text{set } (\text{WEST-and-helper } h \ (a \ \# \ L))$
 $\langle \text{proof} \rangle$

lemma *WEST-and-helper-set-member-converse:*
fixes *regtrace h::trace-regex*
fixes *L::WEST-regex*
assumes *assumption: $(\exists \text{ loc. } \text{loc} < \text{length } L \wedge (\exists \text{ sometrace. } \text{WEST-and-trace } h \ (L \ ! \ \text{loc}) = \text{Some } \text{sometrace} \wedge \text{regtrace} = \text{sometrace}))$*
shows $\text{regtrace} \in \text{set } (\text{WEST-and-helper } h \ L)$
 $\langle \text{proof} \rangle$

lemma *WEST-and-helper-set-member-forward:*
fixes *regtrace h::trace-regex*
fixes *L::WEST-regex*
assumes $\text{regtrace} \in \text{set } (\text{WEST-and-helper } h \ L)$
shows $(\exists \text{ loc. } \text{loc} < \text{length } L \wedge (\exists \text{ sometrace. } \text{WEST-and-trace } h \ (L \ ! \ \text{loc}) = \text{Some } \text{sometrace} \wedge \text{regtrace} = \text{sometrace}))$
 $\langle \text{proof} \rangle$

lemma *WEST-and-helper-set-member:*
fixes *regtrace h::trace-regex*
fixes *L::WEST-regex*
shows $\text{regtrace} \in \text{set } (\text{WEST-and-helper } h \ L) \longleftrightarrow$
 $(\exists \text{ loc. } \text{loc} < \text{length } L \wedge (\exists \text{ sometrace. } \text{WEST-and-trace } h \ (L \ ! \ \text{loc}) = \text{Some } \text{sometrace} \wedge \text{regtrace} = \text{sometrace}))$
 $\langle \text{proof} \rangle$

lemma *WEST-and-set-member-dir1:*
fixes *num-vars::nat*
fixes *L1::WEST-regex*
fixes *L2::WEST-regex*
assumes *L1-of-num-vars: WEST-regex-of-vars L1 num-vars*
assumes *L2-of-num-vars: WEST-regex-of-vars L2 num-vars*
assumes $\text{regtrace} \in \text{set } (\text{WEST-and } L1 \ L2)$
shows $(\exists \text{ loc1 } \text{loc2. } \text{loc1} < \text{length } L1 \wedge \text{loc2} < \text{length } L2 \wedge$
 $(\exists \text{ sometrace. } \text{WEST-and-trace } (L1 \ ! \ \text{loc1}) \ (L2 \ ! \ \text{loc2}) = \text{Some } \text{sometrace} \wedge \text{regtrace} = \text{sometrace}))$
 $\langle \text{proof} \rangle$

lemma *WEST-and-subset:*
shows $\text{set } (\text{WEST-and } T1 \ L2) \subseteq \text{set } (\text{WEST-and } (h1 \ \# \ T1) \ L2)$

$\langle proof \rangle$

lemma *WEST-and-set-member-dir2:*

fixes *num-vars::nat*
fixes *L1::WEST-regex*
fixes *L2::WEST-regex*
assumes *L1-of-num-vars: WEST-regex-of-vars L1 num-vars*
assumes *L2-of-num-vars: WEST-regex-of-vars L2 num-vars*
assumes *exists-locs: ($\exists$ loc1 loc2. loc1 < length L1 $\wedge$ loc2 < length L2 $\wedge$
($\exists$ sometrace. WEST-and-trace (L1 ! loc1) (L2 ! loc2) = Some sometrace $\wedge$
regtrace = sometrace))*
shows *regtrace $\in$ set (WEST-and L1 L2)* $\langle proof \rangle$

lemma *WEST-and-set-member:*

fixes *num-vars::nat*
fixes *L1::WEST-regex*
fixes *L2::WEST-regex*
assumes *L1-of-num-vars: WEST-regex-of-vars L1 num-vars*
assumes *L2-of-num-vars: WEST-regex-of-vars L2 num-vars*
shows *regtrace $\in$ set (WEST-and L1 L2) $\longleftrightarrow$
($\exists$ loc1 loc2. loc1 < length L1 $\wedge$ loc2 < length L2 $\wedge$
($\exists$ sometrace. WEST-and-trace (L1 ! loc1) (L2 ! loc2) = Some sometrace $\wedge$
regtrace = sometrace))*
 $\langle proof \rangle$

lemma *WEST-and-commutative-sets-member:*

fixes *num-vars::nat*
fixes *L1::WEST-regex*
fixes *L2::WEST-regex*
assumes *L1-of-num-vars: WEST-regex-of-vars L1 num-vars*
assumes *L2-of-num-vars: WEST-regex-of-vars L2 num-vars*
assumes *regtrace-in: regtrace $\in$ set (WEST-and L1 L2)*
shows *regtrace $\in$ set (WEST-and L2 L1)*
 $\langle proof \rangle$

lemma *WEST-and-commutative-sets:*

fixes *num-vars::nat*
fixes *L1::WEST-regex*
fixes *L2::WEST-regex*
assumes *L1-of-num-vars: WEST-regex-of-vars L1 num-vars*
assumes *L2-of-num-vars: WEST-regex-of-vars L2 num-vars*
shows *set (WEST-and L1 L2) = set (WEST-and L2 L1)*
 $\langle proof \rangle$

lemma *WEST-and-commutative:*

fixes *num-vars::nat*
fixes *L1::WEST-regex*
fixes *L2::WEST-regex*
assumes *L1-of-num-vars: WEST-regex-of-vars L1 num-vars*

assumes $L2\text{-of-num-vars: WEST-regex-of-vars } L2 \text{ num-vars}$
shows $\text{regex-equiv } (WEST\text{-and } L1 \ L2) \ (WEST\text{-and } L2 \ L1)$
 $\langle \text{proof} \rangle$

3.3.2 Identity and Zero

lemma $WEST\text{-and-helper-identity:}$
shows $WEST\text{-and-helper } [] \ \text{trace} = \text{trace}$
 $\langle \text{proof} \rangle$

lemma $WEST\text{-and-identity: } WEST\text{-and } [[]] \ L = L$
 $\langle \text{proof} \rangle$

lemma $WEST\text{-and-zero: } WEST\text{-and } L \ [] = []$
 $\langle \text{proof} \rangle$

3.3.3 WEST-and-state

Well Defined **fun** $\text{advance-state:: state} \Rightarrow \text{state}$
where $\text{advance-state state} = \{x-1 \mid x. (x \in \text{state} \wedge x \neq 0)\}$

lemma $\text{advance-state-elt-bound:}$
fixes state::state
fixes num-vars::nat
assumes $\forall x \in \text{state}. x < \text{num-vars}$
shows $\forall x \in (\text{advance-state state}). x < (\text{num-vars}-1)$
 $\langle \text{proof} \rangle$

lemma $\text{advance-state-elt-member:}$
fixes state::state
fixes $x::\text{nat}$
assumes $x+1 \in \text{state}$
shows $x \in \text{advance-state state}$
 $\langle \text{proof} \rangle$

lemma $\text{advance-state-match-timestep:}$
fixes $h::WEST\text{-bit}$
fixes $t::\text{state-regex}$
fixes state::state
assumes $\text{match-timestep state } (h\#t)$
shows $\text{match-timestep } (\text{advance-state state}) \ t$
 $\langle \text{proof} \rangle$

lemma $WEST\text{-and-state-well-defined:}$
fixes num-vars::nat
fixes state::state
fixes $s1 \ s2::\text{state-regex}$
assumes $s1\text{-of-num-vars: state-regex-of-vars } s1 \ \text{num-vars}$

assumes $s2\text{-of-num-vars}$: $\text{state-regex-of-vars } s2 \text{ num-vars}$
assumes $\pi\text{-match-r1-r2}$: $\text{match-timestep state } s1 \wedge \text{match-timestep state } s2$
shows $\text{WEST-and-state } s1 \ s2 \neq \text{None}$
 $\langle \text{proof} \rangle$

Correct Forward lemma $\text{WEST-and-state-length}$:

fixes $s1 \ s2::\text{state-regex}$
assumes samelen : $\text{length } s1 = \text{length } s2$
assumes $r\text{-exists}$: $(\text{WEST-and-state } s1 \ s2) \neq \text{None}$
shows $\exists r. \text{length } r = \text{length } s1 \wedge \text{WEST-and-state } s1 \ s2 = \text{Some } r$
 $\langle \text{proof} \rangle$

lemma index-shift :

fixes $a::\text{WEST-bit}$
fixes $t::\text{state-regex}$
fixes $\text{state}::\text{state}$
assumes $(a = \text{One} \longrightarrow 0 \in \text{state}) \wedge (a = \text{Zero} \longrightarrow 0 \notin \text{state})$
assumes $\forall x < \text{length } t. ((t!x) = \text{One} \longrightarrow x + 1 \in \text{state}) \wedge ((t!x) = \text{Zero} \longrightarrow x + 1 \notin \text{state})$
shows $\forall x < \text{length } (a\#t). ((a\#t)!x = \text{One} \longrightarrow x \in \text{state}) \wedge ((a\#t)!x = \text{Zero} \longrightarrow x \notin \text{state})$
 $\langle \text{proof} \rangle$

lemma $\text{index-shift-reverse}$:

fixes $a::\text{WEST-bit}$
fixes $t::\text{state-regex}$
fixes $\text{state}::\text{state}$
assumes $\forall x < \text{length } (a\#t). ((a\#t)!x = \text{One} \longrightarrow x \in \text{state}) \wedge ((a\#t)!x = \text{Zero} \longrightarrow x \notin \text{state})$
shows $\forall x < \text{length } t. ((t!x) = \text{One} \longrightarrow x + 1 \in \text{state}) \wedge ((t!x) = \text{Zero} \longrightarrow x + 1 \notin \text{state})$
 $\langle \text{proof} \rangle$

lemma $\text{WEST-and-state-correct-forward}$:

fixes $\text{num-vars}::\text{nat}$
fixes $\text{state}::\text{state}$
fixes $s1 \ s2::\text{state-regex}$
assumes $s1\text{-of-num-vars}$: $\text{state-regex-of-vars } s1 \text{ num-vars}$
assumes $s2\text{-of-num-vars}$: $\text{state-regex-of-vars } s2 \text{ num-vars}$
assumes match-both : $\text{match-timestep state } s1 \wedge \text{match-timestep state } s2$
shows $\exists \text{somestate}. (\text{match-timestep state somestate}) \wedge (\text{WEST-and-state } s1 \ s2) = \text{Some somestate}$
 $\langle \text{proof} \rangle$

Correct Converse lemma $\text{WEST-and-state-indices}$:

fixes $s \ s1 \ s2::\text{state-regex}$

```

assumes WEST-and-state s1 s2 = Some s
assumes length s1 = length s2
assumes x < length s
shows Some (s!x) = WEST-and-bitwise (s1!x) (s2!x)
<proof>

```

```

lemma WEST-and-state-correct-converse-s1:
  fixes num-vars::nat
  fixes state::state
  fixes s1 s2:: state-regex
  assumes s1-of-num-vars: state-regex-of-vars s1 num-vars
  assumes s2-of-num-vars: state-regex-of-vars s2 num-vars
  assumes match-and: ∃ somestate. (match-timestep state somestate) ∧ (WEST-and-state
s1 s2) = Some somestate
  shows match-timestep state s1
<proof>

```

```

lemma WEST-and-state-correct-converse:
  fixes num-vars::nat
  fixes state::state
  fixes s1 s2:: state-regex
  assumes s1-of-num-vars: state-regex-of-vars s1 num-vars
  assumes s2-of-num-vars: state-regex-of-vars s2 num-vars
  assumes match-and: ∃ somestate. (match-timestep state somestate) ∧ (WEST-and-state
s1 s2) = Some somestate
  shows match-timestep state s1 ∧ match-timestep state s2
<proof>

```

```

lemma WEST-and-state-correct:
  fixes num-vars::nat
  fixes state::state
  fixes s1 s2:: state-regex
  assumes s1-of-num-vars: state-regex-of-vars s1 num-vars
  assumes s2-of-num-vars: state-regex-of-vars s2 num-vars
  shows (match-timestep state s1 ∧ match-timestep state s2) ⟷ (∃ somestate.
match-timestep state somestate ∧ (WEST-and-state s1 s2) = Some somestate)
<proof>

```

3.3.4 WEST-and-trace

```

Well Defined lemma WEST-and-trace-well-defined:
  fixes num-vars::nat
  fixes π::trace
  fixes r1 r2:: trace-regex
  assumes r1-of-num-vars: trace-regex-of-vars r1 num-vars
  assumes r2-of-num-vars: trace-regex-of-vars r2 num-vars
  assumes π-match-r1-r2: match-regex π r1 ∧ match-regex π r2
  shows WEST-and-trace r1 r2 ≠ None

```

$\langle \text{proof} \rangle$

Correct Forward lemma *WEST-and-trace-correct-forward-aux:*

assumes *match-regex* π ($h \# t$)

shows *match-timestep* ($\pi!0$) $h \wedge \text{match-regex}$ ($\text{drop } 1 \ \pi$) t

$\langle \text{proof} \rangle$

lemma *WEST-and-trace-correct-forward-aux-converse:*

assumes $\pi = hxi \# txi$

assumes *match-timestep* (hxi) h

assumes *match-regex* txi t

shows *match-regex* π ($h \# t$)

$\langle \text{proof} \rangle$

lemma *WEST-and-trace-correct-forward-empty-trace:*

fixes *num-vars::nat*

fixes $\pi::\text{trace}$

fixes $r1 \ r2::\text{trace-regex}$

assumes *r1-of-num-vars: trace-regex-of-vars* $r1$ *num-vars*

assumes *r2-of-num-vars: trace-regex-of-vars* $r2$ *num-vars*

assumes *match1: match-regex* $\sqsubseteq r1$

assumes *match2: match-regex* $\sqsubseteq r2$

shows $\exists \text{sometrace. match-regex } \sqsubseteq \text{sometrace} \wedge (\text{WEST-and-trace } r1 \ r2) = \text{Some}$

sometrace

$\langle \text{proof} \rangle$

lemma *WEST-and-trace-correct-forward-nonempty-trace:*

fixes *num-vars::nat*

fixes $\pi::\text{trace}$

fixes $r1 \ r2::\text{trace-regex}$

assumes *r1-of-num-vars: trace-regex-of-vars* $r1$ *num-vars*

assumes *r2-of-num-vars: trace-regex-of-vars* $r2$ *num-vars*

assumes *match-regex* $\pi \ r1 \wedge \text{match-regex } \pi \ r2$

assumes *length* $\pi > 0$

shows $\exists \text{sometrace. match-regex } \pi \ \text{sometrace} \wedge (\text{WEST-and-trace } r1 \ r2) = \text{Some}$

sometrace

$\langle \text{proof} \rangle$

lemma *WEST-and-trace-correct-forward:*

fixes *num-vars::nat*

fixes $\pi::\text{trace}$

fixes $r1 \ r2::\text{trace-regex}$

assumes *r1-of-num-vars: trace-regex-of-vars* $r1$ *num-vars*

assumes *r2-of-num-vars: trace-regex-of-vars* $r2$ *num-vars*

assumes *match-regex* $\pi \ r1 \wedge \text{match-regex } \pi \ r2$

shows $\exists \text{sometrace. match-regex } \pi \ \text{sometrace} \wedge (\text{WEST-and-trace } r1 \ r2) = \text{Some}$

sometrace

$\langle \text{proof} \rangle$

Correct Converse lemma *WEST-and-trace-nonempty-args*:
 fixes $h1\ h2::state\text{-}regex$
 fixes $t\ t1\ t2::trace\text{-}regex$
 assumes $WEST\text{-}and\text{-}trace\ (h1\ \# \ t1)\ (h2\ \# \ t2) = Some\ (h\ \# \ t)$
 shows $WEST\text{-}and\text{-}state\ h1\ h2 = Some\ h \wedge WEST\text{-}and\text{-}trace\ t1\ t2 = Some\ t$
 $\langle proof \rangle$

lemma *WEST-and-trace-lengths-r1*:
 assumes $trace\text{-}regex\text{-}of\text{-}vars\ r1\ n$
 assumes $trace\text{-}regex\text{-}of\text{-}vars\ r2\ n$
 assumes $(WEST\text{-}and\text{-}trace\ r1\ r2) = Some\ sometrace$
 shows $length\ sometrace \geq length\ r1$
 $\langle proof \rangle$

lemma *WEST-and-trace-lengths*:
 assumes $trace\text{-}regex\text{-}of\text{-}vars\ r1\ n$
 assumes $trace\text{-}regex\text{-}of\text{-}vars\ r2\ n$
 assumes $(WEST\text{-}and\text{-}trace\ r1\ r2) = Some\ sometrace$
 shows $length\ sometrace \geq length\ r1 \wedge length\ sometrace \geq length\ r2$
 $\langle proof \rangle$

lemma *WEST-and-trace-correct-converse-r1*:
 fixes $num\text{-}vars::nat$
 fixes $\pi::trace$
 fixes $r1\ r2::trace\text{-}regex$
 assumes $r1\text{-}of\text{-}num\text{-}vars: trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$
 assumes $r2\text{-}of\text{-}num\text{-}vars: trace\text{-}regex\text{-}of\text{-}vars\ r2\ num\text{-}vars$
 assumes $(\exists sometrace. match\text{-}regex\ \pi\ sometrace \wedge (WEST\text{-}and\text{-}trace\ r1\ r2) = Some\ sometrace)$
 shows $match\text{-}regex\ \pi\ r1$
 $\langle proof \rangle$

lemma *WEST-and-trace-correct-converse*:
 fixes $num\text{-}vars::nat$
 fixes $\pi::trace$
 fixes $r1\ r2::trace\text{-}regex$
 assumes $r1\text{-}of\text{-}num\text{-}vars: trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$
 assumes $r2\text{-}of\text{-}num\text{-}vars: trace\text{-}regex\text{-}of\text{-}vars\ r2\ num\text{-}vars$
 assumes $(\exists sometrace. match\text{-}regex\ \pi\ sometrace \wedge (WEST\text{-}and\text{-}trace\ r1\ r2) = Some\ sometrace)$
 shows $match\text{-}regex\ \pi\ r1 \wedge match\text{-}regex\ \pi\ r2$
 $\langle proof \rangle$

lemma *WEST-and-trace-correct*:
 fixes $num\text{-}vars::nat$
 fixes $\pi::trace$
 fixes $r1\ r2::trace\text{-}regex$
 assumes $r1\text{-}of\text{-}num\text{-}vars: trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$

assumes $r2\text{-of-num-vars: trace-regex-of-vars } r2 \text{ num-vars}$
shows $\text{match-regex } \pi \ r1 \wedge \text{match-regex } \pi \ r2 \longleftrightarrow (\exists \text{ sometrace. match-regex } \pi \text{ sometrace} \wedge (\text{WEST-and-trace } r1 \ r2) = \text{Some sometrace})$
 $\langle \text{proof} \rangle$

3.3.5 WEST-and correct

Correct Forward lemma *WEST-and-helper-subset-of-WEST-and:*

assumes $\text{List.member } L1 \ \text{elem}$
shows $\text{set } (\text{WEST-and-helper } \text{elem } (h2\#T2)) \subseteq \text{set } (\text{WEST-and } L1 \ (h2\#T2))$
 $\langle \text{proof} \rangle$

lemma *WEST-and-trace-element-of-WEST-and-helper:*

assumes $\text{List.member } L2 \ \text{elem2}$
assumes $(\text{WEST-and-trace } \text{elem1 } \text{elem2}) = \text{Some sometrace}$
shows $\text{some trace} \in \text{set } (\text{WEST-and-helper } \text{elem1 } L2)$
 $\langle \text{proof} \rangle$

lemma *index-of-L-in-L:*

assumes $i < \text{length } L$
shows $\text{List.member } L \ (L ! i)$
 $\langle \text{proof} \rangle$

lemma *WEST-and-indices:*

fixes $L1 \ L2:: \text{WEST-regex}$
fixes $\text{some trace}:: \text{trace-regex}$
assumes $\exists i1 \ i2. i1 < \text{length } L1 \wedge i2 < \text{length } L2 \wedge \text{WEST-and-trace } (L1 ! i1) \ (L2 ! i2) = \text{Some sometrace}$
shows $\exists i < \text{length } (\text{WEST-and } L1 \ L2). \text{WEST-and } L1 \ L2 ! i = \text{some trace}$
 $\langle \text{proof} \rangle$

lemma *WEST-and-correct-forward:*

fixes $n::\text{nat}$
fixes $\pi::\text{trace}$
fixes $L1 \ L2:: \text{WEST-regex}$
assumes $L1\text{-of-num-vars: WEST-regex-of-vars } L1 \ n$
assumes $L2\text{-of-num-vars: WEST-regex-of-vars } L2 \ n$
assumes $\text{match } \pi \ L1 \wedge \text{match } \pi \ L2$
shows $\text{match } \pi \ (\text{WEST-and } L1 \ L2)$
 $\langle \text{proof} \rangle$

Correct Converse lemma *WEST-and-correct-converse-L1:*

fixes $n::\text{nat}$
fixes $\pi::\text{trace}$
fixes $L1 \ L2:: \text{WEST-regex}$
assumes $L1\text{-of-num-vars: WEST-regex-of-vars } L1 \ n$
assumes $L2\text{-of-num-vars: WEST-regex-of-vars } L2 \ n$
assumes $\text{match } \pi \ (\text{WEST-and } L1 \ L2)$
shows $\text{match } \pi \ L1$

$\langle proof \rangle$

lemma *WEST-and-correct-converse:*

fixes $n::nat$
fixes $\pi::trace$
fixes $L1\ L2::WEST\text{-}regex$
assumes $L1\text{-of-num-vars: } WEST\text{-regex-of-vars } L1\ n$
assumes $L2\text{-of-num-vars: } WEST\text{-regex-of-vars } L2\ n$
assumes $match\ \pi\ (WEST\text{-and } L1\ L2)$
shows $match\ \pi\ L1 \wedge match\ \pi\ L2$

$\langle proof \rangle$

lemma *WEST-and-correct:*

fixes $\pi::trace$
fixes $L1\ L2::WEST\text{-}regex$
assumes $L1\text{-of-num-vars: } WEST\text{-regex-of-vars } L1\ n$
assumes $L2\text{-of-num-vars: } WEST\text{-regex-of-vars } L2\ n$
shows $match\ \pi\ L1 \wedge match\ \pi\ L2 \longleftrightarrow match\ \pi\ (WEST\text{-and } L1\ L2)$

$\langle proof \rangle$

3.4 Facts about the WEST or operator

lemma *WEST-or-correct:*

fixes $\pi::trace$
fixes $L1\ L2::WEST\text{-}regex$
shows $match\ \pi\ (L1 @ L2) \longleftrightarrow (match\ \pi\ L1) \vee (match\ \pi\ L2)$

$\langle proof \rangle$

3.5 Pad and Match Facts

lemma *shift-match-regex:*

assumes $length\ \pi \geq a$
assumes $match\text{-}regex\ \pi\ ((arbitrary\text{-}trace\ num\text{-}vars\ a) @ L)$
shows $match\text{-}regex\ (drop\ a\ \pi)\ (drop\ a\ ((arbitrary\text{-}trace\ num\text{-}vars\ a) @ L))$

$\langle proof \rangle$

lemma *match-regex:*

assumes $length\ \pi \geq a$
assumes $length\ L1 = a$
assumes $match\text{-}regex\ \pi\ (L1 @ L2)$
shows $match\text{-}regex\ (drop\ a\ \pi)\ (drop\ a\ (L1 @ L2))$

$\langle proof \rangle$

lemma *match-regex-converse:*

assumes $length\ \pi \geq a$
assumes $L1 = (arbitrary\text{-}trace\ num\text{-}vars\ a)$
assumes $match\text{-}regex\ (drop\ a\ \pi)\ (drop\ a\ (L1 @ L2))$

shows *match-regex* π ($L1 @ L2$)
 $\langle \text{proof} \rangle$

lemma *shift-match*:
assumes *length* $\pi \geq a$
assumes *match* π (*shift* L *num-vars* a)
shows *match* (*drop* a π) L
 $\langle \text{proof} \rangle$

lemma *shift-match-converse*:
assumes *length* $\pi \geq a$
assumes *match* (*drop* a π) L
shows *match* π (*shift* L *num-vars* a)
 $\langle \text{proof} \rangle$

lemma *pad-zero*:
shows *shift* $L2$ *num-vars* $0 = L2$
 $\langle \text{proof} \rangle$

3.6 Facts about WEST num vars

lemma *retrace-append*:
assumes *trace-regex-of-vars* $L1$ k
assumes *trace-regex-of-vars* $L2$ k
shows *trace-regex-of-vars* ($L1 @ L2$) k
 $\langle \text{proof} \rangle$

lemma *WEST-num-vars-subformulas*:
assumes $G \in \text{subformulas } F$
shows (*WEST-num-vars* F) $\geq$ *WEST-num-vars* G
 $\langle \text{proof} \rangle$

lemma *WEST-num-vars-nnf*:
shows (*WEST-num-vars* φ) = *WEST-num-vars* (*convert-nnf* φ)
 $\langle \text{proof} \rangle$

3.6.1 Facts about num vars for different WEST operators

lemma *length-WEST-and*:
assumes *length* $state1 = k$
assumes *length* $state2 = k$
assumes *WEST-and-state* $state1$ $state2 = \text{Some } state$
shows *length* $state = k$
 $\langle \text{proof} \rangle$

lemma *WEST-and-trace-num-vars*:
assumes *trace-regex-of-vars* $r1$ k
assumes *trace-regex-of-vars* $r2$ k

assumes (*WEST-and-trace* *r1 r2*) = *Some sometrace*
shows *trace-regex-of-vars sometrace k*
 <proof>

lemma *WEST-and-num-vars*:
assumes *WEST-regex-of-vars L1 k*
assumes *WEST-regex-of-vars L2 k*
shows *WEST-regex-of-vars (WEST-and L1 L2) k*
 <proof>

lemma *WEST-or-num-vars*:
assumes *L1-nv: WEST-regex-of-vars L1 k*
assumes *L2-nv: WEST-regex-of-vars L2 k*
shows *WEST-regex-of-vars (L1@L2) k*
 <proof>

lemma *regtraceList-cons-num-vars*:
assumes *trace-regex-of-vars h num-vars*
assumes *WEST-regex-of-vars T num-vars*
shows *WEST-regex-of-vars (h#T) num-vars*
 <proof>

lemma *WEST-simp-state-num-vars*:
assumes *length s1 = num-vars*
assumes *length s2 = num-vars*
shows *length (WEST-simp-state s1 s2) = num-vars*
 <proof>

lemma *WEST-get-state-length*:
assumes *trace-regex-of-vars r num-vars*
shows *length (WEST-get-state r k num-vars) = num-vars*
 <proof>

lemma *WEST-simp-trace-num-vars*:
assumes *trace-regex-of-vars r1 num-vars*
assumes *trace-regex-of-vars r2 num-vars*
shows *trace-regex-of-vars (WEST-simp-trace r1 r2 num-vars) num-vars*
 <proof>

lemma *remove-element-at-index-preserves-nv*:
assumes *i < length L*
assumes *WEST-regex-of-vars L num-vars*
shows *WEST-regex-of-vars (remove-element-at-index i L) num-vars*
 <proof>

lemma *update-L-length*:

assumes $h \in \text{set } (\text{enum-pairs } L)$

shows $\text{length } (\text{update-L } L \ h \ \text{num-var}) = \text{length } L - 1$

<proof>

lemma *update-L-preserves-num-vars*:

assumes *WEST-regex-of-vars* $L \ \text{num-var}$

assumes $h \in \text{set } (\text{enum-pairs } L)$

assumes $K = \text{update-L } L \ h \ \text{num-var}$

shows *WEST-regex-of-vars* $K \ \text{num-var}$

<proof>

lemma *WEST-simp-helper-can-simp*:

assumes $\text{simp-L} = \text{WEST-simp-helper } L \ (\text{enum-pairs } L) \ i \ \text{num-vars}$

assumes $\exists j. j < \text{length } (\text{enum-pairs } L) \wedge j \geq i \wedge$

$\text{check-simp } (L \ ! \ \text{fst } (\text{enum-pairs } L \ ! \ j))$

$(L \ ! \ \text{snd } (\text{enum-pairs } L \ ! \ j))$

assumes $\text{min-j} = \text{Min } \{j. j < \text{length } (\text{enum-pairs } L) \wedge j \geq i \wedge$

$\text{check-simp } (L \ ! \ \text{fst } (\text{enum-pairs } L \ ! \ j))$

$(L \ ! \ \text{snd } (\text{enum-pairs } L \ ! \ j))\}$

assumes $\text{newL} = \text{update-L } L \ (\text{enum-pairs } L \ ! \ \text{min-j}) \ \text{num-vars}$

assumes $i < \text{length } (\text{enum-pairs } L)$

shows $\text{simp-L} = \text{WEST-simp-helper } \text{newL} \ (\text{enum-pairs } \text{newL}) \ 0 \ \text{num-vars}$

<proof>

lemma *WEST-simp-helper-cant-simp*:

assumes $\text{simp-L} = \text{WEST-simp-helper } L \ (\text{enum-pairs } L) \ i \ \text{num-vars}$

assumes $\neg(\exists j. j < \text{length } (\text{enum-pairs } L) \wedge j \geq i \wedge$

$\text{check-simp } (L \ ! \ \text{fst } (\text{enum-pairs } L \ ! \ j))$

$(L \ ! \ \text{snd } (\text{enum-pairs } L \ ! \ j)))$

shows $\text{simp-L} = L$

<proof>

lemma *WEST-simp-helper-length*:

shows $\text{length } (\text{WEST-simp-helper } L \ (\text{enum-pairs } L) \ i \ \text{num-vars}) \leq \text{length } L$

<proof>

lemma *WEST-simp-helper-num-vars*:

assumes *WEST-regex-of-vars* $L \ \text{num-vars}$

shows *WEST-regex-of-vars* $(\text{WEST-simp-helper } L \ (\text{enum-pairs } L) \ i \ \text{num-vars})$

num-vars

<proof>

lemma *WEST-simp-num-vars*:

assumes *WEST-regex-of-vars* $L \ \text{num-vars}$

shows *WEST-regex-of-vars* $(\text{WEST-simp } L \ \text{num-vars}) \ \text{num-vars}$

<proof>

lemma *WEST-and-simp-num-vars*:
assumes *WEST-regex-of-vars L1 k*
assumes *WEST-regex-of-vars L2 k*
shows *WEST-regex-of-vars (WEST-and-simp L1 L2 k) k*
 $\langle proof \rangle$

lemma *WEST-or-simp-num-vars*:
assumes *WEST-regex-of-vars L1 k*
assumes *WEST-regex-of-vars L2 k*
shows *WEST-regex-of-vars (WEST-or-simp L1 L2 k) k*
 $\langle proof \rangle$

lemma *shift-num-vars*:
fixes *L::WEST-regex*
fixes *a k::nat*
assumes *WEST-regex-of-vars L k*
shows *WEST-regex-of-vars (shift L k a) k*
 $\langle proof \rangle$

lemma *WEST-future-num-vars*:
assumes *WEST-regex-of-vars L k*
assumes $a \leq b$
shows *WEST-regex-of-vars (WEST-future L a b k) k*
 $\langle proof \rangle$

lemma *WEST-global-num-vars*:
assumes *WEST-regex-of-vars L k*
assumes $a \leq b$
shows *WEST-regex-of-vars (WEST-global L a b k) k*
 $\langle proof \rangle$

lemma *WEST-until-num-vars*:
assumes *WEST-regex-of-vars L1 k*
assumes *WEST-regex-of-vars L2 k*
assumes $a \leq b$
shows *WEST-regex-of-vars (WEST-until L1 L2 a b k) k*
 $\langle proof \rangle$

lemma *WEST-release-helper-num-vars*:
assumes *WEST-regex-of-vars L1 k*
assumes *WEST-regex-of-vars L2 k*
assumes $a \leq b$

shows *WEST-regex-of-vars* (*WEST-release-helper* *L1 L2 a b k*) *k*
 ⟨*proof*⟩

lemma *WEST-release-num-vars*:
assumes *WEST-regex-of-vars* *L1 k*
assumes *WEST-regex-of-vars* *L2 k*
assumes $a \leq b$
shows *WEST-regex-of-vars* (*WEST-release* *L1 L2 a b k*) *k*
 ⟨*proof*⟩

lemma *WEST-reg-aux-num-vars*:
assumes *is-nnf*: $\exists \psi. F1 = (\text{convert-nnf } \psi)$
assumes $k \geq \text{WEST-num-vars } F1$
assumes *intervals-welldef* *F1*
shows *WEST-regex-of-vars* (*WEST-reg-aux* *F1 k*) *k*
 ⟨*proof*⟩

lemma *nnf-intervals-welldef*:
assumes *intervals-welldef* *F1*
shows *intervals-welldef* (*convert-nnf* *F1*)
 ⟨*proof*⟩

lemma *WEST-reg-num-vars*:
assumes *intervals-welldef* *F1*
shows *WEST-regex-of-vars* (*WEST-reg* *F1*) (*WEST-num-vars* *F1*)
 ⟨*proof*⟩

3.7 Correctness of WEST-simp

3.7.1 WEST-count-diff facts

lemma *count-diff-property-aux*:
assumes $k < \text{length } r1 \wedge k < \text{length } r2$
shows $\text{count-diff } r1 \ r2 \geq \text{count-diff-state } (r1 ! k) (r2 ! k)$
 ⟨*proof*⟩

lemma *count-diff-state-property*:
assumes $\text{count-diff-state } t1 \ t2 = 0$
assumes $ka < \text{length } t1 \wedge ka < \text{length } t2$
shows $t1 ! ka = t2 ! ka$
 ⟨*proof*⟩

lemma *count-diff-property*:
assumes $\text{count-diff } r1 \ r2 = 0$
assumes $k < \text{length } r1 \wedge k < \text{length } r2$
assumes $ka < \text{length } (r1 ! k) \wedge ka < \text{length } (r2 ! k)$
shows $r2 ! k ! ka = r1 ! k ! ka$
 ⟨*proof*⟩

lemma *count-nonS-trace-0-allS*:

assumes *length h = num-vars*

assumes *count-nonS-trace h = 0*

shows *h = map ($\lambda t. S$) [0..*num-vars*]*

<proof>

lemma *trace-tail-num-vars*:

assumes *trace-regex-of-vars (h # trace) num-vars*

shows *trace-regex-of-vars trace num-vars*

<proof>

lemma *count-diff-property-S-aux*:

assumes *count-diff trace [] = 0*

assumes *k < length trace*

assumes *trace-regex-of-vars trace num-vars*

assumes *1 ≤ num-vars*

shows *trace ! k = map ($\lambda t. S$) [0..*num-vars*]*

<proof>

lemma *count-diff-property-S*:

assumes *count-diff r1 r2 = 0*

assumes *k < length r1 ∧ length r2 ≤ k*

assumes *trace-regex-of-vars r1 num-vars*

assumes *num-vars ≥ 1*

assumes *ka < num-vars*

shows *r1 ! k = map ($\lambda t. S$) [0..*num-vars*]*

<proof>

lemma *count-diff-state-commutative*:

shows *count-diff-state e1 e2 = count-diff-state e2 e1*

<proof>

lemma *count-diff-commutative*:

shows *count-diff r1 r2 = count-diff r2 r1*

<proof>

lemma *count-diff-same-trace*:

shows *count-diff trace trace = 0*

<proof>

lemma *count-diff-state-0*:

assumes *count-diff-state h1 h2 = 0*

assumes *length h1 = length h2*

shows *h1 = h2*

<proof>

lemma *count-diff-state-1*:
assumes $\text{length } h1 = \text{length } h2$
assumes $\text{count-diff-state } h1 \ h2 = 1$
shows $\exists ka < \text{length } h1. h1!ka \neq h2!ka$
 $\langle \text{proof} \rangle$

lemma *count-diff-state-other-states*:
assumes $\text{count-diff-state } h1 \ h2 = 1$
assumes $\text{length } h1 = \text{length } h2$
assumes $h1!k \neq h2!k$
assumes $k < \text{length } h1$
shows $\forall i < \text{length } h1. k \neq i \longrightarrow h1!i = h2!i$
 $\langle \text{proof} \rangle$

lemma *count-diff-same-len*:
assumes $\text{trace-regex-of-vars } r1 \ \text{num-vars}$
assumes $\text{trace-regex-of-vars } r2 \ \text{num-vars}$
assumes $\text{count-diff } r1 \ r2 = 0$
assumes $\text{length } r1 = \text{length } r2$
shows $r1 = r2$
 $\langle \text{proof} \rangle$

lemma *count-diff-1*:
assumes $\text{count-diff } r1 \ r2 = 1$
assumes $\text{length } r1 = \text{length } r2$
assumes $\text{trace-regex-of-vars } r1 \ \text{num-vars}$
assumes $\text{trace-regex-of-vars } r2 \ \text{num-vars}$
shows $\exists k < \text{length } r1. \text{count-diff-state } (r1!k) \ (r2!k) = 1$
 $\langle \text{proof} \rangle$

lemma *count-diff-1-other-states*:
assumes $\text{count-diff } r1 \ r2 = 1$
assumes $\text{length } r1 = \text{length } r2$
assumes $\text{trace-regex-of-vars } r1 \ \text{num-vars}$
assumes $\text{trace-regex-of-vars } r2 \ \text{num-vars}$
assumes $\text{count-diff-state } (r1!k) \ (r2!k) = 1$
shows $\forall i < \text{length } r1. k \neq i \longrightarrow r1!i = r2!i$
 $\langle \text{proof} \rangle$

3.7.2 Orsimp-trace Facts

lemma *WEST-simp-bitwise-identity*:
assumes $b1 = b2$
shows $\text{WEST-simp-bitwise } b1 \ b2 = b1$
 $\langle \text{proof} \rangle$

lemma *WEST-simp-bitwise-commutative*:

shows $WEST\text{-}simp\text{-}bitwise\ b1\ b2 = WEST\text{-}simp\text{-}bitwise\ b2\ b1$
 $\langle proof \rangle$

lemma $WEST\text{-}simp\text{-}state\text{-}commutative$:
assumes $length\ s1 = num\text{-}vars$
assumes $length\ s2 = num\text{-}vars$
shows $WEST\text{-}simp\text{-}state\ s1\ s2 = WEST\text{-}simp\text{-}state\ s2\ s1$
 $\langle proof \rangle$

lemma $WEST\text{-}simp\text{-}trace\text{-}commutative$:
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r2\ num\text{-}vars$
shows $WEST\text{-}simp\text{-}trace\ r1\ r2\ num\text{-}vars = WEST\text{-}simp\text{-}trace\ r2\ r1\ num\text{-}vars$
 $\langle proof \rangle$

lemma $WEST\text{-}simp\text{-}trace\text{-}identity$:
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r2\ num\text{-}vars$
assumes $count\text{-}diff\ r1\ r2 = 0$
assumes $length\ r1 \geq length\ r2$
shows $WEST\text{-}simp\text{-}trace\ r1\ r2\ num\text{-}vars = r1$
 $\langle proof \rangle$

lemma $WEST\text{-}simp\text{-}trace\text{-}length$:
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r2\ num\text{-}vars$
assumes $length\ r1 = length\ r2$
shows $length\ (WEST\text{-}simp\text{-}trace\ r1\ r2\ num\text{-}vars) = length\ r1$
 $\langle proof \rangle$

3.7.3 WEST-orsimp-trace-correct

lemma $WEST\text{-}simp\text{-}trace\text{-}correct\text{-}forward$:
assumes $check\text{-}simp\ r1\ r2$
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r2\ num\text{-}vars$
assumes $match\text{-}regex\ \pi\ (WEST\text{-}simp\text{-}trace\ r1\ r2\ num\text{-}vars)$
shows $match\text{-}regex\ \pi\ r1 \vee match\text{-}regex\ \pi\ r2$
 $\langle proof \rangle$

lemma $WEST\text{-}simp\text{-}trace\text{-}correct\text{-}converse$:
assumes $check\text{-}simp\ r1\ r2$
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r1\ num\text{-}vars$
assumes $trace\text{-}regex\text{-}of\text{-}vars\ r2\ num\text{-}vars$
assumes $match\text{-}regex\ \pi\ r1 \vee match\text{-}regex\ \pi\ r2$
shows $match\text{-}regex\ \pi\ (WEST\text{-}simp\text{-}trace\ r1\ r2\ num\text{-}vars)$

$\langle \text{proof} \rangle$

lemma *WEST-simp-trace-correct:*

assumes *check-simp* *r1* *r2*

assumes *trace-regex-of-vars* *r1* *num-vars*

assumes *trace-regex-of-vars* *r2* *num-vars*

shows *match-regex* π (*WEST-simp-trace* *r1* *r2* *num-vars*) $\longleftrightarrow$ *match-regex* π *r1*
 $\vee$ *match-regex* π *r2*

$\langle \text{proof} \rangle$

3.7.4 Simp-helper Correct

lemma *WEST-simp-helper-can-simp-bound:*

assumes *simp-L* = *WEST-simp-helper* *L* (*enum-pairs* *L*) *i* *num-vars*

assumes $\exists j. j < \text{length } (\text{enum-pairs } L) \wedge j \geq i \wedge$
check-simp (*L* ! *fst* (*enum-pairs* *L* ! *j*))
(*L* ! *snd* (*enum-pairs* *L* ! *j*))

assumes *i* < *length* (*enum-pairs* *L*)

shows *length simp-L* < *length L*

$\langle \text{proof} \rangle$

lemma *WEST-simp-helper-same-length:*

assumes *WEST-regex-of-vars* *L* *num-vars*

assumes *K* = *WEST-simp-helper* *L* (*enum-pairs* *L*) 0 *num-vars*

assumes *length K* = *length L*

shows *L* = *K*

$\langle \text{proof} \rangle$

lemma *WEST-simp-helper-less-length:*

assumes *WEST-regex-of-vars* *L* *num-vars*

assumes *length K* < *length L*

assumes *K* = *WEST-simp-helper* *L* (*enum-pairs* *L*) 0 *num-vars*

shows $\exists \text{min-j.}$

(*min-j* < *length* (*enum-pairs* *L*) $\wedge$

K =

WEST-simp-helper (*update-L* *L* (*enum-pairs* *L* ! *min-j*) *num-vars*)

(*enum-pairs*

(*update-L* *L* (*enum-pairs* *L* ! *min-j*) *num-vars*))

0 *num-vars*

$\wedge$ *check-simp* (*L* ! *fst* (*enum-pairs* *L* ! *min-j*)) (*L* ! *snd* (*enum-pairs* *L* !

min-j)))

$\langle \text{proof} \rangle$

lemma *remove-element-at-index-subset:*

fixes *i::nat*

assumes *i* < *length L*

shows *set* (*remove-element-at-index* *i* *L*) $\subseteq$ *set L*

<proof>

lemma *WEST-simp-helper-correct-forward:*
 assumes *WEST-regex-of-vars* *L num-vars*
 assumes *match* π *K*
 assumes $K = \text{WEST-simp-helper } L \text{ (enum-pairs } L) \ 0 \text{ num-vars}$
 shows *match* π *L*
<proof>

lemma *remove-element-at-index-fact:*
 assumes $j1 < j2$
 assumes $j2 < \text{length } L$
 assumes $i < \text{length } L$
 assumes $i \neq j1$
 assumes $i \neq j2$
 shows $L ! i$
 $\in \text{set (remove-element-at-index } j1 \text{ (remove-element-at-index } j2 \text{ } L))$
<proof>

lemma *update-L-match:*
 assumes *WEST-regex-of-vars* *L num-var*
 assumes *match* π *L*
 assumes $h \in \text{set (enum-pairs } L)$
 assumes *check-simp* ($L!(fst \ h)$) ($L!(snd \ h)$)
 shows *match* π (*update-L* *L* *h num-var*)
<proof>

lemma *WEST-simp-helper-correct-converse:*
 assumes *WEST-regex-of-vars* *L num-vars*
 assumes *match* π *L*
 assumes $K = \text{WEST-simp-helper } L \text{ (enum-pairs } L) \ i \text{ num-vars}$
 shows *match* π *K*
<proof>

3.7.5 WEST-simp Correct

lemma *simp-correct-forward:*
 assumes *WEST-regex-of-vars* *L num-vars*
 assumes *match* π (*WEST-simp* *L num-vars*)
 shows *match* π *L*
<proof>

lemma *simp-correct-converse:*
 assumes *WEST-regex-of-vars* *L num-vars*
 assumes *match* π *L*
 shows *match* π (*WEST-simp* *L num-vars*)

$\langle proof \rangle$

lemma *simp-correct*:

assumes *WEST-regex-of-vars* L *num-vars*
shows $match\ \pi\ (WEST\text{-}simp\ L\ num\text{-}vars) \longleftrightarrow match\ \pi\ L$
 $\langle proof \rangle$

3.8 Correctness of WEST-and-simp/WEST-or-simp

lemma *WEST-and-simp-correct*:

fixes $\pi::trace$
fixes $L1\ L2::WEST\text{-}regex$
assumes *L1-of-num-vars*: *WEST-regex-of-vars* $L1\ n$
assumes *L2-of-num-vars*: *WEST-regex-of-vars* $L2\ n$
shows $match\ \pi\ L1 \wedge match\ \pi\ L2 \longleftrightarrow match\ \pi\ (WEST\text{-}and\text{-}simp\ L1\ L2\ n)$
 $\langle proof \rangle$

lemma *WEST-or-simp-correct*:

fixes $\pi::trace$
fixes $L1\ L2::WEST\text{-}regex$
assumes *L1-of-num-vars*: *WEST-regex-of-vars* $L1\ n$
assumes *L2-of-num-vars*: *WEST-regex-of-vars* $L2\ n$
shows $match\ \pi\ L1 \vee match\ \pi\ L2 \longleftrightarrow match\ \pi\ (WEST\text{-}or\text{-}simp\ L1\ L2\ n)$
 $\langle proof \rangle$

3.9 Facts about the WEST future operator

lemma *WEST-future-correct-forward*:

assumes $\bigwedge \pi. (length\ \pi \geq complen\text{-}mltl\ F \longrightarrow (match\ \pi\ L \longleftrightarrow semantics\text{-}mltl\ \pi\ F))$
assumes *WEST-regex-of-vars* L *num-vars*
assumes *WEST-num-vars* $F \leq num\text{-}vars$
assumes $a \leq b$
assumes $length\ \pi \geq (complen\text{-}mltl\ F) + b$
assumes $match\ \pi\ (WEST\text{-}future\ L\ a\ b\ num\text{-}vars)$
shows $\pi \models_m (F_m\ [a, b]\ F)$
 $\langle proof \rangle$

lemma *WEST-future-correct-converse*:

assumes $\bigwedge \pi. (length\ \pi \geq complen\text{-}mltl\ F \longrightarrow (match\ \pi\ L \longleftrightarrow semantics\text{-}mltl\ \pi\ F))$
assumes *WEST-regex-of-vars* L *num-vars*
assumes *WEST-num-vars* $F \leq num\text{-}vars$
assumes $a \leq b$
assumes $length\ \pi \geq (complen\text{-}mltl\ F) + b$
assumes $\pi \models_m (Future\text{-}mltl\ a\ b\ F)$

shows $\text{match } \pi \text{ (WEST-future } L \text{ a b num-vars)}$
 $\langle \text{proof} \rangle$

lemma *WEST-future-correct*:

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F \longrightarrow (\text{match } \pi \text{ } L \longleftrightarrow \text{semantics-mltl } \pi \text{ } F))$
assumes *WEST-regex-of-vars* $L \text{ num-vars}$
assumes *WEST-num-vars* $F \leq \text{num-vars}$
assumes $a \leq b$
assumes $\text{length } \pi \geq (\text{complen-mltl } F) + b$
shows $\text{match } \pi \text{ (WEST-future } L \text{ a b num-vars)} \longleftrightarrow$
 $\text{semantics-mltl } \pi \text{ (Future-mltl } a \text{ b } F)$
 $\langle \text{proof} \rangle$

3.10 Facts about the WEST global operator

lemma *WEST-global-correct-forward*:

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F \longrightarrow (\text{match } \pi \text{ } L \longleftrightarrow \text{semantics-mltl } \pi \text{ } F))$
assumes *WEST-regex-of-vars* $L \text{ num-vars}$
assumes *WEST-num-vars* $F \leq \text{num-vars}$
assumes $a \leq b$
assumes $\text{length } \pi \geq (\text{complen-mltl } F) + b$
assumes $\text{match } \pi \text{ (WEST-global } L \text{ a b num-vars)}$
shows $\text{semantics-mltl } \pi \text{ (Global-mltl } a \text{ b } F)$
 $\langle \text{proof} \rangle$

lemma *WEST-global-correct-converse*:

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F \longrightarrow (\text{match } \pi \text{ } L \longleftrightarrow \text{semantics-mltl } \pi \text{ } F))$
assumes *WEST-regex-of-vars* $L \text{ num-vars}$
assumes *WEST-num-vars* $F \leq \text{num-vars}$
assumes $a \leq b$
assumes $\text{length } \pi \geq (\text{complen-mltl } F) + b$
assumes $\text{semantics-mltl } \pi \text{ (Global-mltl } a \text{ b } F)$
shows $\text{match } \pi \text{ (WEST-global } L \text{ a b num-vars)}$
 $\langle \text{proof} \rangle$

lemma *WEST-global-correct*:

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F \longrightarrow (\text{match } \pi \text{ } L \longleftrightarrow \text{semantics-mltl } \pi \text{ } F))$
assumes *WEST-regex-of-vars* $L \text{ num-vars}$
assumes *WEST-num-vars* $F \leq \text{num-vars}$
assumes $a \leq b$
assumes $\text{length } \pi \geq (\text{complen-mltl } F) + b$

shows $\text{match } \pi \text{ (WEST-global } L \text{ a b num-vars)} \longleftrightarrow$
 $\text{semantics-mltl } \pi \text{ (Global-mltl a b F)}$
 $\langle \text{proof} \rangle$

3.11 Facts about the WEST until operator

lemma *WEST-until-correct-forward:*

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F1 \longrightarrow (\text{match } \pi \text{ L1} \longleftrightarrow \text{semantics-mltl } \pi \text{ F1}))$
assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F2 \longrightarrow (\text{match } \pi \text{ L2} \longleftrightarrow \text{semantics-mltl } \pi \text{ F2}))$
assumes *WEST-regex-of-vars* *L1 num-vars*
assumes *WEST-regex-of-vars* *L2 num-vars*
assumes *WEST-num-vars* *F1* $\leq$ *num-vars*
assumes *WEST-num-vars* *F2* $\leq$ *num-vars*
assumes $a \leq b$
assumes $\text{length } \pi \geq \text{complen-mltl (Until-mltl F1 a b F2)}$
assumes $\text{match } \pi \text{ (WEST-until L1 L2 a b num-vars)}$
shows $\text{semantics-mltl } \pi \text{ (Until-mltl F1 a b F2)}$
 $\langle \text{proof} \rangle$

lemma *WEST-until-correct-converse:*

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F1 \longrightarrow (\text{match } \pi \text{ L1} \longleftrightarrow \text{semantics-mltl } \pi \text{ F1}))$
assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F2 \longrightarrow (\text{match } \pi \text{ L2} \longleftrightarrow \text{semantics-mltl } \pi \text{ F2}))$
assumes *WEST-regex-of-vars* *L1 num-vars*
assumes *WEST-regex-of-vars* *L2 num-vars*
assumes *WEST-num-vars* *F1* $\leq$ *num-vars*
assumes *WEST-num-vars* *F2* $\leq$ *num-vars*
assumes $a \leq b$
assumes $\text{length } \pi \geq (\text{complen-mltl (Until-mltl F1 a b F2)})$
assumes $\text{semantics-mltl } \pi \text{ (Until-mltl F1 a b F2)}$
shows $\text{match } \pi \text{ (WEST-until L1 L2 a b num-vars)}$
 $\langle \text{proof} \rangle$

lemma *WEST-until-correct:*

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F1 \longrightarrow (\text{match } \pi \text{ L1} \longleftrightarrow \text{semantics-mltl } \pi \text{ F1}))$
assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F2 \longrightarrow (\text{match } \pi \text{ L2} \longleftrightarrow \text{semantics-mltl } \pi \text{ F2}))$
assumes *WEST-regex-of-vars* *L1 num-vars*
assumes *WEST-regex-of-vars* *L2 num-vars*
assumes *WEST-num-vars* *F1* $\leq$ *num-vars*
assumes *WEST-num-vars* *F2* $\leq$ *num-vars*
assumes $a \leq b$
assumes $\text{length } \pi \geq \text{complen-mltl (Until-mltl F1 a b F2)}$

shows $\text{match } \pi \text{ (WEST-until } L1 \ L2 \ a \ b \ \text{num-vars)} \longleftrightarrow$
 $\text{semantics-mltl } \pi \text{ (Until-mltl } F1 \ a \ b \ F2)$
 $\langle \text{proof} \rangle$

3.12 Facts about the WEST release Operator

lemma *WEST-release-correct-forward:*

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F1 \longrightarrow (\text{match } \pi \ L1 \longleftrightarrow \text{semantics-mltl } \pi \ F1))$
assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F2 \longrightarrow (\text{match } \pi \ L2 \longleftrightarrow \text{semantics-mltl } \pi \ F2))$
assumes *WEST-regex-of-vars* $L1 \ \text{num-vars}$
assumes *WEST-regex-of-vars* $L2 \ \text{num-vars}$
assumes *WEST-num-vars* $F1 \leq \text{num-vars}$
assumes *WEST-num-vars* $F2 \leq \text{num-vars}$
assumes $a \leq b$
assumes $\text{len: length } \pi \geq \text{complen-mltl (Release-mltl } F1 \ a \ b \ F2)$
assumes $\text{match } \pi \text{ (WEST-release } L1 \ L2 \ a \ b \ \text{num-vars)}$
shows $\text{semantics-mltl } \pi \text{ (Release-mltl } F1 \ a \ b \ F2)$
 $\langle \text{proof} \rangle$

lemma *WEST-release-correct-converse:*

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F1 \longrightarrow (\text{match } \pi \ L1 \longleftrightarrow \text{semantics-mltl } \pi \ F1))$
assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F2 \longrightarrow (\text{match } \pi \ L2 \longleftrightarrow \text{semantics-mltl } \pi \ F2))$
assumes *WEST-regex-of-vars* $L1 \ \text{num-vars}$
assumes *WEST-regex-of-vars* $L2 \ \text{num-vars}$
assumes *WEST-num-vars* $F1 \leq \text{num-vars}$
assumes *WEST-num-vars* $F2 \leq \text{num-vars}$
assumes $a \leq b$
assumes $\text{length } \pi \geq \text{complen-mltl (Release-mltl } F1 \ a \ b \ F2)$
assumes $\text{semantics-mltl } \pi \text{ (Release-mltl } F1 \ a \ b \ F2)$
shows $\text{match } \pi \text{ (WEST-release } L1 \ L2 \ a \ b \ \text{num-vars)}$
 $\langle \text{proof} \rangle$

lemma *WEST-release-correct:*

assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F1 \longrightarrow (\text{match } \pi \ L1 \longleftrightarrow \text{semantics-mltl } \pi \ F1))$
assumes $\bigwedge \pi. (\text{length } \pi \geq \text{complen-mltl } F2 \longrightarrow (\text{match } \pi \ L2 \longleftrightarrow \text{semantics-mltl } \pi \ F2))$
assumes *WEST-regex-of-vars* $L1 \ \text{num-vars}$
assumes *WEST-regex-of-vars* $L2 \ \text{num-vars}$
assumes *WEST-num-vars* $F1 \leq \text{num-vars}$
assumes *WEST-num-vars* $F2 \leq \text{num-vars}$
assumes $a \leq b$
assumes $\text{length } \pi \geq \text{complen-mltl (Release-mltl } F1 \ a \ b \ F2)$

shows *semantics-mltl* π (*Release-mltl* $F1 \ a \ b \ F2$) $\longleftrightarrow$ *match* π (*WEST-release* $L1 \ L2 \ a \ b \ \text{num-vars}$)
 <proof>

3.13 Top level result: Shows that WEST reg is correct

lemma *WEST-reg-aux-correct*:
assumes π -long-enough: $\text{length } \pi \geq \text{complen-mltl } F$
assumes *is-nnf*: $\exists \psi. F = (\text{convert-nnf } \psi)$
assumes φ -nv: *WEST-num-vars* $F \leq \text{num-vars}$
assumes *intervals-welldef* F
shows *match* π (*WEST-reg-aux* $F \ \text{num-vars}$) $\longleftrightarrow$ *semantics-mltl* $\pi \ F$
 <proof>

lemma *complen-convert-nnf*:
shows *complen-mltl* (*convert-nnf* φ) = *complen-mltl* φ
 <proof>

lemma *nnf-int-welldef*:
assumes *intervals-welldef* φ
shows *intervals-welldef* (*convert-nnf* φ)
 <proof>

lemma *WEST-correct*:
fixes $\varphi::(\text{nat}) \ \text{mltl}$
fixes $\pi::\text{trace}$
assumes *int-welldef*: *intervals-welldef* φ
assumes π -long-enough: $\text{length } \pi \geq \text{complen-mltl } (\text{convert-nnf } \varphi)$
shows *match* π (*WEST-reg* φ) $\longleftrightarrow$ *semantics-mltl* $\pi \ \varphi$
 <proof>

lemma *WEST-correct-v2*:
fixes $\varphi::(\text{nat}) \ \text{mltl}$
fixes $\pi::\text{trace}$
assumes *intervals-welldef* φ
assumes π -long-enough: $\text{length } \pi \geq \text{complen-mltl } \varphi$
shows *match* π (*WEST-reg* φ) $\longleftrightarrow$ *semantics-mltl* $\pi \ \varphi$
 <proof>

3.14 Top level result for padded version

lemma *WEST-correct-pad-aux*:
fixes $\varphi::(\text{nat}) \ \text{mltl}$
fixes $\pi::\text{trace}$
assumes *intervals-welldef* φ
assumes π -long-enough: $\text{length } \pi \geq \text{complen-mltl } \varphi$

shows $\text{match } \pi \text{ (pad-WEST-reg } \varphi) \longleftrightarrow \text{semantics-mltl } \pi \varphi$
 $\langle \text{proof} \rangle$

lemma *WEST-correct-pad*:
fixes $\varphi::(\text{nat}) \text{ mltl}$
fixes $\pi::\text{trace}$
assumes *intervals-welldef* φ
assumes $\pi\text{-long-enough}$: $\text{length } \pi \geq \text{complen-mltl } \varphi$
shows $\text{match } \pi \text{ (simp-pad-WEST-reg } \varphi) \longleftrightarrow \text{semantics-mltl } \pi \varphi$
 $\langle \text{proof} \rangle$

end

4 Key algorithms for WEST

theory *Regex-Equivalence*

imports *WEST-Algorithms WEST-Proofs*

begin

fun *depth-datatype-list*:: $\text{state-regex} \Rightarrow \text{nat}$
where *depth-datatype-list* $[] = 0$
 $| \text{depth-datatype-list } (\text{One}\#T) = 1 + \text{depth-datatype-list } T$
 $| \text{depth-datatype-list } (\text{Zero}\#T) = 1 + \text{depth-datatype-list } T$
 $| \text{depth-datatype-list } (S\#T) = 2 + 2 * (\text{depth-datatype-list } T)$

function *enumerate-list*:: $\text{state-regex} \Rightarrow \text{trace-regex}$
where *enumerate-list* $[] = [[]]$
 $| \text{enumerate-list } (\text{One}\#T) = (\text{map } (\lambda x. \text{One}\#x) (\text{enumerate-list } T))$
 $| \text{enumerate-list } (\text{Zero}\#T) = (\text{map } (\lambda x. \text{Zero}\#x) (\text{enumerate-list } T))$
 $| \text{enumerate-list } (S\#T) = (\text{enumerate-list } (\text{Zero}\#T)) @ (\text{enumerate-list } (\text{One}\#T))$
 $\langle \text{proof} \rangle$
termination $\langle \text{proof} \rangle$

fun *flatten-list*:: $'a \text{ list list} \Rightarrow 'a \text{ list}$
where *flatten-list* $L = \text{foldr } (@) L []$

value *flatten-list* $[[12, 13::\text{nat}], [15]]$

value *flatten-list* (let *enumerate-H* = *enumerate-list* $[S, \text{One}]$ in
let *enumerate-T* = $[[]]$ in
 $\text{map } (\lambda t. (\text{map } (\lambda h. h\#t) \text{enumerate-H})) \text{enumerate-T}$)

```

fun enumerate-trace:: trace-regex  $\Rightarrow$  WEST-regex
  where enumerate-trace [] = []
  | enumerate-trace (H#T) = flatten-list
    (let enumerate-H = enumerate-list H in
     let enumerate-T = enumerate-trace T in
     map ( $\lambda t$ . (map ( $\lambda h$ . h#t) enumerate-H)) enumerate-T)

value enumerate-trace [[S, One], [S], [One]]
value enumerate-trace []

fun enumerate-sets:: WEST-regex  $\Rightarrow$  trace-regex set
  where enumerate-sets [] = {}
  | enumerate-sets (h#T) = (set (enumerate-trace h))  $\cup$  (enumerate-sets T)

fun naive-equivalence:: WEST-regex  $\Rightarrow$  WEST-regex  $\Rightarrow$  bool
  where naive-equivalence A B = (A = B  $\vee$  (enumerate-sets A) = (enumerate-sets B))

```

5 Regex Equivalence Correctness

```

lemma enumerate-list-len-alt:
  shows  $\forall$  state  $\in$  set (enumerate-list state-regex).
    length state = length state-regex
  <proof>

```

```

lemma enumerate-list-len:
  assumes state  $\in$  set (enumerate-list state-regex)
  shows length state = length state-regex
  <proof>

```

```

lemma enumerate-list-prop:
  assumes ( $\bigwedge k$ . List.member j k  $\implies$  k  $\neq$  S)
  shows enumerate-list j = [j]
  <proof>

```

```

lemma enumerate-fixed-trace:
  fixes h1:: trace-regex
  assumes  $\bigwedge j$ . List.member h1 j  $\implies$  ( $\bigwedge k$ . List.member j k  $\implies$  k  $\neq$  S)
  shows (enumerate-trace h1) = [h1]
  <proof>

```

If we have two state regexs that don't contain S's, then enumerate trace on each is different.

```

lemma enum-trace-prop:
  fixes h1 h2:: trace-regex

```

assumes $\bigwedge j. \text{List.member } h1 \ j \implies (\bigwedge k. \text{List.member } j \ k \implies k \neq S)$
assumes $\bigwedge j. \text{List.member } h2 \ j \implies (\bigwedge k. \text{List.member } j \ k \implies k \neq S)$
assumes $(\text{set } h1) \neq (\text{set } h2)$
shows $\text{set } (\text{enumerate-trace } h1) \neq \text{set } (\text{enumerate-trace } h2)$
 $\langle \text{proof} \rangle$

lemma *enumerate-list-tail-in*:
assumes $\text{head-}t\#tail-t \in \text{set } (\text{enumerate-list } (h\#trace))$
shows $tail-t \in \text{set } (\text{enumerate-list } trace)$
 $\langle \text{proof} \rangle$

lemma *enumerate-list-fixed*:
assumes $t \in \text{set } (\text{enumerate-list } trace)$
shows $(\forall k. \text{List.member } t \ k \longrightarrow k \neq S)$
 $\langle \text{proof} \rangle$

lemma *map-enum-list-nonempty*:
fixes $t::\text{WEST-bit list list}$
fixes $head::\text{WEST-bit list}$
shows $\text{map } (\lambda h. h \# t) (\text{enumerate-list } head) \neq []$
 $\langle \text{proof} \rangle$

lemma *length-of-flatten-list*:
assumes $\text{flat} =$
 $\text{foldr } (@)$
 $(\text{map } (\lambda t. \text{map } (\lambda h. h \# t) H) T) []$
shows $\text{length flat} = \text{length } T * \text{length } H$
 $\langle \text{proof} \rangle$

lemma *flatten-list-idx*:
assumes $\text{flat} = \text{flatten-list } (\text{map } (\lambda t. \text{map } (\lambda h. h \# t) \text{head}) \text{tail})$
assumes $i < \text{length tail}$
assumes $j < \text{length head}$
shows $(\text{head}!j)\#(\text{tail}!i) = \text{flat}!(i*(\text{length head}) + j) \wedge i*(\text{length head}) + j < \text{length flat}$
 $\langle \text{proof} \rangle$

lemma *flatten-list-shape*:
assumes $\text{List.member flat } x1$
assumes $\text{flat} = \text{flatten-list } (\text{map } (\lambda t. \text{map } (\lambda h. h \# t) H) T)$
shows $\exists x1\text{-head } x1\text{-tail. } x1 = x1\text{-head}\#x1\text{-tail} \wedge \text{List.member } H \ x1\text{-head} \wedge \text{List.member } T \ x1\text{-tail}$
 $\langle \text{proof} \rangle$

lemma *flatten-list-len*:

assumes $\bigwedge t. \text{List.member } T \ t \implies \text{length } t = n$
assumes $\text{flat} = \text{flatten-list } (\text{map } (\lambda t. \text{map } (\lambda h. h \# t) \ H) \ T)$
shows $\bigwedge x1. \text{List.member flat } x1 \implies \text{length } x1 = n+1$
 $\langle \text{proof} \rangle$

lemma *flatten-list-lemma*:

assumes $\bigwedge x1. \text{List.member to-flatten } x1 \implies (\bigwedge x2. \text{List.member } x1 \ x2 \implies \text{length } x2 = \text{length trace})$
assumes $a \in \text{set } (\text{flatten-list to-flatten})$
shows $\text{length } a = \text{length trace}$
 $\langle \text{proof} \rangle$

lemma *enumerate-trace-len*:

assumes $a \in \text{set } (\text{enumerate-trace trace})$
shows $\text{length } a = \text{length trace}$
 $\langle \text{proof} \rangle$

definition *regex-zeros-and-ones*:: $\text{trace-regex} \Rightarrow \text{bool}$

where $\text{regex-zeros-and-ones tr} =$
 $(\forall j. \text{List.member tr } j \longrightarrow (\forall k. \text{List.member } j \ k \longrightarrow k \neq S))$

lemma *match-enumerate-state-aux-first-bit*:

assumes $a\text{-head} = \text{Zero} \vee a\text{-head} = \text{One}$
assumes $a\text{-head} \# a\text{-tail} \in \text{set } (\text{enumerate-list } (h\text{-head} \# h))$
shows $h\text{-head} = a\text{-head} \vee h\text{-head} = S$
 $\langle \text{proof} \rangle$

lemma *advance-state-iff*:

assumes $x > 0$
shows $x \in \text{state} \longleftrightarrow (x-1) \in \text{advance-state state}$
 $\langle \text{proof} \rangle$

lemma *match-enumerate-state-aux*:

assumes $a \in \text{set } (\text{enumerate-list } h)$
assumes $\text{match-timestep state } a$
shows $\text{match-timestep state } h$
 $\langle \text{proof} \rangle$

lemma *enumerate-list-index-one*:

assumes $j < \text{length } (\text{enumerate-list } a)$
shows $\text{One} \# \text{enumerate-list } a \ ! \ j = \text{enumerate-list } (S \# a) \ ! \ (\text{length } (\text{enumerate-list } a) + j) \wedge$
 $(\text{length } (\text{enumerate-list } a) + j < \text{length } (\text{enumerate-list } (S \# a)))$

$\langle \text{proof} \rangle$

lemma *list-concat-index*:

assumes $j < \text{length } L1$

shows $(L1 @ L2)!j = L1!j$

$\langle \text{proof} \rangle$

lemma *enumerate-list-index-zero*:

assumes $j < \text{length } (\text{enumerate-list } a)$

shows $\text{Zero} \# \text{enumerate-list } a ! j = \text{enumerate-list } (S \# a) ! j \wedge$
 $j < \text{length } (\text{enumerate-list } (S \# a))$

$\langle \text{proof} \rangle$

lemma *match-enumerate-list*:

assumes $\text{match-timestep state } a$

shows $\exists j < \text{length } (\text{enumerate-list } a).$

$\text{match-timestep state } (\text{enumerate-list } a ! j)$

$\langle \text{proof} \rangle$

lemma *enumerate-trace-head-in*:

assumes $a\text{-head} \# a\text{-tail} \in \text{set } (\text{enumerate-trace } (h \# \text{trace}))$

shows $a\text{-head} \in \text{set } (\text{enumerate-list } h)$

$\langle \text{proof} \rangle$

lemma *enumerate-trace-tail-in*:

assumes $a\text{-head} \# a\text{-tail} \in \text{set } (\text{enumerate-trace } (h \# \text{trace}))$

shows $a\text{-tail} \in \text{set } (\text{enumerate-trace trace})$

$\langle \text{proof} \rangle$

Intuitively, this says that the traces in enumerate trace h are “more specific” than h, which is “more generic”—i.e., h matches everything that each element of enumerate trace h matches.

lemma *match-enumerate-trace-aux*:

assumes $a \in \text{set } (\text{enumerate-trace trace})$

assumes $\text{match-regex } \pi a$

shows $\text{match-regex } \pi \text{trace}$

$\langle \text{proof} \rangle$

lemma *match-enumerate-trace*:

assumes $a \in \text{set } (\text{enumerate-trace } h) \wedge \text{match-regex } \pi a$

shows $\text{match } \pi (h \# T)$

$\langle \text{proof} \rangle$

lemma *match-enumerate-sets1*:

```

assumes ( $\exists r \in (\text{enumerate-sets } R). \text{match-regex } \pi \ r$ )
shows ( $\text{match } \pi \ R$ )
 $\langle \text{proof} \rangle$ 

lemma match-cases:
  assumes  $\text{match } \pi \ (a \ \# \ R)$ 
  shows  $\text{match } \pi \ [a] \ \vee \ \text{match } \pi \ R$ 
 $\langle \text{proof} \rangle$ 

lemma enumerate-trace-decompose:
  assumes  $\text{state} \in \text{set } (\text{enumerate-list } h)$ 
  assumes  $\text{trace} \in \text{set } (\text{enumerate-trace } T)$ 
  shows  $\text{state} \# \text{trace} \in \text{set } (\text{enumerate-trace } (h \# T))$ 
 $\langle \text{proof} \rangle$ 

lemma match-enumerate-trace-aux-converse:
  assumes  $\text{match-regex } \pi \ \text{trace}$ 
  shows  $\text{match } \pi \ (\text{enumerate-trace } \text{trace})$ 
 $\langle \text{proof} \rangle$ 

lemma match-enumerate-sets2:
  assumes ( $\text{match } \pi \ R$ )
  shows ( $\exists r \in \text{enumerate-sets } R. \text{match-regex } \pi \ r$ )
 $\langle \text{proof} \rangle$ 

lemma match-enumerate-sets:
  shows ( $\exists r \in \text{enumerate-sets } R. \text{match-regex } \pi \ r$ )  $\longleftrightarrow (\text{match } \pi \ R)$ 
 $\langle \text{proof} \rangle$ 

lemma regex-equivalence-correct1:
  assumes ( $\text{naive-equivalence } A \ B$ )
  shows  $\text{match } \pi \ A = \text{match } \pi \ B$ 
 $\langle \text{proof} \rangle$ 

lemma regex-equivalence-correct:
  shows ( $\text{naive-equivalence } A \ B$ )  $\longrightarrow (\text{regex-equiv } A \ B)$ 
 $\langle \text{proof} \rangle$ 

export-code naive-equivalence in Haskell module-name regex-equiv

end

```

References

- [1] J. Elwing, L. Gamboa-Guzman, J. Sorkin, C. Travesset, Z. Wang, and K. Y. Rozier. Mission-time LTL (MLTL) formula validation via regular expressions. In P. Herber and A. Wijs, editors, *iFM*, volume 14300 of *LNCS*, pages 279–301. Springer, 2023.
- [2] Z. Wang, L. P. Gamboa-Guzman, and K. Y. Rozier. WEST: Interactive Validation of Mission-time Linear Temporal Logic (MLTL). 2024.