Formalization of Randomized Approximation Algorithms for Frequency Moments

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Abstract

In 1999 Alon et. al. introduced the still active research topic of approximating the frequency moments of a data stream using randomized algorithms with minimal space usage. This includes the problem of estimating the cardinality of the stream elements—the zeroth frequency moment. But, also higher-order frequency moments that provide information about the skew of the data stream. (The k-th frequency moment of a data stream is the sum of the k-th powers of the occurrence counts of each element in the stream.) This entry formalizes three randomized algorithms for the approximation of F_0, F_2 and F_k for $k \geq 3$ based on [1, 2] and verifies their expected accuracy, success probability and space usage.

Contents

1 Preliminary Results

theory *Frequency-Moments-Preliminary-Results*
imports
HOL.Transcendental
HOL−Computational-Algebra.Primes
HOL−Library.Extended-Real
HOL−Library.Multiset
HOL−Library.Sublist
Prefix-Free-Code-Combinators.Prefix-Free-Code-Combinators
Bertrands-Postulate.Bertrand
Expander-Graphs.Expander-Graphs-Multiset-Extras
begin

This section contains various preliminary results.

lemma *card-ordered-pairs*:
fixes M :: ('a ::linorder) set
assumes finite M
shows $2 * card \{(x,y) \in M \times M.\ x < y\} = card\ M * (card\ M - 1)$
proof −
have a: finite $(M \times M)$ **using** *assms* **by** *simp*

have inj-swap: inj $(\lambda x.\ (snd\ x,\ fst\ x))$
by (rule inj-onI, simp add: prod-eq-iff)

have $2 * card \{(x,y) \in M \times M.\ x < y\} =$
$card \{(x,y) \in M \times M.\ x < y\} + card ((\lambda x.\ (snd\ x,\ fst\ x))`\{(x,y) \in M \times M.\ x < y\})$
by (simp add: card-image[OF inj-on-subset[OF inj-swap]])
also have ... $= card \{(x,y) \in M \times M.\ x < y\} + card \{(x,y) \in M \times M.\ y < x\}$
by (auto intro: arg-cong[**where** f=card] simp add:set-eq-iff image-iff)
also have ... $= card (\{(x,y) \in M \times M.\ x < y\} \cup \{(x,y) \in M \times M.\ y < x\})$
by (intro card-Un-disjoint[symmetric] a finite-subset[**where** $B=M \times M$] subsetI) auto
also have ... $= card ((M \times M) - \{(x,y) \in M \times M.\ x = y\})$
by (auto intro: arg-cong[**where** f=card] simp add:set-eq-iff)
also have ... $= card (M \times M) - card \{(x,y) \in M \times M.\ x = y\}$
by (intro card-Diff-subset a finite-subset[**where** $B=M \times M$] subsetI) auto
also have ... $= card\ M \hat{}\ 2 - card ((\lambda x.\ (x,x))\ `\ M)$
using *assms*
by (intro arg-cong2[**where** f=(−)] arg-cong[**where** f=card])
(auto simp:power2-eq-square set-eq-iff image-iff)
also have ... $= card\ M \hat{}\ 2 - card\ M$
by (intro arg-cong2[**where** f=(−)] card-image inj-onI, auto)

```
  also have ... = card M * (card M - 1)
    by (cases card M ≥ 0, auto simp:power2-eq-square algebra-simps)
  finally show ?thesis by simp
qed

lemma ereal-mono: x ≤ y ⟹ ereal x ≤ ereal y
  by simp

lemma abs-ge-iff: ((x::real) ≤ abs y) = (x ≤ y ∨ x ≤ -y)
  by linarith

lemma count-list-gr-1:
  (x ∈ set xs) = (count-list xs x ≥ 1)
  by (induction xs, simp, simp)

lemma count-list-append: count-list (xs@ys) v = count-list xs v + count-list ys v
  by (induction xs, simp, simp)

lemma count-list-lt-suffix:
  assumes suffix a b
  assumes x ∈ {b ! i| i. i <  length b - length a}
  shows  count-list a x < count-list b x
proof -
  have length a ≤ length b using assms(1)
    by (simp add: suffix-length-le)
  hence x ∈ set (nths b {i. i < length b - length a})
    using assms diff-commute by (auto simp add:set-nths)
  hence a:x ∈ set (take (length b - length a) b)
    by (subst (asm) lessThan-def[symmetric], simp)
  have b = (take (length b - length a) b)@drop (length b - length a) b
    by simp
  also have ... = (take (length b - length a) b)@a
    using assms(1) suffix-take by auto
  finally have b:b = (take (length b - length a) b)@a by simp

  have count-list a x < 1 + count-list a x by simp
  also have ... ≤ count-list (take (length b - length a) b) x + count-list a x
    using a count-list-gr-1
    by (intro add-mono, fast, simp)
  also have ... = count-list b x
    using b count-list-append by metis
  finally show ?thesis by simp
qed

lemma suffix-drop-drop:
  assumes x ≥ y
  shows suffix (drop x a) (drop y a)
proof -
  have drop y a = take (x - y) (drop y a)@drop (x- y) (drop y a)
```

```
    by (subst append-take-drop-id, simp)
  also have  ... = take (x−y) (drop y a)@drop x a
    using assms by simp
  finally have drop y a = take (x−y) (drop y a)@drop x a by simp
  thus ?thesis
    by (auto simp add:suffix-def)
qed

lemma count-list-card: count-list xs x = card {k. k < length xs ∧ xs ! k = x}
proof −
  have count-list xs x = length (filter ((=) x) xs)
    by (induction xs, simp, simp)
  also have ... = card {k. k < length xs ∧ xs ! k = x}
    by (subst length-filter-conv-card, metis)
  finally show ?thesis by simp
qed

lemma card-gr-1-iff:
  assumes finite S  x ∈ S  y ∈ S  x ≠ y
  shows card S > 1
  using assms card-le-Suc0-iff-eq leI by auto

lemma count-list-ge-2-iff:
  assumes y < z
  assumes z < length xs
  assumes xs ! y = xs ! z
  shows count-list xs (xs ! y) > 1
proof −
  have  1 < card {k. k < length xs ∧ xs ! k = xs ! y}
    using assms by (intro card-gr-1-iff[where x=y and y=z], auto)

  thus ?thesis
    by (simp add: count-list-card)
qed
```

Results about multisets and sorting

```
lemmas disj-induct-mset = disj-induct-mset

lemma prod-mset-conv:
  fixes f :: 'a ⇒ 'b::{comm-monoid-mult}
  shows prod-mset (image-mset f A) = prod (λx. f x^(count A x)) (set-mset A)
proof (induction A rule: disj-induct-mset)
  case 1
  then show ?case by simp
next
  case (2 n M x)
  moreover have count M x = 0 using 2 by (simp add: count-eq-zero-iff)
  moreover have ⋀y. y ∈ set-mset M ⟹ y ≠ x using 2 by blast
  ultimately show ?case by (simp add:algebra-simps)
```

qed

There is a version *sum-list-map-eq-sum-count* but it doesn't work if the function maps into the reals.

lemma *sum-list-eval*:
  fixes f :: $'a \Rightarrow 'b$::{*ring*,*semiring-1*}
  shows *sum-list* (*map f xs*) = ($\sum$ x ∈ *set xs*. *of-nat* (*count-list xs x*) ∗ *f x*)
proof −
  define M **where** M = *mset xs*
  have *sum-mset* (*image-mset f M*) = ($\sum$ x ∈ *set-mset M*. *of-nat* (*count M x*) ∗ *f x*)
  proof (*induction M rule*:*disj-induct-mset*)
    case *1*
    then show *?case* **by** *simp*
  next
    case (*2 n M x*)
    have *a*:$\bigwedge y.\ y \in$ *set-mset* $M \Longrightarrow y \neq x$ **using** *2*(*2*) **by** *blast*
    show *?case* **using** *2* **by** (*simp add*:*a count-eq-zero-iff*[*symmetric*])
  qed
  moreover have $\bigwedge x$. *count-list xs x* = *count* (*mset xs*) *x*
    by (*induction xs*, *simp*, *simp*)
  ultimately show *?thesis*
    by (*simp add*:*M-def sum-mset-sum-list*[*symmetric*])
qed

lemma *prod-list-eval*:
  fixes f :: $'a \Rightarrow 'b$::{*ring*,*semiring-1*,*comm-monoid-mult*}
  shows *prod-list* (*map f xs*) = ($\prod$ x ∈ *set xs*. (*f x*)^(*count-list xs x*))
proof −
  define M **where** M = *mset xs*
  have *prod-mset* (*image-mset f M*) = ($\prod$ x ∈ *set-mset M*. *f x* ^ (*count M x*))
  proof (*induction M rule*:*disj-induct-mset*)
    case *1*
    then show *?case* **by** *simp*
  next
    case (*2 n M x*)
    have *a*:$\bigwedge y.\ y \in$ *set-mset* $M \Longrightarrow y \neq x$ **using** *2*(*2*) **by** *blast*
    have *b*:*count M x* = *0* **using** *2* **by** (*subst count-eq-zero-iff*) *blast*
    show *?case* **using** *2* **by** (*simp add*:*a b mult.commute*)
  qed
  moreover have $\bigwedge x$. *count-list xs x* = *count* (*mset xs*) *x*
    by (*induction xs*, *simp*, *simp*)
  ultimately show *?thesis*
    by (*simp add*:*M-def prod-mset-prod-list*[*symmetric*])
qed

lemma *sorted-sorted-list-of-multiset*: *sorted* (*sorted-list-of-multiset M*)
  by (*induction M*, *auto simp*:*sorted-insort*)

lemma *count-mset*: *count* (*mset xs*) *a* = *count-list xs a*
  by (*induction xs*, *auto*)

lemma *swap-filter-image*: *filter-mset g* (*image-mset f A*) = *image-mset f* (*filter-mset* (*g* ∘ *f*) *A*)
  by (*induction A*, *auto*)

lemma *list-eq-iff*:
  assumes *mset xs* = *mset ys*
  assumes *sorted xs*
  assumes *sorted ys*
  shows *xs* = *ys*
  using *assms properties-for-sort* **by** *blast*

lemma *sorted-list-of-multiset-image-commute*:
  assumes *mono f*
  shows *sorted-list-of-multiset* (*image-mset f M*) = *map f* (*sorted-list-of-multiset M*)
proof −
  have *sorted* (*sorted-list-of-multiset* (*image-mset f M*))
    by (*simp add*:*sorted-sorted-list-of-multiset*)
  moreover have *sorted-wrt* (*λx y. f x ≤ f y*) (*sorted-list-of-multiset M*)
    by (*rule sorted-wrt-mono-rel*[**where** *P*=*λx y. x ≤ y*])
      (*auto intro*: *monoD*[*OF assms*] *sorted-sorted-list-of-multiset*)
  hence *sorted* (*map f* (*sorted-list-of-multiset M*))
    by (*subst sorted-wrt-map*)
  ultimately show *?thesis*
    by (*intro list-eq-iff*, *auto*)
qed

Results about rounding and floating point numbers

lemma *round-down-ge*:
  x ≤ *round-down prec x* + *2 powr* (−*prec*)
  using *round-down-correct* **by** (*simp*, *meson diff-diff-eq diff-eq-diff-less-eq*)

lemma *truncate-down-ge*:
  x ≤ *truncate-down prec x* + *abs x* ∗ *2 powr* (−*prec*)
proof (*cases abs x* > *0*)
  case *True*
  have *x* ≤ *round-down* (*int prec* − ⌊*log 2* |*x*|⌋) *x* + *2 powr* (−*real-of-int*(*int prec* − ⌊*log 2* |*x*|⌋))
    by (*rule round-down-ge*)
  also have ... ≤ *truncate-down prec x* + *2 powr* (⌊*log 2* |*x*|⌋) ∗ *2 powr* (−*real prec*)
    by (*rule add-mono*, *simp-all add*:*powr-add*[*symmetric*] *truncate-down-def*)
  also have ... ≤ *truncate-down prec x* + |*x*| ∗ *2 powr* (−*real prec*)
    using *True*
    by (*intro add-mono mult-right-mono*, *simp-all add*:*le-log-iff*[*symmetric*])
  finally show *?thesis* **by** *simp*

next
  case *False*
  then show *?thesis* **by** *simp*
qed

lemma *truncate-down-pos*:
  assumes $x \geq 0$
  shows $x * (1 - 2\ powr\ (-prec)) \leq truncate\text{-}down\ prec\ x$
  by (*simp add:right-diff-distrib diff-le-eq*)
   (*metis truncate-down-ge assms abs-of-nonneg*)

lemma *truncate-down-eq*:
  assumes *truncate-down r x = truncate-down r y*
  shows *abs (x−y) ≤ max (abs x) (abs y) ∗ 2 powr (−real r)*
proof −
  have *x − y ≤ truncate-down r x + abs x ∗ 2 powr (−real r) − y*
    by (*rule diff-right-mono*, *rule truncate-down-ge*)
  also have ... ≤ *y + abs x ∗ 2 powr (−real r) − y*
    using *truncate-down-le*
    by (*intro diff-right-mono add-mono*, *subst assms*(*1*), *simp-all*)
  also have ... ≤ *abs x ∗ 2 powr (−real r)* **by** *simp*
  also have ... ≤ *max (abs x) (abs y) ∗ 2 powr (−real r)* **by** *simp*
  finally have *a:x − y ≤ max (abs x) (abs y) ∗ 2 powr (−real r)* **by** *simp*

  have *y − x ≤ truncate-down r y + abs y ∗ 2 powr (−real r) − x*
    by (*rule diff-right-mono*, *rule truncate-down-ge*)
  also have ... ≤ *x + abs y ∗ 2 powr (−real r) − x*
    using *truncate-down-le*
    by (*intro diff-right-mono add-mono*, *subst assms*(*1*)[*symmetric*], *auto*)
  also have ... ≤ *abs y ∗ 2 powr (−real r)* **by** *simp*
  also have ... ≤ *max (abs x) (abs y) ∗ 2 powr (−real r)* **by** *simp*
  finally have *b:y − x ≤ max (abs x) (abs y) ∗ 2 powr (−real r)* **by** *simp*

  show *?thesis*
    using *abs-le-iff a b* **by** *linarith*
qed

definition *rat-of-float* :: *float* ⇒ *rat* **where**
  rat-of-float f = of-int (mantissa f) ∗
    (*if exponent f ≥ 0 then 2 ^ (nat (exponent f)) else 1 / 2 ^ (nat (−exponent f)))*)

lemma *real-of-rat-of-float*: *real-of-rat (rat-of-float x) = real-of-float x*
proof −
  have *real-of-rat (rat-of-float x) = mantissa x ∗ (2 powr (exponent x))*
   by (*simp add:rat-of-float-def of-rat-mult of-rat-divide of-rat-power powr-realpow*[*symmetric*] *powr-minus-divide*)
  also have ... = *real-of-float x*
    using *mantissa-exponent* **by** *simp*

```
  finally show ?thesis by simp
qed

lemma log-est: log 2 (real n + 1) ≤ n
proof −
  have 1 + real n = real (n + 1)
    by simp
  also have ... ≤ real (2 ^ n)
    by (intro of-nat-mono suc-n-le-2-pow-n)
  also have ... = 2 powr (real n)
    by (simp add:powr-realpow)
  finally have 1 + real n ≤ 2 powr (real n)
    by simp
  thus ?thesis
    by (simp add: Transcendental.log-le-iff)
qed

lemma truncate-mantissa-bound:
  abs (⌊x ∗ 2 powr (real r − real-of-int ⌊log 2 |x|⌋)⌋) ≤ 2 ^(r+1) (is ?lhs ≤ -)
proof −
  define q where q = ⌊x ∗ 2 powr (real r − real-of-int (⌊log 2 |x|⌋))⌋

  have abs q ≤ 2 ^(r + 1) if a:x > 0
  proof −
    have abs q = q
      using a by (intro abs-of-nonneg, simp add:q-def)
    also have ... ≤ x ∗ 2 powr (real r − real-of-int ⌊log 2 |x|⌋)
      unfolding q-def using of-int-floor-le by blast
    also have ... = x ∗ 2 powr real-of-int (int r − ⌊log 2 |x|⌋)
      by auto
    also have ... = 2 powr (log 2 x + real-of-int (int r − ⌊log 2 |x|⌋))
      using a by (simp add:powr-add)
    also have ... ≤ 2 powr (real r + 1)
      using a by (intro powr-mono, linarith+)
    also have ... = 2 ^(r+1)
      by (subst powr-realpow[symmetric], simp-all add:add.commute)
    finally show abs q ≤ 2 ^(r+1)
      by (metis of-int-le-iff of-int-numeral of-int-power)
  qed

  moreover have abs q ≤ (2 ^(r + 1)) if a: x < 0
  proof −
    have −(2 ^(r+1) + 1) = −(2 powr (real r + 1)+1)
      by (subst powr-realpow[symmetric], simp-all add: add.commute)
    also have  ... < −(2 powr (log 2 (− x) + (r − ⌊log 2 |x|⌋)) + 1)
      using a by (simp, linarith)
    also have ... = x ∗ 2 powr (r − ⌊log 2 |x|⌋) − 1
      using a by (simp add:powr-add)
    also have ... ≤ q
```

by (*simp add*:*q-def*)
also have ... = − *abs q*
using *a*
by (*subst abs-of-neg*, *simp-all add*: *mult-pos-neg2 q-def*)
finally have $-(2\ \hat{}\,(r+1)+1) < -$ *abs q* **using** *of-int-less-iff* **by** *fastforce*
hence $-(2\ \hat{}\,(r+1)) \leq -$ *abs q* **by** *linarith*
thus *abs q* $\leq 2\hat{}(r+1)$ **by** *linarith*
qed

moreover have $x = 0 \Longrightarrow$ *abs q* $\leq 2\hat{}(r+1)$
by (*simp add*:*q-def*)
ultimately have *abs q* $\leq 2\hat{}(r+1)$
by *fastforce*
thus *?thesis* **using** *q-def* **by** *blast*
qed

lemma *truncate-float-bit-count*:
bit-count (F_e (*float-of* (*truncate-down r x*))) $\leq 10 + 4 *$ *real r* $+ 2*$*log* 2 (2 + |*log* 2 |*x*||)
(**is** *?lhs* ≤ *?rhs*)
proof −
define *m* **where** $m = \lfloor x * 2$ *powr* (*real r* − *real-of-int* ⌊*log* 2 |*x*|⌋)⌋
define *e* **where** $e = \lfloor$*log* 2 |*x*|⌋ − *int r*

have *a*: (*real-of-int* ⌊*log* 2 |*x*|⌋ − *real r*) = *e*
by (*simp add*:*e-def*)
have *abs m* $+ 2 \leq 2\ \hat{}\,(r + 1) + 2\hat{}1$
using *truncate-mantissa-bound*
by (*intro add-mono*, *simp-all add*:*m-def*)
also have ... $\leq 2\ \hat{}\,(r+2)$
by *simp*
finally have *b*:*abs m* $+ 2 \leq 2\ \hat{}\,(r+2)$ **by** *simp*
hence *real-of-int* (|*m*| + 2) ≤ *real-of-int* ($4 * 2\ \hat{}\ r$)
by (*subst of-int-le-iff*, *simp*)
hence |*real-of-int m*| $+ 2 \leq 4 * 2\ \hat{}\ r$
by *simp*
hence *c*:*log* 2 (*real-of-int* (|*m*| + 2)) ≤ *r*+2
by (*simp add*: *Transcendental.log-le-iff powr-add powr-realpow*)

have *real-of-int* (*abs e* + 1) ≤ *real-of-int* |⌊*log* 2 |*x*|⌋| + *real-of-int r* + 1
by (*simp add*:*e-def*)
also have ... ≤ 1 + *abs* (*log* 2 (*abs x*)) + *real-of-int r* + 1
by (*simp add*:*abs-le-iff*, *linarith*)
also have ... ≤ (*real-of-int r*+ 1) ∗ (2 + *abs* (*log* 2 (*abs x*)))
by (*simp add*:*distrib-left distrib-right*)
finally have *d*:*real-of-int* (*abs e* + 1) ≤ (*real-of-int r*+ 1) ∗ (2 + *abs* (*log* 2 (*abs x*))) **by** *simp*

have *log* 2 (*real-of-int* (*abs e* + 1)) ≤ *log* 2 (*real-of-int r* + 1) + *log* 2 (2 + *abs*

$(log\ 2\ (abs\ x)))$
  using *d* **by** (*simp flip*: *log-mult-pos*)
 also have ... $\leq r + log\ 2\ (2 + abs\ (log\ 2\ (abs\ x)))$
  using *log-est* **by** (*intro add-mono*, *simp-all add*:*add.commute*)
 finally have *e*: $log\ 2\ (real\text{-}of\text{-}int\ (abs\ e + 1)) \leq r + log\ 2\ (2 + abs\ (log\ 2\ (abs\ x)))$ **by** *simp*

 have *?lhs* = *bit-count* (F_e (*float-of* (*real-of-int m* $*$ *2 powr real-of-int e*)))
  by (*simp add*:*truncate-down-def round-down-def m-def*[*symmetric*] *a*)
 also have ... $\leq$ *ereal* $(6 + (2 * log\ 2\ (real\text{-}of\text{-}int\ (|m| + 2)) + 2 * log\ 2\ (real\text{-}of\text{-}int\ (|e| + 1))))$
  using *float-bit-count-2* **by** *simp*
 also have ... $\leq$ *ereal* $(6 + (2 * real\ (r+2) + 2 * (r + log\ 2\ (2 + abs\ (log\ 2\ (abs\ x))))))$
  using *c e*
  by (*subst ereal-less-eq*, *intro add-mono mult-left-mono*, *linarith+*)
 also have ... = *?rhs* **by** *simp*
 finally show *?thesis* **by** *simp*
qed

definition *prime-above* :: *nat* $\Rightarrow$ *nat*
 where *prime-above n* = (*SOME x*. $x \in \{n..(2{*}n{+}2)\} \wedge$ *prime x*)

The term *prime-above n* returns a prime between n and $2 * n + 2$. Because of Bertrand's postulate there always is such a value. In a refinement of the algorithms, it may make sense to replace this with an algorithm, that finds such a prime exactly or approximately.

The definition is intentionally inexact, to allow refinement with various algorithms, without modifying the high-level mathematical correctness proof.

lemma *ex-subset*:
 assumes $\exists x \in A.\ P\ x$
 assumes $A \subseteq B$
 shows $\exists x \in B.\ P\ x$
 using *assms* **by** *auto*

lemma
 shows *prime-above-prime*: *prime* (*prime-above n*)
 and *prime-above-range*: *prime-above n* $\in \{n..(2{*}n{+}2)\}$
proof −
 define *r* **where** $r = (\lambda x.\ x \in \{n..(2{*}n{+}2)\} \wedge prime\ x)$
 have $\exists x.\ r\ x$
 proof (*cases n>2*)
  case *True*
  hence $n{-}1 > 1$ **by** *simp*
  hence $\exists x \in \{(n{-}1){<}..{<}(2{*}(n{-}1))\}.\ prime\ x$
   using *bertrand* **by** *simp*
  moreover have $\{n - 1{<}..{<}2 * (n - 1)\} \subseteq \{n..2 * n + 2\}$
   by (*intro subsetI*, *auto*)

```
    ultimately have ∃x ∈ {n..(2*n+2)}. prime x
      by (rule ex-subset)
    then show ?thesis by (simp add:r-def Bex-def)
  next
    case False
    hence 2 ∈ {n..(2*n+2)}
      by simp
    moreover have prime (2::nat)
      using two-is-prime-nat by blast
    ultimately have r 2
      using r-def by simp
    then show ?thesis by (rule exI)
  qed
  moreover have prime-above n = (SOME x. r x)
    by (simp add:prime-above-def r-def)
  ultimately have a:r (prime-above n)
    using someI-ex by metis
  show prime (prime-above n)
    using a unfolding r-def by blast
  show prime-above n ∈ {n..(2*n+2)}
    using a unfolding r-def by blast
qed

lemma prime-above-min:  prime-above n ≥ 2
  using prime-above-prime
  by (simp add: prime-ge-2-nat)

lemma prime-above-lower-bound: prime-above n ≥ n
  using prime-above-range
  by simp

lemma prime-above-upper-bound: prime-above n ≤ 2*n+2
  using prime-above-range
  by simp

end
```

2 Frequency Moments

```
theory Frequency-Moments
  imports
    Frequency-Moments-Preliminary-Results
    Finite-Fields.Finite-Fields-Mod-Ring-Code
    Interpolation-Polynomials-HOL-Algebra.Interpolation-Polynomial-Cardinalities
begin
```

This section contains a definition of the frequency moments of a stream and a few general results about frequency moments..

```
definition F where
```

$F\ k\ xs = (\sum\ x \in set\ xs.\ (rat\text{-}of\text{-}nat\ (count\text{-}list\ xs\ x)\hat{}k))$

lemma *F-ge-0*: $F\ k\ as \geq 0$
unfolding *F-def* **by** (*rule sum-nonneg*, *simp*)

lemma *F-gr-0*:
assumes $as \neq []$
shows $F\ k\ as > 0$
proof −
have $rat\text{-}of\text{-}nat\ 1 \leq rat\text{-}of\text{-}nat\ (card\ (set\ as))$
using *assms card-0-eq*[**where** A=*set as*]
by (*intro of-nat-mono*)
(*metis List.finite-set One-nat-def Suc-leI neq0-conv set-empty*)
also have ... $= (\sum x{\in}set\ as.\ 1)$ **by** *simp*
also have ... $\leq (\sum x{\in}set\ as.\ rat\text{-}of\text{-}nat\ (count\text{-}list\ as\ x)\ \hat{}\ k)$
by (*intro sum-mono one-le-power*)
(*metis count-list-gr-1 of-nat-1 of-nat-le-iff*)
also have ... $\leq F\ k\ as$
by (*simp add*:*F-def*)
finally show *?thesis* **by** *simp*
qed

definition P_e :: *nat* ⇒ *nat* ⇒ *nat list* ⇒ *bool list option* **where**
$P_e\ p\ n\ f = (if\ p > 1 \wedge f \in bounded\text{-}degree\text{-}polynomials\ (ring\text{-}of\ (mod\text{-}ring\ p))\ n$ *then*
$([0..{<}n] \rightarrow_e Nb_e\ p)\ (\lambda i \in \{..{<}n\}.\ ring.coeff\ (ring\text{-}of\ (mod\text{-}ring\ p))\ f\ i)$ *else None*)

lemma *poly-encoding*:
is-encoding $(P_e\ p\ n)$
proof (*cases* $p > 1$)
case *True*
interpret *cring ring-of* (*mod-ring p*)
using *mod-ring-is-cring True* **by** *blast*
have *a*:*inj-on* $(\lambda x.\ (\lambda i{\in}\{..{<}n\}.\ coeff\ x\ i))$ (*bounded-degree-polynomials* (*ring-of* (*mod-ring p*)) *n*)
proof (*rule inj-onI*)
fix *x y*
assume *b*:$x \in bounded\text{-}degree\text{-}polynomials\ (ring\text{-}of\ (mod\text{-}ring\ p))\ n$
assume *c*:$y \in bounded\text{-}degree\text{-}polynomials\ (ring\text{-}of\ (mod\text{-}ring\ p))\ n$
assume *d*:$restrict\ (coeff\ x)\ \{..{<}n\} = restrict\ (coeff\ y)\ \{..{<}n\}$
have $coeff\ x\ i = coeff\ y\ i$ **for** *i*
proof (*cases* $i < n$)
case *True*
then show *?thesis* **by** (*metis lessThan-iff restrict-apply d*)
next
case *False*
hence *e*: $i \geq n$ **by** *linarith*
have $coeff\ x\ i = \mathbf{0}_{ring\text{-}of\ (mod\text{-}ring\ p)}$

```
      using b e by (subst coeff-length, auto simp:bounded-degree-polynomials-length)
      also have ... = coeff y i
      using c e by (subst coeff-length, auto simp:bounded-degree-polynomials-length)
      finally show ?thesis by simp
    qed
    then show x = y
      using b c univ-poly-carrier
     by (subst coeff-iff-polynomial-cond) (auto simp:bounded-degree-polynomials-length)
  qed

  have is-encoding (λf. P_e p n f)
    unfolding P_e-def using a True
   by (intro encoding-compose[where f=([0..<n] →_e Nb_e p)] fun-encoding bounded-nat-encoding)
     auto
  thus ?thesis by simp
next
  case False
  hence is-encoding (λf. P_e p n f)
    unfolding P_e-def using encoding-triv by simp
  then show ?thesis by simp
qed

lemma bounded-degree-polynomial-bit-count:
  assumes p > 1
  assumes x ∈ bounded-degree-polynomials (ring-of (mod-ring p)) n
  shows bit-count (P_e p n x) ≤ ereal (real n * (log 2 p + 1))
proof −
  interpret cring ring-of (mod-ring p)
    using mod-ring-is-cring assms by blast

  have a: x ∈ carrier (poly-ring (ring-of (mod-ring p)))
    using assms(2) by (simp add:bounded-degree-polynomials-def)

  have real-of-int ⌊log 2 (p−1)⌋+1 ≤ log 2 (p−1) + 1
    using floor-eq-iff by (intro add-mono, auto)
  also have ... ≤ log 2 p + 1
    using assms by (intro add-mono, auto)
  finally have b: ⌊log 2 (p−1)⌋+1 ≤ log 2 p + 1
    by simp

  have bit-count (P_e p n x) = (∑ k ← [0..<n]. bit-count (Nb_e p (coeff x k)))
    using assms restrict-extensional
   by (auto intro!:arg-cong[where f=sum-list] simp add:P_e-def fun-bit-count lessThan-atLeast0)
  also have ... = (∑ k ← [0..<n]. ereal (floorlog 2 (p−1)))
    using coeff-in-carrier[OF a] mod-ring-carr
    by (subst bounded-nat-bit-count-2, auto)
  also have ... = n * ereal (floorlog 2 (p−1))
    by (simp add: sum-list-triv)
  also have ... = n * real-of-int (⌊log 2 (p−1)⌋+1)
```

using *assms*(*1*) **by** (*simp add*:*floorlog-def*)
also have ... $\leq$ *ereal* (*real n* $*$ (*log 2 p* + *1*))
by (*subst ereal-less-eq*, *intro mult-left-mono b*, *auto*)
finally show *?thesis* **by** *simp*
qed

end

3 Ranks, k smallest element and elements

theory *K-Smallest*
imports
Frequency-Moments-Preliminary-Results
Interpolation-Polynomials-HOL-Algebra.Interpolation-Polynomial-Cardinalities
begin

This section contains definitions and results for the selection of the k smallest elements, the k-th smallest element, rank of an element in an ordered set.

definition *rank-of* :: *'a* :: *linorder* $\Rightarrow$ *'a set* $\Rightarrow$ *nat* **where** *rank-of x S* = *card* $\{y \in S.\ y < x\}$

The function *rank-of* returns the rank of an element within a set.

lemma *rank-mono*:
assumes *finite S*
shows $x \leq y \Longrightarrow$ *rank-of x S* $\leq$ *rank-of y S*
unfolding *rank-of-def* **using** *assms* **by** (*intro card-mono*, *auto*)

lemma *rank-mono-2*:
assumes *finite S*
shows $S' \subseteq S \Longrightarrow$ *rank-of x S'* $\leq$ *rank-of x S*
unfolding *rank-of-def* **using** *assms* **by** (*intro card-mono*, *auto*)

lemma *rank-mono-commute*:
assumes *finite S*
assumes $S \subseteq T$
assumes *strict-mono-on T f*
assumes $x \in T$
shows *rank-of x S* = *rank-of* (*f x*) (*f* ‘ *S*)
proof −
have *a*: *inj-on f T*
by (*metis assms*(*3*) *strict-mono-on-imp-inj-on*)

have *rank-of* (*f x*) (*f* ‘ *S*) = *card* (*f* ‘ $\{y \in S.\ f\ y < f\ x\}$)
unfolding *rank-of-def* **by** (*intro arg-cong*[**where** *f*=*card*], *auto*)
also have ... = *card* (*f* ‘ $\{y \in S.\ y < x\}$)
using *assms* **by** (*intro arg-cong*[**where** *f*=*card*] *arg-cong*[**where** *f*=(‘) *f*])
(*meson in-mono linorder-not-le strict-mono-onD strict-mono-on-leD set-eq-iff*)
also have ... = *card* $\{y \in S.\ y < x\}$

using *assms* **by** (*intro card-image inj-on-subset*[*OF a*], *blast*)
also have ... = *rank-of x S*
by (*simp add*:*rank-of-def*)
finally show *?thesis*
by *simp*
qed

definition *least* **where** *least k S* = $\{y \in S.$ *rank-of y S* $< k\}$

The function *K-Smallest.least* returns the k smallest elements of a finite set.

lemma *rank-strict-mono*:
assumes *finite S*
shows *strict-mono-on S* ($\lambda x.$ *rank-of x S*)
proof −
have $\bigwedge x\ y.\ x \in S \Longrightarrow y \in S \Longrightarrow x < y \Longrightarrow$ *rank-of x S* < *rank-of y S*
unfolding *rank-of-def* **using** *assms*
by (*intro psubset-card-mono*, *auto*)

thus *?thesis*
by (*simp add*:*rank-of-def strict-mono-on-def*)
qed

lemma *rank-of-image*:
assumes *finite S*
shows ($\lambda x.$ *rank-of x S*) ‘ *S* = {*0..<card S*}
proof (*rule card-seteq*)
show *finite* {*0..<card S*} **by** *simp*

have $\bigwedge x.\ x \in S \Longrightarrow$ *card* $\{y \in S.\ y < x\}$ < *card S*
by (*rule psubset-card-mono*, *metis assms*, *blast*)
thus ($\lambda x.$ *rank-of x S*) ‘ *S* ⊆ {*0..<card S*}
by (*intro image-subsetI*, *simp add*:*rank-of-def*)

have *inj-on* ($\lambda x.$ *rank-of x S*) *S*
by (*metis strict-mono-on-imp-inj-on rank-strict-mono assms*)
thus *card* {*0..<card S*} ≤ *card* (($\lambda x.$ *rank-of x S*) ‘ *S*)
by (*simp add*:*card-image*)
qed

lemma *card-least*:
assumes *finite S*
shows *card* (*least k S*) = *min k* (*card S*)
proof (*cases card S* < *k*)
case *True*
have $\bigwedge t.$ *rank-of t S* ≤ *card S*
unfolding *rank-of-def* **using** *assms*
by (*intro card-mono*, *auto*)
hence $\bigwedge t.$ *rank-of t S* < *k*
by (*metis True not-less-iff-gr-or-eq order-less-le-trans*)

```
  hence least k S = S
    by (simp add:least-def)
  then show ?thesis using True by simp
next
  case False
  hence a:card S ≥ k using leI by blast
  hence card ((λx. rank-of x S) -` {0..<k} ∩ S) = card {0..<k}
    using assms
    by (intro card-vimage-inj-on strict-mono-on-imp-inj-on rank-strict-mono)
     (simp-all add: rank-of-image)
  hence card (least k S) = k
    by (simp add: Collect-conj-eq Int-commute least-def vimage-def)
  then show ?thesis using a by linarith
qed

lemma least-subset: least k S ⊆ S
  by (simp add:least-def)

lemma least-mono-commute:
  assumes finite S
  assumes strict-mono-on S f
  shows f ` least k S = least k (f ` S)
proof -
  have a:inj-on f S
    using strict-mono-on-imp-inj-on[OF assms(2)] by simp

  have card (least k (f ` S)) = min k (card (f ` S))
    by (subst card-least, auto simp add:assms)
  also have ... = min k (card S)
    by (subst card-image, metis a, auto)
  also have ... = card (least k S)
    by (subst card-least, auto simp add:assms)
  also have ... = card (f ` least k S)
    by (subst card-image[OF inj-on-subset[OF a]], simp-all add:least-def)
  finally have b: card (least k (f ` S)) ≤ card (f ` least k S) by simp

  have c: f ` least k S ⊆least k (f ` S)
    using assms by (intro image-subsetI)
      (simp add:least-def rank-mono-commute[symmetric, where T=S])

  show ?thesis
    using b c assms by (intro card-seteq, simp-all add:least-def)
qed

lemma least-eq-iff:
  assumes finite B
  assumes A ⊆ B
  assumes ⋀x. x ∈ B ⟹ rank-of x B < k ⟹ x ∈ A
  shows least k A = least k B
```

proof −
  have *least k B* ⊆ *least k A*
    using *assms rank-mono-2*[*OF assms*(*1*,*2*)] *order-le-less-trans*
    by (*simp add*:*least-def*, *blast*)
  moreover have *card* (*least k B*) ≥ *card* (*least k A*)
    using *assms finite-subset*[*OF assms*(*2*,*1*)] *card-mono*[*OF assms*(*1*,*2*)]
    by (*simp add*: *card-least min-le-iff-disj*)
  moreover have *finite* (*least k A*)
    using *finite-subset least-subset assms*(*1*,*2*) **by** *metis*
  ultimately show *?thesis*
    by (*intro card-seteq*[*symmetric*], *simp-all*)
qed

lemma *least-insert*:
  assumes *finite S*
  shows *least k* (*insert x* (*least k S*)) = *least k* (*insert x S*) (**is** *?lhs* = *?rhs*)
proof (*rule least-eq-iff*)
  show *finite* (*insert x S*)
    using *assms*(*1*) **by** *simp*
  show *insert x* (*least k S*) ⊆ *insert x S*
    using *least-subset* **by** *blast*
  show *y* ∈ *insert x* (*least k S*) **if** *a*: *y* ∈ *insert x S* **and** *b*: *rank-of y* (*insert x S*) < *k* **for** *y*
  proof −
    have *rank-of y S* ≤ *rank-of y* (*insert x S*)
      using *assms* **by** (*intro rank-mono-2*, *auto*)
    also have ... < *k* **using** *b* **by** *simp*
    finally have *rank-of y S* < *k* **by** *simp*
    hence *y* = *x* ∨ (*y* ∈ *S* ∧ *rank-of y S* < *k*)
      using *a* **by** *simp*
    thus *?thesis* **by** (*simp add*:*least-def*)
  qed
qed

definition *count-le* **where** *count-le x M* = *size* {#*y* ∈# *M*. *y* ≤ *x*#}
definition *count-less* **where** *count-less x M* = *size* {#*y* ∈# *M*. *y* < *x*#}

definition *nth-mset* :: *nat* ⇒ (′*a* :: *linorder*) *multiset* ⇒ ′*a* **where**
  nth-mset k M = *sorted-list-of-multiset M* ! *k*

lemma *nth-mset-bound-left*:
  assumes *k* < *size M*
  assumes *count-less x M* ≤ *k*
  shows *x* ≤ *nth-mset k M*
proof (*rule ccontr*)
  define *xs* **where** *xs* = *sorted-list-of-multiset M*
  have *s-xs*: *sorted xs* **by** (*simp add*:*xs-def sorted-sorted-list-of-multiset*)
  have *l-xs*: *k* < *length xs*

```
    using  assms(1) by (simp add:xs-def size-mset[symmetric])
  have M-xs: M = mset xs by (simp add:xs-def)
  hence a:⋀i. i ≤ k ⟹ xs ! i ≤ xs ! k
    using s-xs l-xs sorted-iff-nth-mono by blast

  assume ¬(x ≤ nth-mset k M)
  hence x > nth-mset k M by simp
  hence b:x > xs ! k by (simp add:nth-mset-def xs-def[symmetric])

  have k < card {0..k} by simp
  also have ... ≤ card {i. i < length xs ∧ xs ! i < x}
    using a b l-xs order-le-less-trans
    by (intro card-mono subsetI) auto
  also have ... = length (filter (λy. y < x) xs)
    by (subst length-filter-conv-card, simp)
  also have ... = size (mset (filter (λy. y < x) xs))
    by (subst size-mset, simp)
  also have ... = count-less x M
    by (simp add:count-less-def M-xs)
  also have ... ≤ k
    using assms by simp
  finally show False by simp
qed

lemma nth-mset-bound-left-excl:
  assumes k < size M
  assumes count-le x M ≤ k
  shows x < nth-mset k M
proof (rule ccontr)
  define xs where xs = sorted-list-of-multiset M
  have s-xs: sorted xs by (simp add:xs-def sorted-sorted-list-of-multiset)
  have l-xs: k < length xs
    using  assms(1) by (simp add:xs-def size-mset[symmetric])
  have M-xs: M = mset xs by (simp add:xs-def)
  hence a:⋀i. i ≤ k ⟹ xs ! i ≤ xs ! k
    using s-xs l-xs sorted-iff-nth-mono by blast

  assume ¬(x < nth-mset k M)
  hence x ≥ nth-mset k M by simp
  hence b:x ≥ xs ! k by (simp add:nth-mset-def xs-def[symmetric])

  have k+1 ≤ card {0..k} by simp
  also have ... ≤ card {i. i < length xs ∧ xs ! i ≤ xs ! k}
    using a b l-xs order-le-less-trans
    by (intro card-mono subsetI, auto)
  also have ... ≤ card {i. i < length xs ∧ xs ! i ≤ x}
    using b by (intro card-mono subsetI, auto)
  also have ... = length (filter (λy. y ≤ x) xs)
    by (subst length-filter-conv-card, simp)
```

```
  also have ... = size (mset (filter (λy. y ≤ x) xs))
    by (subst size-mset, simp)
  also have ... = count-le x M
    by (simp add:count-le-def M-xs)
  also have ... ≤ k
    using assms by simp
  finally show False by simp
qed

lemma nth-mset-bound-right:
  assumes k < size M
  assumes count-le x M > k
  shows nth-mset k M ≤ x
proof (rule ccontr)
  define xs where xs = sorted-list-of-multiset M
  have s-xs: sorted xs by (simp add:xs-def sorted-sorted-list-of-multiset)
  have l-xs: k < length xs
    using  assms(1) by (simp add:xs-def size-mset[symmetric])
  have M-xs: M = mset xs by (simp add:xs-def)

  assume ¬(nth-mset k M ≤ x)
  hence x < nth-mset k M by simp
  hence x < xs ! k
    by (simp add:nth-mset-def xs-def[symmetric])
  hence a:⋀i. i < length xs ∧ xs ! i ≤ x ⟹ i < k
    using s-xs l-xs sorted-iff-nth-mono leI by fastforce
  have count-le x M = size (mset (filter (λy. y ≤ x) xs))
    by (simp add:count-le-def M-xs)
  also have ... = length (filter (λy. y ≤ x) xs)
    by (subst size-mset, simp)
  also have ... = card {i. i < length xs ∧ xs ! i ≤ x}
    by (subst length-filter-conv-card, simp)
  also have ... ≤ card {i. i < k}
    using a by (intro card-mono subsetI, auto)
  also have ... = k by simp
  finally have count-le x M ≤ k by simp
  thus False using assms by simp
qed

lemma nth-mset-commute-mono:
  assumes mono f
  assumes k < size M
  shows f (nth-mset k M) = nth-mset k (image-mset f M)
proof −
  have a:k < length (sorted-list-of-multiset M)
    by (metis assms(2) mset-sorted-list-of-multiset size-mset)
  show ?thesis
     using a by (simp add:nth-mset-def sorted-list-of-multiset-image-commute[OF
assms(1)])
```

```
qed

lemma nth-mset-max:
  assumes size A > k
  assumes ⋀x. x ≤ nth-mset k A ⟹ count A x ≤ 1
  shows nth-mset k A = Max (least (k+1) (set-mset A)) and card (least (k+1)
(set-mset A)) = k+1
proof −
  define xs where xs = sorted-list-of-multiset A
  have k-bound: k < length xs unfolding xs-def
    by (metis size-mset mset-sorted-list-of-multiset assms(1))

  have A-def: A = mset xs by (simp add:xs-def)
  have s-xs: sorted xs by (simp add:xs-def sorted-sorted-list-of-multiset)
  have ⋀x. x ≤ xs ! k ⟹ count A x ≤ Suc 0
    using assms(2) by (simp add:xs-def[symmetric] nth-mset-def)
  hence no-col: ⋀x. x ≤ xs ! k ⟹ count-list xs x ≤ 1
    by (simp add:A-def count-mset)

  have inj-xs: inj-on (λk. xs ! k) {0..k}
    by (rule inj-onI, simp) (metis (full-types) count-list-ge-2-iff k-bound no-col
       le-neq-implies-less linorder-not-le order-le-less-trans s-xs sorted-iff-nth-mono)

  have ⋀y. y < length xs ⟹ rank-of (xs ! y) (set xs) < k+1 ⟹ y < k+1
  proof (rule ccontr)
    fix y
    assume b:y < length xs
    assume ¬y < k +1
    hence a:k + 1 ≤ y by simp

    have d:Suc k < length xs using a b by simp

    have k+1 = card ((!) xs ‘ {0..k})
      by (subst card-image[OF inj-xs], simp)
    also have ... ≤ rank-of (xs ! (k+1)) (set xs)
      unfolding rank-of-def using k-bound
       by (intro card-mono image-subsetI conjI, simp-all) (metis count-list-ge-2-iff
no-col not-le le-imp-less-Suc s-xs
          sorted-iff-nth-mono d order-less-le)
    also have ... ≤ rank-of (xs ! y) (set xs)
      unfolding rank-of-def
      by (intro card-mono subsetI, simp-all)
       (metis Suc-eq-plus1 a b s-xs order-less-le-trans sorted-iff-nth-mono)
    also assume ... < k+1
    finally show False by force
  qed

  moreover have rank-of (xs ! y) (set xs) < k+1 if a:y < k + 1 for y
  proof −
```

have *rank-of* (*xs* ! *y*) (*set xs*) $\leq$ *card* ((λ*k*. *xs* ! *k*) ‘ {*k*. *k* < *length xs* $\wedge$ *xs* ! *k* < *xs* ! *y*})
unfolding *rank-of-def*
by (*intro card-mono subsetI*, *simp*)
(*metis* (*no-types*, *lifting*) *imageI in-set-conv-nth mem-Collect-eq*)
also have ... $\leq$ *card* {*k*. *k* < *length xs* $\wedge$ *xs* ! *k* < *xs* ! *y*}
by (*rule card-image-le*, *simp*)
also have ... $\leq$ *card* {*k*. *k* < *y*}
by (*intro card-mono subsetI*, *simp-all add*:*not-less*)
(*metis sorted-iff-nth-mono s-xs linorder-not-less*)
also have ... = *y* **by** *simp*
also have ... < *k* + *1* **using** *a* **by** *simp*
finally show *rank-of* (*xs* ! *y*) (*set xs*) < *k*+*1* **by** *simp*
qed

ultimately have *rank-conv*: $\bigwedge$*y*. *y* < *length xs* $\Longrightarrow$ *rank-of* (*xs* ! *y*) (*set xs*) < *k*+*1* $\longleftrightarrow$ *y* < *k*+*1*
by *blast*

have *y* $\leq$ *xs* ! *k* **if** *a*:*y* $\in$ *least* (*k*+*1*) (*set xs*) **for** *y*
proof −
have *y* $\in$ *set xs* **using** *a least-subset* **by** *blast*
then obtain *i* **where** *i-bound*: *i* < *length xs* **and** *y-def*: *y* = *xs* ! *i* **using** *in-set-conv-nth* **by** *metis*
hence *rank-of* (*xs* ! *i*) (*set xs*) < *k*+*1*
using *a y-def i-bound* **by** (*simp add*: *least-def*)
hence *i* < *k*+*1*
using *rank-conv i-bound* **by** *blast*
hence *i* $\leq$ *k* **by** *linarith*
hence *xs* ! *i* $\leq$ *xs* ! *k*
using *s-xs i-bound k-bound sorted-nth-mono* **by** *blast*
thus *y* $\leq$ *xs* ! *k* **using** *y-def* **by** *simp*
qed

moreover have *xs* ! *k* $\in$ *least* (*k*+*1*) (*set xs*)
using *k-bound rank-conv* **by** (*simp add*:*least-def*)

ultimately have *Max* (*least* (*k*+*1*) (*set xs*)) = *xs* ! *k*
by (*intro Max-eqI finite-subset*[*OF least-subset*], *auto*)

hence *nth-mset k A* = *Max* (*K-Smallest.least* (*Suc k*) (*set xs*))
by (*simp add*:*nth-mset-def xs-def*[*symmetric*])
also have ... = *Max* (*least* (*k*+*1*) (*set-mset A*))
by (*simp add*:*A-def*)
finally show *nth-mset k A* = *Max* (*least* (*k*+*1*) (*set-mset A*)) **by** *simp*

have *k* + *1* = *card* ((λ*i*. *xs* ! *i*) ‘ {*0*..*k*})
by (*subst card-image*[*OF inj-xs*], *simp*)
also have ... $\leq$ *card* (*least* (*k*+*1*) (*set xs*))

using *rank-conv k-bound*
by (*intro card-mono image-subsetI finite-subset*[*OF least-subset*], *simp-all add*:*least-def*)
finally have *card* (*least* (*k+1*) (*set xs*)) ≥ *k+1* **by** *simp*
moreover have *card* (*least* (*k+1*) (*set xs*)) ≤ *k+1*
by (*subst card-least*, *simp*, *simp*)
ultimately have *card* (*least* (*k+1*) (*set xs*)) = *k+1* **by** *simp*
thus *card* (*least* (*k+1*) (*set-mset A*)) = *k+1* **by** (*simp add*:*A-def*)
qed

end

4 Landau Symbols

theory *Landau-Ext*
imports
HOL−Library.Landau-Symbols
HOL.Topological-Spaces
begin

This section contains results about Landau Symbols in addition to "HOL-Library.Landau".

lemma *landau-sum*:
assumes *eventually* (λ*x*. *g1 x* ≥ (*0*::*real*)) *F*
assumes *eventually* (λ*x*. *g2 x* ≥ *0*) *F*
assumes *f1* ∈ *O*[*F*](*g1*)
assumes *f2* ∈ *O*[*F*](*g2*)
shows (λ*x*. *f1 x* + *f2 x*) ∈ *O*[*F*](λ*x*. *g1 x* + *g2 x*)
proof −
obtain *c1* **where** *a1*: *c1* > *0* **and** *b1*: *eventually* (λ*x*. *abs* (*f1 x*) ≤ *c1* ∗ *abs* (*g1 x*)) *F*
using *assms*(*3*) **by** (*simp add*:*bigo-def*, *blast*)
obtain *c2* **where** *a2*: *c2* > *0* **and** *b2*: *eventually* (λ*x*. *abs* (*f2 x*) ≤ *c2* ∗ *abs* (*g2 x*)) *F*
using *assms*(*4*) **by** (*simp add*:*bigo-def*, *blast*)
have *eventually* (λ*x*. *abs* (*f1 x* + *f2 x*) ≤ (*max c1 c2*) ∗ *abs* (*g1 x* + *g2 x*)) *F*
proof (*rule eventually-mono*[*OF eventually-conj*[*OF b1 eventually-conj*[*OF b2 eventually-conj*[*OF assms*(*1*,*2*)]]]])
fix *x*
assume *a*: |*f1 x*| ≤ *c1* ∗ |*g1 x*| ∧ |*f2 x*| ≤ *c2* ∗ |*g2 x*| ∧ *0* ≤ *g1 x* ∧ *0* ≤ *g2 x*
have |*f1 x* + *f2 x*| ≤ |*f1 x* | + |*f2 x*| **using** *abs-triangle-ineq* **by** *blast*
also have ... ≤ *c1* ∗ |*g1 x*| + *c2* ∗ |*g2 x*| **using** *a add-mono* **by** *blast*
also have ... ≤ *max c1 c2* ∗ |*g1 x*| + *max c1 c2* ∗ |*g2 x*|
by (*intro add-mono mult-right-mono*) *auto*
also have ... = *max c1 c2* ∗ (|*g1 x*| + |*g2 x*|)
by (*simp add*:*algebra-simps*)
also have ... ≤ *max c1 c2* ∗ (|*g1 x* + *g2 x*|)
using *a a1 a2* **by** (*intro mult-left-mono*) *auto*
finally show |*f1 x* + *f2 x*| ≤ *max c1 c2* ∗ |*g1 x* + *g2 x*|
by (*simp add*:*algebra-simps*)

qed
hence $0 < max\ c1\ c2 \wedge (\forall_F\ x\ in\ F.\ |f1\ x + f2\ x| \leq max\ c1\ c2 * |g1\ x + g2\ x|)$
using *a1 a2* **by** *linarith*
thus *?thesis*
by (*simp add*: *bigo-def*, *blast*)
qed

lemma *landau-sum-1*:
assumes *eventually* $(\lambda x.\ g1\ x \geq (0::real))\ F$
assumes *eventually* $(\lambda x.\ g2\ x \geq 0)\ F$
assumes $f \in O[F](g1)$
shows $f \in O[F](\lambda x.\ g1\ x + g2\ x)$
proof −
have $f = (\lambda x.\ f\ x + 0)$ **by** *simp*
also have ... $\in O[F](\lambda x.\ g1\ x + g2\ x)$
using *assms zero-in-bigo* **by** (*intro landau-sum*)
finally show *?thesis* **by** *simp*
qed

lemma *landau-sum-2*:
assumes *eventually* $(\lambda x.\ g1\ x \geq (0::real))\ F$
assumes *eventually* $(\lambda x.\ g2\ x \geq 0)\ F$
assumes $f \in O[F](g2)$
shows $f \in O[F](\lambda x.\ g1\ x + g2\ x)$
proof −
have $f = (\lambda x.\ 0 + f\ x)$ **by** *simp*
also have ... $\in O[F](\lambda x.\ g1\ x + g2\ x)$
using *assms zero-in-bigo* **by** (*intro landau-sum*)
finally show *?thesis* **by** *simp*
qed

lemma *landau-ln-3*:
assumes *eventually* $(\lambda x.\ (1::real) \leq f\ x)\ F$
assumes $f \in O[F](g)$
shows $(\lambda x.\ ln\ (f\ x)) \in O[F](g)$
proof −
have $1 \leq x \Longrightarrow |ln\ x| \leq |x|$ **for** $x :: real$
using *ln-bound* **by** *auto*
hence $(\lambda x.\ ln\ (f\ x)) \in O[F](f)$
by (*intro landau-o.big-mono eventually-mono*[*OF assms*(*1*)]) *simp*
thus *?thesis*
using *assms*(*2*) *landau-o.big-trans* **by** *blast*
qed

lemma *landau-ln-2*:
assumes $a > (1::real)$
assumes *eventually* $(\lambda x.\ 1 \leq f\ x)\ F$
assumes *eventually* $(\lambda x.\ a \leq g\ x)\ F$
assumes $f \in O[F](g)$

```
  shows (λx. ln (f x)) ∈ O[F](λx. ln (g x))
proof −
  obtain c where a: c > 0 and b: eventually (λx. abs (f x) ≤ c ∗ abs (g x)) F
    using assms(4) by (simp add:bigo-def, blast)
  define d where d = 1 + (max 0 (ln c)) / ln a
  have d:eventually (λx. abs (ln (f x)) ≤ d ∗ abs (ln (g x))) F
  proof (rule eventually-mono[OF eventually-conj[OF b eventually-conj[OF assms(3,2)]]])
    fix x
    assume c:|f x| ≤ c ∗ |g x| ∧ a ≤ g x ∧ 1 ≤ f x
    have abs (ln (f x)) = ln (f x)
      by (subst abs-of-nonneg, rule ln-ge-zero, metis c, simp)
    also have ... ≤ ln (c ∗ abs (g x))
      using c assms(1)  mult-pos-pos[OF a] by auto
    also have ... ≤ ln c + ln (abs (g x))
      using c assms(1) by (simp add: a ln-mult-pos)
    also have ... ≤ (d−1)∗ln a + ln (g x)
      using assms(1) c
      by (intro add-mono iffD2[OF ln-le-cancel-iff], simp-all add:d-def)
    also have ... ≤ (d−1)∗ ln (g x) + ln (g x)
      using assms(1) c
     by (intro add-mono mult-left-mono iffD2[OF ln-le-cancel-iff], simp-all add:d-def)
    also have ... = d ∗ ln (g x) by (simp add:algebra-simps)
    also have ... = d ∗ abs (ln (g x))
      using c assms(1) by auto
    finally show abs (ln (f x)) ≤ d ∗ abs (ln (g x)) by simp
  qed
  hence ∀F x in F. |ln (f x)| ≤ d ∗ |ln (g x)|
    by simp
  moreover have 0 < d
    unfolding d-def using assms(1)
    by (intro add-pos-nonneg divide-nonneg-pos, auto)
  ultimately show ?thesis
    by (auto simp:bigo-def)
qed

lemma landau-real-nat:
  fixes f :: 'a ⇒ int
  assumes (λx. of-int (f x)) ∈ O[F](g)
  shows (λx. real (nat (f x))) ∈ O[F](g)
proof −
  obtain c where a: c > 0 and b: eventually (λx. abs (of-int (f x)) ≤ c ∗ abs (g
x)) F
    using assms(1) by (simp add:bigo-def, blast)
  have ∀F x in F. real (nat (f x)) ≤ c ∗ |g x|
    by (rule eventually-mono[OF b],  simp)
  thus ?thesis using a
    by (auto simp:bigo-def)
qed
```

```
lemma landau-ceil:
  assumes (λ-. 1) ∈ O[F'](g)
  assumes f ∈ O[F'](g)
  shows (λx. real-of-int ⌈f x⌉) ∈ O[F'](g)
proof −
  have (λx. real-of-int ⌈f x⌉) ∈ O[F'](λx. 1 + abs (f x))
    by (intro landau-o.big-mono always-eventually allI, simp, linarith)
  also have (λx. 1 + abs(f x)) ∈ O[F'](g)
    using assms(2) by (intro sum-in-bigo assms(1), auto)
  finally show ?thesis by simp
qed

lemma landau-rat-ceil:
  assumes (λ-. 1) ∈ O[F'](g)
  assumes (λx. real-of-rat (f x)) ∈ O[F'](g)
  shows (λx. real-of-int ⌈f x⌉) ∈ O[F'](g)
proof −
  have a:|real-of-int ⌈x⌉| ≤ 1 + real-of-rat |x| for x :: rat
  proof (cases x ≥ 0)
    case True
    then show ?thesis
      by (simp, metis add.commute of-int-ceiling-le-add-one of-rat-ceiling)
  next
    case False
    have real-of-rat x − 1 ≤ real-of-rat x
      by simp
    also have ... ≤ real-of-int ⌈x⌉
      by (metis ceiling-correct of-rat-ceiling)
    finally have real-of-rat (x)−1 ≤ real-of-int ⌈x⌉ by simp

    hence − real-of-int ⌈x⌉ ≤ 1 + real-of-rat (− x)
      by (simp add: of-rat-minus)
    then show ?thesis using False by simp
  qed
  have (λx. real-of-int ⌈f x⌉) ∈ O[F'](λx. 1 + abs (real-of-rat (f x)))
    using a
    by (intro landau-o.big-mono always-eventually allI, simp)
  also have (λx. 1 + abs (real-of-rat (f x))) ∈ O[F'](g)
    using assms
    by (intro sum-in-bigo assms(1), subst landau-o.big.abs-in-iff, simp)
  finally show ?thesis by simp
qed

lemma landau-nat-ceil:
  assumes (λ-. 1) ∈ O[F'](g)
  assumes f ∈ O[F'](g)
  shows (λx. real (nat ⌈f x⌉)) ∈ O[F'](g)
  using assms
  by (intro landau-real-nat landau-ceil, auto)
```

lemma *eventually-prod1'*:
  assumes $B \neq bot$
  assumes $(\forall_F\ x\ in\ A.\ P\ x)$
  shows $(\forall_F\ x\ in\ A \times_F B.\ P\ (fst\ x))$
proof −
  have $(\forall_F\ x\ in\ A \times_F B.\ P\ (fst\ x)) = (\forall_F\ (x,y)\ in\ A \times_F B.\ P\ x)$
    by (*simp add*:*case-prod-beta'*)
  also have ... $= (\forall_F\ x\ in\ A.\ P\ x)$
    by (*subst eventually-prod1*[*OF assms*(*1*)], *simp*)
  finally show *?thesis* **using** *assms*(*2*) **by** *simp*
qed

lemma *eventually-prod2'*:
  assumes $A \neq bot$
  assumes $(\forall_F\ x\ in\ B.\ P\ x)$
  shows $(\forall_F\ x\ in\ A \times_F B.\ P\ (snd\ x))$
proof −
  have $(\forall_F\ x\ in\ A \times_F B.\ P\ (snd\ x)) = (\forall_F\ (x,y)\ in\ A \times_F B.\ P\ y)$
    by (*simp add*:*case-prod-beta'*)
  also have ... $= (\forall_F\ x\ in\ B.\ P\ x)$
    by (*subst eventually-prod2*[*OF assms*(*1*)], *simp*)
  finally show *?thesis* **using** *assms*(*2*) **by** *simp*
qed

lemma *sequentially-inf*: $\forall_F\ x\ in\ sequentially.\ n \leq real\ x$
  by (*meson eventually-at-top-linorder nat-ceiling-le-eq*)

instantiation *rat* :: *linorder-topology*
begin

definition *open-rat* :: *rat set* ⇒ *bool*
  where *open-rat* = *generate-topology* (*range* ($\lambda a.$ {..< *a*}) ∪ *range* ($\lambda a.$ {*a* <..}))

instance
  by *standard* (*rule open-rat-def*)
end

lemma *inv-at-right-0-inf*:
  $\forall_F\ x\ in\ at\text{-}right\ 0.\ c \leq 1\ /\ real\text{-}of\text{-}rat\ x$
proof −
  have *a*: $c \leq 1\ /\ real\text{-}of\text{-}rat\ x$ **if** *b*: $x \in \{0 <..< 1\ /\ rat\text{-}of\text{-}int\ (max\ \lceil c \rceil\ 1)\}$ **for** *x*
  proof −
    have $c * real\text{-}of\text{-}rat\ x \leq real\text{-}of\text{-}int\ (max\ \lceil c \rceil\ 1) * real\text{-}of\text{-}rat\ x$
      using *b* **by** (*intro mult-right-mono*, *linarith*, *auto*)
    also have ... $< real\text{-}of\text{-}int\ (max\ \lceil c \rceil\ 1) * real\text{-}of\text{-}rat\ (1/rat\text{-}of\text{-}int\ (max\ \lceil c \rceil\ 1))$
      using *b* **by** (*intro mult-strict-left-mono iffD2*[*OF of-rat-less*], *auto*)
    also have ... ≤ 1

by (*simp add*:*of-rat-divide*)
finally have $c * \textit{real-of-rat}\ x \le 1$ **by** *simp*
moreover have $0 < \textit{real-of-rat}\ x$
using *b* **by** *simp*
ultimately show *?thesis* **by** (*subst pos-le-divide-eq*, *auto*)
qed

show *?thesis*
using *a*
by (*intro eventually-at-rightI*[**where** $b=1/\textit{rat-of-int}\ (\max \lceil c \rceil\ 1)$], *simp-all*)
qed

end

5 Probability Spaces

Some additional results about probability spaces in addition to "HOL-Probability".

theory *Probability-Ext*
imports
HOL−Probability.Stream-Space
Concentration-Inequalities.Bienaymes-Identity
Universal-Hash-Families.Carter-Wegman-Hash-Family
Frequency-Moments-Preliminary-Results
begin

context *prob-space*
begin

lemma *pmf-mono*:
assumes $M = \textit{measure-pmf}\ p$
assumes $\bigwedge x.\ x \in P \Longrightarrow x \in \textit{set-pmf}\ p \Longrightarrow x \in Q$
shows $\textit{prob}\ P \le \textit{prob}\ Q$
proof −
have $\textit{prob}\ P = \textit{prob}\ (P \cap (\textit{set-pmf}\ p))$
by (*rule measure-pmf-eq*[*OF assms*(*1*)], *blast*)
also have ... $\le \textit{prob}\ Q$
using *assms* **by** (*intro finite-measure.finite-measure-mono*, *auto*)
finally show *?thesis* **by** *simp*
qed

lemma *pmf-add*:
assumes $M = \textit{measure-pmf}\ p$
assumes $\bigwedge x.\ x \in P \Longrightarrow x \in \textit{set-pmf}\ p \Longrightarrow x \in Q \lor x \in R$
shows $\textit{prob}\ P \le \textit{prob}\ Q + \textit{prob}\ R$
proof −
have [*simp*]:$\textit{events} = \textit{UNIV}$ **by** (*subst assms*(*1*), *simp*)
have $\textit{prob}\ P \le \textit{prob}\ (Q \cup R)$
using *assms* **by** (*intro pmf-mono*[*OF assms*(*1*)], *blast*)

also have ... $\leq$ *prob Q* + *prob R*
by (*rule measure-subadditive*, *auto*)
finally show *?thesis* **by** *simp*
qed

lemma *pmf-add-2*:
assumes M = *measure-pmf p*
assumes *prob* $\{\omega.\ P\ \omega\} \leq r1$
assumes *prob* $\{\omega.\ Q\ \omega\} \leq r2$
shows *prob* $\{\omega.\ P\ \omega \vee Q\ \omega\} \leq r1 + r2$ (**is** *?lhs* $\leq$ *?rhs*)
proof −
have *?lhs* $\leq$ *prob* $\{\omega.\ P\ \omega\}$ + *prob* $\{\omega.\ Q\ \omega\}$
by (*intro pmf-add*[*OF assms*(*1*)], *auto*)
also have ... $\leq$ *?rhs*
by (*intro add-mono assms*(*2*−*3*))
finally show *?thesis*
by *simp*
qed

end

end

6 Frequency Moment 0

theory *Frequency-Moment-0*
imports
Frequency-Moments-Preliminary-Results
Median-Method.Median
K-Smallest
Universal-Hash-Families.Carter-Wegman-Hash-Family
Frequency-Moments
Landau-Ext
Probability-Ext
Universal-Hash-Families.Universal-Hash-Families-More-Product-PMF
begin

This section contains a formalization of a new algorithm for the zero-th frequency moment inspired by ideas described in [2]. It is a KMV-type (k-minimum value) algorithm with a rounding method and matches the space complexity of the best algorithm described in [2].

In addition to the Isabelle proof here, there is also an informal hand-written proof in Appendix A.

type-synonym *f0-state* = *nat* × *nat* × *nat* × *nat* × (*nat* ⇒ *nat list*) × (*nat* ⇒ *float set*)

definition *hash* **where** *hash p* = *ring.hash* (*ring-of* (*mod-ring p*))

```
fun f0-init :: rat ⇒ rat ⇒ nat ⇒ f0-state pmf where
  f0-init δ ε n =
    do {
      let s = nat ⌈−18 ∗ ln (real-of-rat ε)⌉;
      let t = nat ⌈80 / (real-of-rat δ)^2⌉;
      let p = prime-above (max n 19);
      let r = nat (4 ∗ ⌈log 2 (1 / real-of-rat δ)⌉ + 23);
        h ← prod-pmf {..<s} (λ-. pmf-of-set (bounded-degree-polynomials (ring-of
(mod-ring p)) 2));
      return-pmf (s, t, p, r, h, (λ- ∈ {0..<s}. {}))
    }

fun f0-update :: nat ⇒ f0-state ⇒ f0-state pmf where
  f0-update x (s, t, p, r, h, sketch) =
    return-pmf (s, t, p, r, h, λi ∈ {..<s}.
      least t (insert (float-of (truncate-down r (hash p x (h i)))) (sketch i)))

fun f0-result :: f0-state ⇒ rat pmf where
  f0-result (s, t, p, r, h, sketch) = return-pmf (median s (λi ∈ {..<s}.
      (if card (sketch i) < t then of-nat (card (sketch i)) else
        rat-of-nat t∗ rat-of-nat p / rat-of-float (Max (sketch i)))
    ))

fun f0-space-usage :: (nat × rat × rat) ⇒ real where
  f0-space-usage (n, ε, δ) = (
    let s = nat ⌈−18 ∗ ln (real-of-rat ε)⌉ in
    let r = nat (4 ∗ ⌈log 2 (1 / real-of-rat δ)⌉ + 23) in
    let t = nat ⌈80 / (real-of-rat δ)^2 ⌉ in
    6 +
    2 ∗ log 2 (real s + 1) +
    2 ∗ log 2 (real t + 1) +
    2 ∗ log 2 (real n + 21) +
    2 ∗ log 2 (real r + 1) +
    real s ∗ (5 + 2 ∗ log 2 (21 + real n) +
    real t ∗ (13 + 4 ∗ r + 2 ∗ log 2 (log 2 (real n + 13)))))

definition encode-f0-state :: f0-state ⇒ bool list option where
  encode-f0-state =
    N_e ⋈_e (λs.
    N_e ×_e (
    N_e ⋈_e (λp.
    N_e ×_e (
    ([0..<s] →_e (P_e p 2)) ×_e
    ([0..<s] →_e (S_e F_e))))))

lemma inj-on encode-f0-state (dom encode-f0-state)
proof −
  have is-encoding encode-f0-state
    unfolding encode-f0-state-def
```

```
    by (intro dependent-encoding exp-golomb-encoding poly-encoding fun-encoding
set-encoding float-encoding)
  thus ?thesis by (rule encoding-imp-inj)
qed

context
  fixes ε δ :: rat
  fixes n :: nat
  fixes as :: nat list
  fixes result
  assumes ε-range: ε ∈ {0<..<1}
  assumes δ-range: δ ∈ {0<..<1}
  assumes as-range: set as ⊆ {..<n}
  defines result ≡ fold (λa state. state ≫= f0-update a) as (f0-init δ ε n) ≫=
f0-result
begin

private definition t where t = nat ⌈80 / (real-of-rat δ)^2⌉
private lemma t-gt-0: t > 0 using δ-range by (simp add:t-def)

private definition s where s = nat ⌈−(18 ∗ ln (real-of-rat ε))⌉
private lemma s-gt-0: s > 0 using ε-range by (simp add:s-def)

private definition p where p = prime-above (max n 19)

private lemma p-prime:Factorial-Ring.prime p
  using p-def prime-above-prime by presburger

private lemma p-ge-18: p ≥ 18
proof −
  have p ≥ 19
    by (metis p-def prime-above-lower-bound max.bounded-iff)
  thus ?thesis by simp
qed

private lemma p-gt-0: p > 0 using p-ge-18 by simp
private lemma p-gt-1: p > 1 using p-ge-18 by simp

private lemma n-le-p: n ≤ p
proof −
  have n ≤ max n 19 by simp
  also have ... ≤ p
    unfolding p-def by (rule prime-above-lower-bound)
  finally show ?thesis by simp
qed

private lemma p-le-n: p ≤ 2∗n + 40
proof −
  have p ≤ 2 ∗ (max n 19) + 2
```

```
    by (subst p-def, rule prime-above-upper-bound)
  also have ... ≤ 2 ∗ n + 40
    by (cases n ≥ 19, auto)
  finally show ?thesis by simp
qed

private lemma as-lt-p: ⋀x. x ∈ set as ⟹ x < p
  using as-range atLeastLessThan-iff
  by (intro order-less-le-trans[OF - n-le-p]) blast

private lemma as-subset-p: set as ⊆ {..<p}
   using as-lt-p  by (simp add: subset-iff)

private definition r where r = nat (4 ∗ ⌈log 2 (1 / real-of-rat δ)⌉ + 23)

private lemma r-bound: 4 ∗ log 2 (1 / real-of-rat δ) + 23 ≤ r
proof −
  have 0 ≤ log 2 (1 / real-of-rat δ) using δ-range by simp
  hence 0 ≤ ⌈log 2 (1 / real-of-rat δ)⌉ by simp
  hence 0 ≤ 4 ∗ ⌈log 2 (1 / real-of-rat δ)⌉ + 23
    by (intro add-nonneg-nonneg mult-nonneg-nonneg, auto)
  thus ?thesis by (simp add:r-def)
qed

private lemma r-ge-23: r ≥ 23
proof −
  have (23::real) = 0 + 23 by simp
  also have ... ≤ 4 ∗ log 2 (1 / real-of-rat δ) + 23
    using δ-range by (intro add-mono mult-nonneg-nonneg, auto)
  also have ... ≤ r using r-bound by simp
  finally show 23 ≤ r by simp
qed

private lemma two-pow-r-le-1: 0 < 1 − 2 powr − real r
proof −
  have a: 2 powr (0::real) = 1
    by simp
  show ?thesis using r-ge-23
    by (simp, subst a[symmetric], intro powr-less-mono, auto)
qed

interpretation carter-wegman-hash-family ring-of (mod-ring p) 2
  rewrites ring.hash (ring-of (mod-ring p)) = Frequency-Moment-0.hash p
  using carter-wegman-hash-familyI[OF mod-ring-is-field mod-ring-finite]
  using hash-def p-prime by auto

private definition tr-hash where tr-hash x ω = truncate-down r (hash x ω)

private definition sketch-rv where
```

sketch-rv ω = *least t* ((λx. *float-of* (*tr-hash x* ω)) ' *set as*)

private definition *estimate*
where *estimate S* = (*if card S* < *t then of-nat* (*card S*) *else of-nat t* $*$ *of-nat p* / *rat-of-float* (*Max S*))

private definition *sketch-rv'* **where** *sketch-rv'* ω = *least t* ((λx. *tr-hash x* ω) ' *set as*)
private definition *estimate'* **where** *estimate' S* = (*if card S* < *t then real* (*card S*) *else real t* $*$ *real p* / *Max S*)

private definition Ω_0 **where** Ω_0 = *prod-pmf* {..<*s*} (λ-. *pmf-of-set space*)

private lemma *f0-alg-sketch*:
defines *sketch* ≡ *fold* (λ*a state. state* ≫= *f0-update a*) *as* (*f0-init* δ ε *n*)
shows *sketch* = *map-pmf* (λx. (*s*,*t*,*p*,*r*, *x*, λi ∈ {..<*s*}. *sketch-rv* (*x i*))) Ω_0
unfolding *sketch-rv-def*
proof (*subst sketch-def*, *induction as rule*:*rev-induct*)
case *Nil*
then show *?case*
by (*simp add*:*s-def p-def*[*symmetric*] *map-pmf-def t-def r-def Let-def least-def restrict-def space-def* Ω_0-*def*)
next
case (*snoc x xs*)
let *?sketch* = $\lambda\omega$ *xs*. *least t* ((λa. *float-of* (*tr-hash a* ω)) ' *set xs*)
have *fold* (λ*a state. state* ≫= *f0-update a*) (*xs* @ [*x*]) (*f0-init* δ ε *n*) =
(*map-pmf* ($\lambda\omega$. (*s*, *t*, *p*, *r*, ω, λi ∈ {..<*s*}. *?sketch* (ω *i*) *xs*)) Ω_0) ≫= *f0-update x*
by (*simp add*: *restrict-def snoc del*:*f0-init.simps*)
also have ... = Ω_0 ≫= ($\lambda\omega$. *f0-update x* (*s*, *t*, *p*, *r*, ω, λi∈{..<*s*}. *?sketch* (ω *i*) *xs*))
by (*simp add*:*map-pmf-def bind-assoc-pmf bind-return-pmf del*:*f0-update.simps*)
also have ... = *map-pmf* ($\lambda\omega$. (*s*, *t*, *p*, *r*, ω, λi∈{..<*s*}. *?sketch* (ω *i*) (*xs*@[*x*]))) Ω_0
by (*simp add*:*least-insert map-pmf-def tr-hash-def cong*:*restrict-cong*)
finally show *?case* **by** *blast*
qed

private lemma *card-nat-in-ball*:
fixes *x* :: *nat*
fixes *q* :: *real*
assumes $q \geq 0$
defines *A* ≡ {*k*. *abs* (*real x* − *real k*) ≤ *q* ∧ *k* ≠ *x*}
shows *real* (*card A*) ≤ *2* $*$ *q* **and** *finite A*
proof −
have *a*: *of-nat x* ∈ {⌈*real x*−*q*⌉..⌊*real x*+*q*⌋}
using *assms*
by (*simp add*: *ceiling-le-iff*)

have *card A* = *card* (*int* ' *A*)
  by (*rule card-image*[*symmetric*], *simp*)
also have ... $\le$ *card* ({⌈*real x*−*q*⌉..⌊*real x*+*q*⌋} − {*of-nat x*})
  by (*intro card-mono image-subsetI*, *simp-all add*:*A-def abs-le-iff*, *linarith*)
also have ... = *card* {⌈*real x*−*q*⌉..⌊*real x*+*q*⌋} − *1*
  by (*rule card-Diff-singleton*, *rule a*)
also have ... = *int* (*card* {⌈*real x*−*q*⌉..⌊*real x*+*q*⌋}) − *int 1*
  by (*intro of-nat-diff*)
   (*metis a card-0-eq empty-iff finite-atLeastAtMost-int less-one linorder-not-le*)
also have ... $\le$ ⌊*q*+*real x*⌋+*1* −⌈*real x*−*q*⌉ − *1*
  using *assms* **by** (*simp*, *linarith*)
also have ... $\le$ *2*∗*q*
  by *linarith*
finally show *card A* $\le$ *2* ∗ *q*
  by *simp*

have *A* ⊆ {..*x* + *nat* ⌈*q*⌉}
  by (*rule subsetI*, *simp add*:*A-def abs-le-iff*, *linarith*)
thus *finite A*
  by (*rule finite-subset*, *simp*)
qed

private lemma *prob-degree-lt-1*:
  prob {ω. *degree* ω < *1*} $\le$ *1*/*real p*
proof −
  have *space* ∩ {ω. *length* ω $\le$ *Suc 0*} = *bounded-degree-polynomials* (*ring-of* (*mod-ring p*)) *1*
  by (*auto simp*:*set-eq-iff bounded-degree-polynomials-def space-def*)
moreover have *field-size* = *p* **by** (*simp add*:*ring-of-def mod-ring-def*)
  hence *real* (*card* (*bounded-degree-polynomials* (*ring-of* (*mod-ring p*)) *1*))/*card space* = *1* / *real p*
  by (*simp add*:*space-def bounded-degree-polynomials-card power2-eq-square*)
ultimately show *?thesis*
  by (*simp add*:*M-def measure-pmf-of-set*)
qed

private lemma *collision-prob*:
  assumes *c* $\ge$ *1*
  shows *prob* {ω. ∃ *x* ∈ *set as*. ∃ *y* ∈ *set as*. *x* ≠ *y* ∧ *tr-hash x* ω $\le$ *c* ∧ *tr-hash x* ω = *tr-hash y* ω} $\le$
    (*5*/*2*) ∗ (*real* (*card* (*set as*)))^2 ∗ c^2 ∗ *2 powr* −(*real r*) / (*real p*)^2 + *1*/*real p*
(**is** *prob* {ω. *?l* ω} $\le$ *?r1* + *?r2*)
proof −
  define ϱ :: *real* **where** ϱ = *9*/*8*

  have *rho-c-ge-0*: ϱ ∗ *c* $\ge$ *0* **unfolding** *ϱ-def* **using** *assms* **by** *simp*

  have *c-ge-0*: *c*$\ge$*0* **using** *assms* **by** *simp*

have *degree* $\omega \geq 1 \Longrightarrow \omega \in$ *space* $\Longrightarrow$ *degree* $\omega = 1$ **for** ω
by (*simp add*:*bounded-degree-polynomials-def space-def*)
(*metis One-nat-def Suc-1 le-less-Suc-eq less-imp-diff-less list.size*(*3*) *pos2*)

hence *a*: $\bigwedge\omega\ x\ y.\ x < p \Longrightarrow y < p \Longrightarrow x \neq y \Longrightarrow$ *degree* $\omega \geq 1 \Longrightarrow \omega \in$ *space* $\Longrightarrow$ *hash* $x\ \omega \neq$ *hash* $y\ \omega$
using *inj-onD*[*OF inj-if-degree-1*] *mod-ring-carr* **by** *blast*

have *b*: *prob* $\{\omega.$ *degree* $\omega \geq 1 \wedge$ *tr-hash* $x\ \omega \leq c \wedge$ *tr-hash* $x\ \omega =$ *tr-hash* $y\ \omega\}$ $\leq 5 * c^2 * 2$ *powr* $(-$*real* $r)$ /(*real* $p)^2$
if *b-assms*: $x \in$ *set as* $y \in$ *set as* $x < y$ **for** $x\ y$
proof −
have *c*: *real* $u \leq \varrho * c \wedge |$*real* $u -$ *real* $v| \leq \varrho * c * 2$ *powr* $(-$*real* $r)$
if *c-assms*:*truncate-down* r (*real* u) $\leq c$ *truncate-down* r (*real* u) $=$ *truncate-down* r (*real* v) **for** $u\ v$
proof −
have $9 * 2$ *powr* $-$ *real* $r \leq 9 * 2$ *powr* $(-$ *real* $23)$
using *r-ge-23* **by** (*intro mult-left-mono powr-mono*, *auto*)

also have ... ≤ 1 **by** *simp*

finally have $9 * 2$ *powr* $-$ *real* $r \leq 1$ **by** *simp*

hence $1 \leq \varrho * (1 - 2$ *powr* $(-$ *real* $r))$
by (*simp add*:*ϱ-def*)

hence *d*: $(c*1)\ /\ (1 - 2$ *powr* $(-$*real* $r)) \leq c * \varrho$
using *assms two-pow-r-le-1* **by** (*simp add*: *pos-divide-le-eq*)

have $\bigwedge x.$ *truncate-down* r (*real* x) $\leq c \Longrightarrow$ *real* $x * (1 - 2$ *powr* $-$ *real* $r) \leq c * 1$
using *truncate-down-pos*[*OF of-nat-0-le-iff*] *order-trans* **by** (*simp*, *blast*)

hence $\bigwedge x.$ *truncate-down* r (*real* x) $\leq c \Longrightarrow$ *real* $x \leq c * \varrho$
using *two-pow-r-le-1* **by** (*intro order-trans*[*OF - d*], *simp add*: *pos-le-divide-eq*)

hence *e*: *real* $u \leq c * \varrho$ *real* $v \leq c * \varrho$
using *c-assms* **by** *auto*

have $|$*real* $u -$ *real* $v| \leq ($*max* $|$*real* $u|\ |$*real* $v|) * 2$ *powr* $(-$*real* $r)$
using *c-assms* **by** (*intro truncate-down-eq*, *simp*)

also have ... $\leq (c * \varrho) * 2$ *powr* $(-$*real* $r)$
using *e* **by** (*intro mult-right-mono*, *auto*)

finally have $|$*real* $u -$ *real* $v| \leq \varrho * c * 2$ *powr* $(-$*real* $r)$
by (*simp add*:*algebra-simps*)

thus *?thesis* **using** *e* **by** (*simp add*:*algebra-simps*)

qed

have *prob* $\{\omega.\ degree\ \omega \geq 1 \wedge tr\text{-}hash\ x\ \omega \leq c \wedge tr\text{-}hash\ x\ \omega = tr\text{-}hash\ y\ \omega\} \leq$
prob $(\bigcup\ i \in \{(u,v) \in \{..<p\} \times \{..<p\}.\ u \neq v \wedge truncate\text{-}down\ r\ u \leq c \wedge truncate\text{-}down\ r\ u = truncate\text{-}down\ r\ v\}.$
$\{\omega.\ hash\ x\ \omega = fst\ i \wedge hash\ y\ \omega = snd\ i\})$
using *a* **by** (*intro pmf-mono*[*OF M-def*], *simp add*:*tr-hash-def*)
(*metis hash-range mod-ring-carr b-assms as-subset-p lessThan-iff nat-neq-iff subset-eq*)

also have ... $\leq (\sum\ i \in \{(u,v) \in \{..<p\} \times \{..<p\}.\ u \neq v \wedge$
$truncate\text{-}down\ r\ u \leq c \wedge truncate\text{-}down\ r\ u = truncate\text{-}down\ r\ v\}.$
$prob\ \{\omega.\ hash\ x\ \omega = fst\ i \wedge hash\ \ y\ \omega = snd\ i\})$
by (*intro measure-UNION-le finite-cartesian-product finite-subset*[**where** $B=\{0..<p\} \times \{0..<p\}$])
(*auto simp add*:*M-def*)

also have ... $\leq (\sum\ i \in \{(u,v) \in \{..<p\} \times \{..<p\}.\ u \neq v \wedge$
$truncate\text{-}down\ r\ u \leq c \wedge truncate\text{-}down\ r\ u = truncate\text{-}down\ r\ v\}.$
$prob\ \{\omega.\ (\forall\, u \in \{x,y\}.\ hash\ u\ \omega = (if\ u = x\ then\ (fst\ i)\ else\ (snd\ i)))\})$
by (*intro sum-mono pmf-mono*[*OF M-def*]) *force*

also have ... $\leq (\sum\ i \in \{(u,v) \in \{..<p\} \times \{..<p\}.\ u \neq v \wedge$
$truncate\text{-}down\ r\ u \leq c \wedge truncate\text{-}down\ r\ u = truncate\text{-}down\ r\ v\}.\ 1/(real\ p)^2)$
using *assms as-subset-p b-assms*
by (*intro sum-mono*, *subst hash-prob*) (*auto simp*: *ring-of-def mod-ring-def power2-eq-square*)

also have ... $= 1/(real\ p)^2\ *$
$card\ \{(u,v) \in \{0..<p\} \times \{0..<p\}.\ u \neq v \wedge truncate\text{-}down\ r\ u \leq c \wedge truncate\text{-}down\ r\ u = truncate\text{-}down\ r\ v\}$
by *simp*

also have ... $\leq 1/(real\ p)^2\ *$
$card\ \{(u,v) \in \{..<p\} \times \{..<p\}.\ u \neq v \wedge real\ u \leq \varrho * c \wedge abs\ (real\ u - real\ v) \leq \varrho * c * 2\ powr\ (-real\ r)\}$
using *c*
by (*intro mult-mono of-nat-mono card-mono finite-cartesian-product finite-subset*[**where** $B=\{..<p\}\times\{..<p\}$])
auto

also have ... $\leq 1/(real\ p)^2 * card\ (\bigcup u' \in \{u.\ u < p \wedge real\ u \leq \varrho * c\}.$
$\{(u::nat,v::nat).\ u = u' \wedge abs\ (real\ u - real\ v) \leq \varrho * c * 2\ powr\ (-real\ r) \wedge v < p \wedge v \neq u'\})$
by (*intro mult-left-mono of-nat-mono card-mono finite-cartesian-product finite-subset*[**where** $B=\{..<p\}\times\{..<p\}$])
auto

also have ... $\leq$ *1/(real p)*^{2} $*$ ($\sum$ $u' \in \{u.\ u < p \wedge$ *real* $u \leq \varrho * c\}$.
card $\{(u,v).\ u = u' \wedge$ *abs* (*real* u $-$ *real* v) $\leq \varrho * c * 2$ *powr* ($-$*real* r) $\wedge$ $v < p \wedge v \neq u'\}$)
by (*intro mult-left-mono of-nat-mono card-UN-le*, *auto*)

also have ... $=$ *1/(real p)*^{2} $*$ ($\sum$ $u' \in \{u.\ u < p \wedge$ *real* $u \leq \varrho * c\}$.
card (($\lambda x.\ (u',x)$) ` $\{v.$ *abs* (*real* u' $-$ *real* v) $\leq \varrho * c * 2$ *powr* ($-$*real* r) $\wedge$ $v < p \wedge v \neq u'\}$))
by (*intro arg-cong2*[**where** *f*=($*$)] *arg-cong*[**where** *f*=*real*] *sum.cong arg-cong*[**where** *f*=*card*])
(*auto simp add*:*set-eq-iff*)

also have ... $\leq$ *1/(real p)*^{2} $*$ ($\sum$ $u' \in \{u.\ u < p \wedge$ *real* $u \leq \varrho * c\}$.
card $\{v.$ *abs* (*real* u' $-$ *real* v) $\leq \varrho * c * 2$ *powr* ($-$*real* r) $\wedge$ $v < p \wedge v \neq u'\}$)
by (*intro mult-left-mono of-nat-mono sum-mono card-image-le*, *auto*)

also have ... $\leq$ *1/(real p)*^{2} $*$ ($\sum$ $u' \in \{u.\ u < p \wedge$ *real* $u \leq \varrho * c\}$.
card $\{v.$ *abs* (*real* u' $-$ *real* v) $\leq \varrho * c * 2$ *powr* ($-$*real* r) $\wedge$ $v \neq u'\}$)
by (*intro mult-left-mono sum-mono of-nat-mono card-mono card-nat-in-ball subsetI*) *auto*

also have ... $\leq$ *1/(real p)*^{2} $*$ ($\sum$ $u' \in \{u.\ u < p \wedge$ *real* $u \leq \varrho * c\}$.
real (*card* $\{v.$ *abs* (*real* u' $-$ *real* v) $\leq \varrho * c * 2$ *powr* ($-$*real* r) $\wedge$ $v \neq u'\}$))
by *simp*

also have ... $\leq$ *1/(real p)*^{2} $*$ ($\sum$ $u' \in \{u.\ u < p \wedge$ *real* $u \leq \varrho * c\}$. *2* $*$ ($\varrho * c * 2$ *powr* ($-$*real* r)))
by (*intro mult-left-mono sum-mono card-nat-in-ball*(*1*), *auto*)

also have ... $=$ *1/(real p)*^{2} $*$ (*real* (*card* $\{u.\ u < p \wedge$ *real* $u \leq \varrho * c\}$) $*$ (*2* $*$ ($\varrho * c * 2$ *powr* ($-$*real* r))))
by *simp*

also have ... $\leq$ *1/(real p)*^{2} $*$ (*real* (*card* $\{u.\ u \leq$ *nat* ($\lfloor \varrho * c \rfloor$)$\}$) $*$ (*2* $*$ ($\varrho * c * 2$ *powr* ($-$*real* r))))
using *rho-c-ge-0 le-nat-floor*
by (*intro mult-left-mono mult-right-mono of-nat-mono card-mono subsetI*) *auto*

also have ... $\leq$ *1/(real p)*^{2} $*$ ((*1*+$\varrho * c$) $*$ (*2* $*$ ($\varrho * c * 2$ *powr* ($-$*real* r))))
using *rho-c-ge-0* **by** (*intro mult-left-mono mult-right-mono*, *auto*)

also have ... $\leq$ *1/(real p)*^{2} $*$ (((*1*+ϱ) $*$ c) $*$ (*2* $*$ ($\varrho * c * 2$ *powr* ($-$*real* r))))
using *assms* **by** (*intro mult-mono*, *auto simp add*:*distrib-left distrib-right* *ϱ-def*)

also have ... $=$ (ϱ $*$ (*2* $+$ ϱ $*$ *2*)) $*$ c^2 $*$ *2 powr* ($-$*real* r) /(*real* p)^2
by (*simp add*:*ac-simps power2-eq-square*)

also have ... $\leq$ *5* $*$ c^2 $*$ *2 powr* $(-$*real r*$)$ $/($*real p*$)^2$
by (*intro divide-right-mono mult-right-mono*) (*auto simp add*:*ϱ-def*)

finally show *?thesis* **by** *simp*
qed

have *prob* $\{\omega.$ *?l* $\omega \wedge$ *degree* $\omega \geq$ *1*$\} \leq$
prob $(\bigcup\ i \in \{(x,y) \in$ (*set as*) $\times$ (*set as*). $x < y\}$. $\{\omega.$ *degree* $\omega \geq$ *1* $\wedge$ *tr-hash* (*fst i*) $\omega \leq c \wedge$
tr-hash (*fst i*) $\omega =$ *tr-hash* (*snd i*) $\omega\})$
by (*rule pmf-mono*[*OF M-def*], *simp*, *metis linorder-neqE-nat*)

also have ... $\leq (\sum\ i \in \{(x,y) \in$ (*set as*) $\times$ (*set as*). $x < y\}$. *prob*
$\{\omega.$ *degree* $\omega \geq$ *1* $\wedge$ *tr-hash* (*fst i*) $\omega \leq c \wedge$ *tr-hash* (*fst i*) $\omega =$ *tr-hash* (*snd i*) $\omega\})$
unfolding *M-def*
by (*intro measure-UNION-le finite-cartesian-product finite-subset*[**where** *B*=(*set as*) $\times$ (*set as*)])
auto

also have ... $\leq (\sum\ i \in \{(x,y) \in$ (*set as*) $\times$ (*set as*). $x < y\}$. *5* $*$ c^2 $*$ *2 powr* $(-$*real r*$)$ $/($*real p*$)^2)$
using *b* **by** (*intro sum-mono*, *simp add*:*case-prod-beta*)

also have ... $=$ $((5/2) * c^2$ $*$ *2 powr* $(-$*real r*$)$ $/($*real p*$)^2) * (2 *$ *card* $\{(x,y) \in$ (*set as*) $\times$ (*set as*). $x < y\})$
by *simp*

also have ... $=$ $((5/2) * c^2$ $*$ *2 powr* $(-$*real r*$)$ $/($*real p*$)^2) *$ (*card* (*set as*) $*$ (*card* (*set as*) $-$ *1*))
by (*subst card-ordered-pairs*, *auto*)

also have ... $\leq$ $((5/2) * c^2$ $*$ *2 powr* $(-$*real r*$)$ $/($*real p*$)^2) *$ (*real* (*card* (*set as*)))^2
by (*intro mult-left-mono*) (*auto simp add*:*power2-eq-square mult-left-mono*)

also have ... $=$ $(5/2) *$ (*real* (*card* (*set as*)))^2 $* c^2 *$ *2 powr* $(-$*real r*$)$ $/($*real p*$)^2$
by (*simp add*:*algebra-simps*)

finally have *f*:*prob* $\{\omega.$ *?l* $\omega \wedge$ *degree* $\omega \geq$ *1*$\} \leq$ *?r1* **by** *simp*

have *prob* $\{\omega.$ *?l* $\omega\} \leq$ *prob* $\{\omega.$ *?l* $\omega \wedge$ *degree* $\omega \geq$ *1*$\}$ $+$ *prob* $\{\omega.$ *degree* $\omega <$ *1*$\}$
by (*rule pmf-add*[*OF M-def*], *auto*)
also have ... $\leq$ *?r1* $+$ *?r2*
by (*intro add-mono f prob-degree-lt-1*)
finally show *?thesis* **by** *simp*
qed

private lemma *of-bool-square*: (*of-bool x*)^2 $=$ ((*of-bool x*)::*real*)

by (*cases x*, *auto*)

private definition *Q* **where** $Q\ y\ \omega = card\ \{x \in set\ as.\ int\ (hash\ x\ \omega) < y\}$

private definition *m* **where** $m = card\ (set\ as)$

private lemma
assumes $a \geq 0$
assumes $a \leq int\ p$
shows *exp-Q*: $expectation\ (\lambda\omega.\ real\ (Q\ a\ \omega)) = real\ m * (of\text{-}int\ a)\ /\ p$
and *var-Q*: $variance\ (\lambda\omega.\ real\ (Q\ a\ \omega)) \leq real\ m * (of\text{-}int\ a)\ /\ p$
proof −
have *exp-single*: $expectation\ (\lambda\omega.\ of\text{-}bool\ (int\ (hash\ x\ \omega) < a)) = real\text{-}of\text{-}int\ a\ /real\ p$
if $a{:}x \in set\ as$ **for** *x*
proof −
have *x-le-p*: $x < p$ **using** *a as-lt-p* **by** *simp*
have $expectation\ (\lambda\omega.\ of\text{-}bool\ (int\ (hash\ x\ \omega) < a)) = expectation\ (indicat\text{-}real\ \{\omega.\ int\ (Frequency\text{-}Moment\text{-}0.hash\ p\ x\ \omega) < a\})$
by (*intro arg-cong2*[**where** $f{=}integral^L$] *ext*, *simp-all*)
also have ... $= prob\ \{\omega.\ hash\ x\ \omega \in \{k.\ int\ k < a\}\}$
by (*simp add*:*M-def*)
also have ... $= card\ (\{k.\ int\ k < a\} \cap \{..{<}p\})\ /\ real\ p$
by (*subst prob-range*) (*simp-all add*: *x-le-p ring-of-def mod-ring-def lessThan-def*)
also have ... $= card\ \{..{<}nat\ a\}\ /\ real\ p$
using *assms* **by** (*intro arg-cong2*[**where** $f{=}(/)$] *arg-cong*[**where** $f{=}real$] *arg-cong*[**where** $f{=}card$])
(*auto simp add*:*set-eq-iff*)
also have ... $= real\text{-}of\text{-}int\ a/real\ p$
using *assms* **by** *simp*
finally show $expectation\ (\lambda\omega.\ of\text{-}bool\ (int\ (hash\ x\ \omega) < a)) = real\text{-}of\text{-}int\ a\ /real\ p$
by *simp*
qed

have $expectation(\lambda\omega.\ real\ (Q\ a\ \omega)) = expectation\ (\lambda\omega.\ (\sum x \in set\ as.\ of\text{-}bool\ (int\ (hash\ x\ \omega) < a)))$
by (*simp add*:*Q-def Int-def*)
also have ... $= (\sum x \in set\ as.\ expectation\ (\lambda\omega.\ of\text{-}bool\ (int\ (hash\ x\ \omega) < a)))$
by (*rule Bochner-Integration.integral-sum*, *simp*)
also have ... $= (\sum\ x \in set\ as.\ a\ /real\ p)$
by (*rule sum.cong*, *simp*, *subst exp-single*, *simp*, *simp*)
also have ... $= real\ m *\ real\text{-}of\text{-}int\ a\ /\ real\ p$
by (*simp add*:*m-def*)
finally show $expectation\ (\lambda\omega.\ real\ (Q\ a\ \omega)) = real\ m * real\text{-}of\text{-}int\ a\ /\ p$ **by** *simp*

have *indep*: $J \subseteq set\ as \Longrightarrow card\ J = 2 \Longrightarrow indep\text{-}vars\ (\lambda\text{-}.\ borel)\ (\lambda i\ x.\ of\text{-}bool\ (int\ (hash\ i\ x) < a))\ J$ **for** *J*
using *as-subset-p mod-ring-carr*

by (*intro indep-vars-compose2*[**where** Y=λi x. *of-bool* (*int* x < a) **and** M'=λ-. *discrete*]
k-wise-indep-vars-subset[*OF k-wise-indep*] *finite-subset*[*OF - finite-set*]) *auto*

have *rv*: $\bigwedge x$. $x \in$ *set as* $\Longrightarrow$ *random-variable borel* ($\lambda\omega$. *of-bool* (*int* (*hash* x ω) < a))
by (*simp add*:*M-def*)

have *variance* ($\lambda\omega$. *real* (Q a ω)) = *variance* ($\lambda\omega$. ($\sum x \in$ *set as*. *of-bool* (*int* (*hash* x ω) < a)))
by (*simp add*:*Q-def Int-def*)
also have ... = ($\sum x \in$ *set as*. *variance* ($\lambda\omega$. *of-bool* (*int* (*hash* x ω) < a)))
by (*intro bienaymes-identity-pairwise-indep-2 indep rv*) *auto*
also have ... $\le$ ($\sum$ $x \in$ *set as*. a / *real* p)
by (*rule sum-mono*, *simp add*: *variance-eq of-bool-square*, *simp add*: *exp-single*)
also have ... = *real* m $*$ *real-of-int* a /*real* p
by (*simp add*:*m-def*)
finally show *variance* ($\lambda\omega$. *real* (Q a ω)) $\le$ *real* m $*$ *real-of-int* a / p
by *simp*
qed

private lemma *t-bound*: $t \le$ *81* / (*real-of-rat* δ)^2
proof −
have $t \le$ *80* / (*real-of-rat* δ)^2 + *1* **using** *t-def t-gt-0* **by** *linarith*
also have ... $\le$ *80* / (*real-of-rat* δ)^2 + *1* / (*real-of-rat* δ)^2
using δ*-range* **by** (*intro add-mono*, *simp*, *simp add*:*power-le-one*)
also have ... = *81* / (*real-of-rat* δ)^2 **by** *simp*
finally show *?thesis* **by** *simp*
qed

private lemma *t-r-bound*:
18 $*$ *40* $*$ (*real* t)^2 $*$ *2 powr* ($-$*real* r) $\le$ *1*
proof −
have *720* $*$ (*real* t)^2 $*$ *2 powr* ($-$*real* r) $\le$ *720* $*$ (*81* / (*real-of-rat* δ)^2)^2 $*$ *2 powr* ($-$*4* $*$ *log 2* (*1* / *real-of-rat* δ) $-$ *23*)
using *r-bound t-bound* **by** (*intro mult-left-mono mult-mono power-mono powr-mono*, *auto*)

also have ... $\le$ *720* $*$ (*81* / (*real-of-rat* δ)^2)^2 $*$ (*2 powr* ($-$*4* $*$ *log 2* (*1* / *real-of-rat* δ)) $*$ *2 powr* ($-$*23*))
using δ*-range* **by** (*intro mult-left-mono mult-mono power-mono add-mono*)
(*simp-all add*:*power-le-one powr-diff*)

also have ... = *720* $*$ (*81*^{2} / (*real-of-rat* δ)^*4*) $*$ (*2 powr* (*log 2* ((*real-of-rat* δ)^*4*)) $*$ *2 powr* ($-$*23*))
using δ*-range* **by** (*intro arg-cong2*[**where** f=($*$)])
(*simp-all add*:*power2-eq-square power4-eq-xxxx log-divide log-powr*[*symmetric*])

also have ... = *720* $*$ *81*^{2} $*$ *2 powr* ($-$*23*) **using** δ*-range* **by** *simp*

```
  also have ... ≤ 1 by simp

  finally show ?thesis by simp
qed

private lemma m-eq-F-0: real m = of-rat (F 0 as)
  by (simp add:m-def F-def)

private lemma estimate′-bounds:
  prob {ω. of-rat δ ∗ real-of-rat (F 0 as) < |estimate′ (sketch-rv′ ω) − of-rat (F 0
as)|} ≤ 1/3
proof (cases card (set as) ≥ t)
  case True
  define δ′ where δ′ = 3 ∗ real-of-rat δ / 4
  define u where u = ⌈real t ∗ p / (m ∗ (1+δ′))⌉
  define v where v = ⌊real t ∗ p / (m ∗ (1−δ′))⌋

  define has-no-collision where
    has-no-collision = (λω. ∀x∈ set as. ∀y ∈ set as. (tr-hash x ω = tr-hash y ω
⟶ x = y) ∨ tr-hash x ω > v)

  have 2 powr (−real r) ≤ 2 powr (−(4 ∗ log 2 (1 / real-of-rat δ) + 23))
    using r-bound by (intro powr-mono, linarith, simp)
  also have ... = 2 powr (−4 ∗ log 2 (1 /real-of-rat δ) −23)
    by (rule arg-cong2[where f=(powr)], auto simp add:algebra-simps)
  also have ... ≤ 2 powr ( −1 ∗ log 2 (1 /real-of-rat δ) −4)
    using δ-range by (intro powr-mono diff-mono, auto)
  also have ... = 2 powr ( −1 ∗ log 2 (1 /real-of-rat δ)) /  16
    by (simp add: powr-diff)
  also have ... = real-of-rat δ / 16
    using δ-range by (simp add:log-divide)
  also have ... < real-of-rat δ / 8
    using δ-range by (subst pos-divide-less-eq, auto)
  finally have r-le-δ: 2 powr (−real r) < real-of-rat δ / 8
    by simp

  have δ′-gt-0: δ′ > 0 using δ-range by (simp add:δ′-def)
  have δ′ < 3/4 using δ-range by (simp add:δ′-def)+
  also have ... < 1 by simp
  finally have δ′-lt-1: δ′ < 1 by simp

  have t ≤ 81 / (real-of-rat δ)^2
    using t-bound by simp
  also have ... = (81∗9/16) / (δ′)^2
    by (simp add:δ′-def power2-eq-square)
  also have ... ≤ 46 / δ′^2
    by (intro divide-right-mono, simp, simp)
  finally have t-le-δ′: t ≤ 46/ δ′^2 by simp
```

have $80 \le (\textit{real-of-rat}\ \delta)^2 * (80\ /\ (\textit{real-of-rat}\ \delta)^2)$ **using** *δ-range* **by** *simp*
also have ... $\le (\textit{real-of-rat}\ \delta)^2 * t$
by (*intro mult-left-mono*, *simp add:t-def of-nat-ceiling*, *simp*)
finally have $80 \le (\textit{real-of-rat}\ \delta)^2 * t$ **by** *simp*
hence *t-ge-δ'*: $45 \le t * \delta' * \delta'$ **by** (*simp add:δ'-def power2-eq-square*)

have $m \le \textit{card}\ \{..<n\}$ **unfolding** *m-def* **using** *as-range* **by** (*intro card-mono*, *auto*)
also have ... $\le p$ **using** *n-le-p* **by** *simp*
finally have *m-le-p*: $m \le p$ **by** *simp*

hence *t-le-m*: $t \le \textit{card}\ (\textit{set as})$ **using** *True* **by** *simp*
have *m-ge-0*: $\textit{real}\ m > 0$ **using** *m-def True t-gt-0* **by** *simp*

have $v \le \textit{real}\ t * \textit{real}\ p\ /\ (\textit{real}\ m * (1 - \delta'))$ **by** (*simp add:v-def*)

also have ... $\le \textit{real}\ t * \textit{real}\ p\ /\ (\textit{real}\ m * (1/4))$
using *δ'-lt-1 m-ge-0 δ-range*
by (*intro divide-left-mono mult-left-mono mult-nonneg-nonneg mult-pos-pos*, *simp-all add:δ'-def*)

finally have *v-ubound*: $v \le 4 * \textit{real}\ t * \textit{real}\ p\ /\ \textit{real}\ m$ **by** (*simp add:algebra-simps*)

have *a-ge-1*: $u \ge 1$ **using** *δ'-gt-0 p-gt-0 m-ge-0 t-gt-0*
by (*auto intro*!:*mult-pos-pos divide-pos-pos simp add:u-def*)
hence *a-ge-0*: $u \ge 0$ **by** *simp*
have $\textit{real}\ m * (1 - \delta') < \textit{real}\ m$ **using** *δ'-gt-0 m-ge-0* **by** *simp*
also have ... $\le 1 * \textit{real}\ p$ **using** *m-le-p* **by** *simp*
also have ... $\le \textit{real}\ t * \textit{real}\ p$ **using** *t-gt-0* **by** (*intro mult-right-mono*, *auto*)
finally have $\textit{real}\ m * (1 - \delta') < \textit{real}\ t * \textit{real}\ p$ **by** *simp*
hence *v-gt-0*: $v > 0$ **using** *mult-pos-pos m-ge-0 δ'-lt-1* **by** (*simp add:v-def*)
hence *v-ge-1*: $\textit{real-of-int}\ v \ge 1$ **by** *linarith*

have $\textit{real}\ t \le \textit{real}\ m$ **using** *True m-def* **by** *linarith*
also have ... $< (1 + \delta') * \textit{real}\ m$ **using** *δ'-gt-0 m-ge-0* **by** *force*
finally have *a-le-p-aux*: $\textit{real}\ t < (1 + \delta') * \textit{real}\ m$ **by** *simp*

have $u \le \textit{real}\ t * \textit{real}\ p\ /\ (\textit{real}\ m * (1 + \delta'))+1$ **by** (*simp add:u-def*)
also have ... $< \textit{real}\ p + 1$
using *m-ge-0 δ'-gt-0 a-le-p-aux a-le-p-aux p-gt-0*
by (*simp add*: *pos-divide-less-eq ac-simps*)
finally have $u \le \textit{real}\ p$
by (*metis int-less-real-le not-less of-int-le-iff of-int-of-nat-eq*)
hence *u-le-p*: $u \le \textit{int}\ p$ **by** *linarith*

have $\textit{prob}\ \{\omega.\ Q\ u\ \omega \ge t\} \le \textit{prob}\ \{\omega \in \textit{Sigma-Algebra.space}\ M.\ \textit{abs}\ (\textit{real}\ (Q\ u\ \omega) -$
$\textit{expectation}\ (\lambda\omega.\ \textit{real}\ (Q\ u\ \omega))) \ge 3 * \textit{sqrt}\ (m * \textit{real-of-int}\ u\ /\ p)\}$

proof (*rule pmf-mono*[*OF M-def*])
fix ω
assume $\omega \in \{\omega.\ t \leq Q\ u\ \omega\}$
hence *t-le*: $t \leq Q\ u\ \omega$ **by** *simp*
have *real m* $*$ *real-of-int u* / *real p* $\leq$ *real m* $*$ (*real t* $*$ *real p* / (*real m* $*$ (1 + δ'))+1) / *real p*
using *m-ge-0 p-gt-0* **by** (*intro divide-right-mono mult-left-mono*, *simp-all add*: *u-def*)
also have ... = *real m* $*$ *real t* $*$ *real p* / (*real m* $*$ (1+δ') $*$ *real p*) + *real m* / *real p*
by (*simp add*:*distrib-left add-divide-distrib*)
also have ... = *real t* / (1+δ') + *real m* / *real p*
using *p-gt-0 m-ge-0* **by** *simp*
also have ... $\leq$ *real t* / (1+δ') + 1
using *m-le-p p-gt-0* **by** (*intro add-mono*, *auto*)
finally have *real m* $*$ *real-of-int u* / *real p* $\leq$ *real t* / (1 + δ') + 1
by *simp*

hence 3 $*$ *sqrt* (*real m* $*$ *of-int u* / *real p*) + *real m* $*$ *of-int u* / *real p* $\leq$
3 $*$ *sqrt* (t / (1+δ')+1)+(t/(1+δ')+1)
by (*intro add-mono mult-left-mono real-sqrt-le-mono*, *auto*)
also have ... $\leq$ 3 $*$ *sqrt* (*real t*+1) + ((t $*$ (1 $-$ δ' / (1+δ'))) + 1)
using δ'*-gt-0 t-gt-0* **by** (*intro add-mono mult-left-mono real-sqrt-le-mono*)
(*simp-all add*: *pos-divide-le-eq left-diff-distrib*)
also have ... = 3 $*$ *sqrt* (*real t*+1) + (t $-$ δ' $*$ t / (1+δ')) + 1 **by** (*simp add*:*algebra-simps*)
also have ... $\leq$ 3 $*$ *sqrt* (46 / δ'^2 + 1 / δ'^2) + (t $-$ δ' $*$ $t/2$) + 1 / δ'
using δ'*-gt-0 t-gt-0* δ'*-lt-1 add-pos-pos t-le-*δ'
by (*intro add-mono mult-left-mono real-sqrt-le-mono add-mono*)
(*simp-all add*: *power-le-one pos-le-divide-eq*)
also have ... $\leq$ (21 / δ' + (t $-$ 45 / ($2$$*$$\delta'$))) + 1 / δ'
using δ'*-gt-0 t-ge-*δ' **by** (*intro add-mono*)
(*simp-all add*:*real-sqrt-divide divide-le-cancel real-le-lsqrt pos-divide-le-eq ac-simps*)
also have ... $\leq$ t **using** δ'*-gt-0* **by** *simp*
also have ... $\leq$ $Q\ u\ \omega$ **using** *t-le* **by** *simp*
finally have 3 $*$ *sqrt* (*real m* $*$ *of-int u* / *real p*) + *real m* $*$ *of-int u* / *real p* $\leq$ $Q\ u\ \omega$
by *simp*
hence 3 $*$ *sqrt* (*real m* $*$ *real-of-int u* / *real p*) $\leq$ |*real* ($Q\ u\ \omega$) $-$ *expectation* ($\lambda\omega$. *real* ($Q\ u\ \omega$))|
using *a-ge-0 u-le-p True* **by** (*simp add*:*exp-Q abs-ge-iff*)

thus $\omega \in \{\omega \in$ *Sigma-Algebra.space M*. 3 $*$ *sqrt* (*real m* $*$ *real-of-int u* / *real p*) $\leq$
|*real* ($Q\ u\ \omega$) $-$ *expectation* ($\lambda\omega$. *real* ($Q\ u\ \omega$))|$\}$
by (*simp add*: *M-def*)
qed
also have ... $\leq$ *variance* ($\lambda\omega$. *real* ($Q\ u\ \omega$)) / (3 $*$ *sqrt* (*real m* $*$ *of-int u* / *real*

p))^2
  using *a-ge-1 p-gt-0 m-ge-0*
  by (*intro Chebyshev-inequality*, *simp add*:*M-def*, *auto*)

also have ... $\leq$ (*real m* $*$ *real-of-int u* / *real p*) / (*3* $*$ *sqrt* (*real m* $*$ *of-int u* / *real p*))^2
  using *a-ge-0 u-le-p* **by** (*intro divide-right-mono var-Q*, *auto*)

also have ... $\leq$ *1/9* **using** *a-ge-0* **by** *simp*

finally have *case-1*: *prob* {ω. *Q u* ω $\geq$ *t*} $\leq$ *1/9* **by** *simp*

have *case-2*: *prob* {ω. *Q v* ω < *t*} $\leq$ *1/9*
proof (*cases v* $\leq$ *p*)
  case *True*
  have *prob* {ω. *Q v* ω < *t*} $\leq$ *prob* {ω ∈ *Sigma-Algebra.space M*. *abs* (*real* (*Q v* ω) − *expectation* (λω. *real* (*Q v* ω)))
    $\geq$ *3* $*$ *sqrt* (*m* $*$ *real-of-int v* / *p*)}
  proof (*rule pmf-mono*[*OF M-def*])
    fix ω
    assume ω ∈ *set-pmf* (*pmf-of-set space*)
    have (*real t* + *3* $*$ *sqrt* (*real t* / (*1* − δ′))) $*$ (*1* − δ′) = *real t* − δ′ $*$ *t* + *3* $*$ ((*1*−δ′) $*$ *sqrt*(*real t* / (*1*−δ′)))
      by (*simp add*:*algebra-simps*)

    also have ... = *real t* − δ′ $*$ *t* + *3* $*$ *sqrt* ((*1*−δ′)^2 $*$ (*real t* / (*1*−δ′)))
      using δ′*-lt-1* **by** (*subst real-sqrt-mult*, *simp*)

    also have ... = *real t* − δ′ $*$ *t* + *3* $*$ *sqrt* (*real t* $*$ (*1*− δ′))
      by (*simp add*:*power2-eq-square distrib-left*)

    also have ... $\leq$ *real t* − *45*/ δ′ + *3* $*$ *sqrt* (*real t*)
    using δ′*-gt-0 t-ge-*δ′ δ′*-lt-1* **by** (*intro add-mono mult-left-mono real-sqrt-le-mono*)
      (*simp-all add*:*pos-divide-le-eq ac-simps left-diff-distrib power-le-one*)

    also have ... $\leq$ *real t* − *45*/ δ′ + *3* $*$ *sqrt* (*46* / δ′^{2})
      using *t-le-*δ′ δ′*-lt-1* δ′*-gt-0*
    by (*intro add-mono mult-left-mono real-sqrt-le-mono*, *simp-all add*:*pos-divide-le-eq power-le-one*)

    also have ... = *real t* + (*3* $*$ *sqrt*(*46*) − *45*)/ δ′
      using δ′*-gt-0* **by** (*simp add*:*real-sqrt-divide diff-divide-distrib*)

    also have ... $\leq$ *t*
      using δ′*-gt-0* **by** (*simp add*:*pos-divide-le-eq real-le-lsqrt*)

    finally have *aux*: (*real t* + *3* $*$ *sqrt* (*real t* / (*1* − δ′))) $*$ (*1* − δ′) $\leq$ *real t*
      by *simp*

assume $\omega \in \{\omega.\ Q\ v\ \omega < t\}$
hence $Q\ v\ \omega < t$ **by** *simp*

hence *real* $(Q\ v\ \omega) + 3 * $ *sqrt* (*real m* $*$ *real-of-int v* / *real p*)
$\le$ *real* $t - 1 + 3 *$ *sqrt* (*real m* $*$ *real-of-int v* / *real p*)
using *m-le-p p-gt-0* **by** (*intro add-mono*, *auto simp add*: *algebra-simps add-divide-distrib*)

also have ... $\le$ (*real* $t-1$) $+ 3 *$ *sqrt* (*real m* $*$ (*real t* $*$ *real p* / (*real m* $*$ $(1-\delta'))$) / *real p*)
by (*intro add-mono mult-left-mono real-sqrt-le-mono divide-right-mono*)
(*auto simp add*:*v-def*)

also have ... $\le$ *real* $t + 3 *$ *sqrt*(*real t* / $(1-\delta')) - 1$
using *m-ge-0 p-gt-0* **by** *simp*

also have ... $\le$ *real t* / $(1-\delta')-1$
using δ'*-lt-1 aux* **by** (*simp add*: *pos-le-divide-eq*)
also have ... $\le$ *real m* $*$ (*real t* $*$ *real p* / (*real m* $*$ $(1-\delta')$)) / *real p* $- 1$
using *p-gt-0 m-ge-0* **by** *simp*
also have ... $\le$ *real m* $*$ (*real t* $*$ *real p* / (*real m* $*$ $(1-\delta')$)) / *real p* $-$ *real m* / *real p*
using *m-le-p p-gt-0*
by (*intro diff-mono*, *auto*)
also have ... $=$ *real m* $*$ (*real t* $*$ *real p* / (*real m* $*$ $(1-\delta'))-1$) / *real p*
by (*simp add*: *left-diff-distrib right-diff-distrib diff-divide-distrib*)
also have ... $\le$ *real m* $*$ *real-of-int v* / *real p*
by (*intro divide-right-mono mult-left-mono*, *simp-all add*:*v-def*)

finally have *real* $(Q\ v\ \omega) + 3 *$ *sqrt* (*real m* $*$ *real-of-int v* / *real p*)
$\le$ *real m* $*$ *real-of-int v* / *real p* **by** *simp*

hence $3 *$ *sqrt* (*real m* $*$ *real-of-int v* / *real p*) $\le$ |*real* $(Q\ v\ \omega)$ $-$*expectation* ($\lambda\omega.$ *real* $(Q\ v\ \omega)$)|
using *v-gt-0 True* **by** (*simp add*: *exp-Q abs-ge-iff*)

thus $\omega \in \{\omega\in$ *Sigma-Algebra.space M*. $3 *$ *sqrt* (*real m* $*$ *real-of-int v* / *real p*) $\le$
|*real* $(Q\ v\ \omega)$ $-$ *expectation* ($\lambda\omega.$ *real* $(Q\ v\ \omega)$)|}
by (*simp add*:*M-def*)
qed
also have ... $\le$ *variance* ($\lambda\omega.$ *real* $(Q\ v\ \omega)$) / ($3 *$ *sqrt* (*real m* $*$ *real-of-int v* / *real p*))^2
using *v-gt-0 p-gt-0 m-ge-0*
by (*intro Chebyshev-inequality*, *simp add*:*M-def*, *auto*)

also have ... $\le$ (*real m* $*$ *real-of-int v* / *real p*) / ($3 *$ *sqrt* (*real m* $*$ *real-of-int v* / *real p*))^2
using *v-gt-0 True* **by** (*intro divide-right-mono var-Q*, *auto*)

also have ... = *1/9*
using *p-gt-0 v-gt-0 m-ge-0* **by** (*simp add*:*power2-eq-square*)

finally show *?thesis* **by** *simp*
next
case *False*
have *prob* {ω. *Q v* ω < *t*} ≤ *prob* {ω. *False*}
proof (*rule pmf-mono*[*OF M-def*])
fix ω
assume *a*:ω ∈ {ω. *Q v* ω < *t*}
assume ω ∈ *set-pmf* (*pmf-of-set space*)
hence *b*:⋀*x*. *x* < *p* ⟹ *hash x* ω < *p*
using *hash-range mod-ring-carr* **by** (*simp add*:*M-def measure-pmf-inverse*)
have *t* ≤ *card* (*set as*) **using** *True* **by** *simp*
also have ... ≤ *Q v* ω
unfolding *Q-def* **using** *b False as-lt-p* **by** (*intro card-mono subsetI*, *simp*, *force*)
also have ... < *t* **using** *a* **by** *simp*
finally have *False* **by** *auto*
thus ω ∈ {ω. *False*} **by** *simp*
qed
also have ... = *0* **by** *auto*
finally show *?thesis* **by** *simp*
qed

have *prob* {ω. ¬*has-no-collision* ω} ≤
prob {ω. ∃ *x* ∈ *set as*. ∃ *y* ∈ *set as*. *x* ≠ *y* ∧ *tr-hash x* ω ≤ *real-of-int v* ∧ *tr-hash x* ω = *tr-hash y* ω}
by (*rule pmf-mono*[*OF M-def*]) (*simp add*:*has-no-collision-def M-def*, *force*)

also have ... ≤ (*5/2*) ∗ (*real* (*card* (*set as*)))^2 ∗ (*real-of-int v*)^2 ∗ *2 powr* − *real r* / (*real p*)^2 + *1* / *real p*
using *collision-prob v-ge-1* **by** *blast*

also have ... ≤ (*5/2*) ∗ (*real m*)^2 ∗ (*real-of-int v*)^2 ∗ *2 powr* − *real r* / (*real p*)^2 + *1* / *real p*
by (*intro divide-right-mono add-mono mult-right-mono mult-mono power-mono*, *simp-all add*:*m-def*)

also have ... ≤ (*5/2*) ∗ (*real m*)^2 ∗ (*4* ∗ *real t* ∗ *real p* / *real m*)^2 ∗ (*2 powr* − *real r*) / (*real p*)^2 + *1* / *real p*
using *v-def v-ge-1 v-ubound*
by (*intro add-mono divide-right-mono mult-right-mono mult-left-mono*, *auto*)

also have ... = *40* ∗ (*real t*)^2 ∗ (*2 powr* −*real r*) + *1* / *real p*
using *p-gt-0 m-ge-0 t-gt-0* **by** (*simp add*:*algebra-simps power2-eq-square*)

also have ... ≤ *1/18* + *1/18*

using *t-r-bound p-ge-18* **by** (*intro add-mono*, *simp-all add*: *pos-le-divide-eq*)

also have ... = *1/9* **by** *simp*

finally have *case-3*: *prob* {ω. ¬*has-no-collision* ω} ≤ *1/9* **by** *simp*

have *prob* {ω. *real-of-rat* δ ∗ *of-rat* (*F 0 as*) < |*estimate′* (*sketch-rv′* ω) − *of-rat* (*F 0 as*)|} ≤
prob {ω. *Q u* ω ≥ *t* ∨ *Q v* ω < *t* ∨ ¬(*has-no-collision* ω)}
proof (*rule pmf-mono*[*OF M-def*], *rule ccontr*)
fix ω
assume ω ∈ *set-pmf* (*pmf-of-set space*)
assume ω ∈ {ω. *real-of-rat* δ ∗ *real-of-rat* (*F 0 as*) < |*estimate′* (*sketch-rv′* ω) − *real-of-rat* (*F 0 as*)|}
hence *est*: *real-of-rat* δ ∗ *real-of-rat* (*F 0 as*) < |*estimate′* (*sketch-rv′* ω) − *real-of-rat* (*F 0 as*)| **by** *simp*
assume ω ∉ {ω. *t* ≤ *Q u* ω ∨ *Q v* ω < *t* ∨ ¬ *has-no-collision* ω}
hence ¬(*t* ≤ *Q u* ω ∨ *Q v* ω < *t* ∨ ¬ *has-no-collision* ω) **by** *simp*
hence *lb*: *Q u* ω < *t* **and** *ub*: *Q v* ω ≥ *t* **and** *no-col*: *has-no-collision* ω **by** *simp+*

define *y* **where** *y* = *nth-mset* (*t*−*1*) {#*int* (*hash x* ω). *x* ∈# *mset-set* (*set as*)#}
define *y′* **where** *y′* = *nth-mset* (*t*−*1*) {#*tr-hash x* ω. *x* ∈# *mset-set* (*set as*)#}

have *rank-t-lb*: *u* ≤ *y*
unfolding *y-def* **using** *True t-gt-0 lb*
by (*intro nth-mset-bound-left*, *simp-all add*:*count-less-def swap-filter-image Q-def*)

have *rank-t-ub*: *y* ≤ *v* − *1*
unfolding *y-def* **using** *True t-gt-0 ub*
by (*intro nth-mset-bound-right*, *simp-all add*:*Q-def swap-filter-image count-le-def*)

have *y-ge-0*: *real-of-int y* ≥ *0* **using** *rank-t-lb a-ge-0* **by** *linarith*

have *mono* (λ*x*. *truncate-down r* (*real-of-int x*))
by (*metis truncate-down-mono mono-def of-int-le-iff*)
hence *y′-eq*: *y′* = *truncate-down r y*
unfolding *y-def y′-def* **using** *True t-gt-0*
by (*subst nth-mset-commute-mono*[**where** *f*=(λ*x*. *truncate-down r* (*of-int x*))])
(*simp-all add*: *multiset.map-comp comp-def tr-hash-def*)

have *real-of-int u* ∗ (*1* − *2 powr* −*real r*) ≤ *real-of-int y* ∗ (*1* − *2 powr* (−*real r*))
using *rank-t-lb of-int-le-iff two-pow-r-le-1*
by (*intro mult-right-mono*, *auto*)
also have ... ≤ *y′*

using y'*-eq truncate-down-pos*[*OF y-ge-0*] **by** *simp*
finally have *rank-t-lb'*: $u * (1 - 2\ powr\ -real\ r) \leq y'$ **by** *simp*

have $y' \leq real\text{-}of\text{-}int\ y$
by (*subst* y'*-eq*, *rule truncate-down-le*, *simp*)
also have ... $\leq real\text{-}of\text{-}int\ (v-1)$
using *rank-t-ub of-int-le-iff* **by** *blast*
finally have *rank-t-ub'*: $y' \leq v-1$
by *simp*

have $0 < u * (1-2\ powr\ -real\ r)$
using *a-ge-1 two-pow-r-le-1* **by** (*intro mult-pos-pos*, *auto*)
hence y'*-pos*: $y' > 0$ **using** *rank-t-lb'* **by** *linarith*

have *no-col'*: $\bigwedge x.\ x \leq y' \Longrightarrow count\ \{\#tr\text{-}hash\ x\ \omega.\ x \in\#\ mset\text{-}set\ (set\ as)\#\}\ x \leq 1$
using *rank-t-ub' no-col*
by (*simp add:vimage-def card-le-Suc0-iff-eq count-image-mset has-no-collision-def*) *force*

have *h-1*: $Max\ (sketch\text{-}rv'\ \omega) = y'$
using *True t-gt-0 no-col'*
by (*simp add:sketch-rv'-def y'-def nth-mset-max*)

have $card\ (sketch\text{-}rv'\ \omega) = card\ (least\ ((t-1)+1)\ (set\text{-}mset\ \{\#tr\text{-}hash\ x\ \omega.\ x \in\#\ mset\text{-}set\ (set\ as)\#\}))$
using *t-gt-0* **by** (*simp add:sketch-rv'-def*)
also have ... $= (t-1) + 1$
using *True t-gt-0 no-col'* **by** (*intro nth-mset-max*(2), *simp-all add:y'-def*)
also have ... $= t$ **using** *t-gt-0* **by** *simp*
finally have $card\ (sketch\text{-}rv'\ \omega) = t$ **by** *simp*
hence *h-3*: $estimate'\ (sketch\text{-}rv'\ \omega) = real\ t * real\ p\ /\ y'$
using *h-1* **by** (*simp add:estimate'-def*)

have $(real\ t) * real\ p \leq (1 + \delta') * real\ m * ((real\ t) * real\ p\ /\ (real\ m * (1 + \delta')))$
using δ'*-lt-1 m-def True t-gt-0* δ'*-gt-0* **by** *auto*
also have ... $\leq (1+\delta') * m * u$
using δ'*-gt-0* **by** (*intro mult-left-mono*, *simp-all add:u-def*)
also have ... $< ((1 + real\text{-}of\text{-}rat\ \delta)*(1-real\text{-}of\text{-}rat\ \delta/8)) * m * u$
using *True m-def t-gt-0 a-ge-1* δ*-range*
by (*intro mult-strict-right-mono*, *auto simp add:*δ'*-def right-diff-distrib*)
also have ... $\leq ((1 + real\text{-}of\text{-}rat\ \delta)*(1-2\ powr\ (-r))) * m * u$
using *r-le-*δ δ*-range a-ge-0* **by** (*intro mult-right-mono mult-left-mono*, *auto*)
also have ... $= (1 + real\text{-}of\text{-}rat\ \delta) * m * (u * (1-2\ powr\ -real\ r))$
by *simp*
also have ... $\leq (1 + real\text{-}of\text{-}rat\ \delta) * m * y'$
using δ*-range* **by** (*intro mult-left-mono rank-t-lb'*, *simp*)
finally have $real\ t * real\ p < (1 + real\text{-}of\text{-}rat\ \delta) * m * y'$ **by** *simp*

```
    hence f-1: estimate′ (sketch-rv′ ω) < (1 + real-of-rat δ) ∗ m
      using y′-pos by (simp add: h-3 pos-divide-less-eq)

    have (1 − real-of-rat δ) ∗ m ∗ y′ ≤ (1 − real-of-rat δ) ∗ m ∗ v
      using δ-range rank-t-ub′ y′-pos by (intro mult-mono rank-t-ub′, simp-all)
    also have ... = (1−real-of-rat δ) ∗ (real m ∗ v)
      by simp
    also have ... < (1−δ′) ∗ (real m ∗ v)
      using δ-range m-ge-0 v-ge-1
      by (intro mult-strict-right-mono mult-pos-pos, simp-all add:δ′-def)
    also have ... ≤ (1−δ′) ∗ (real m ∗ (real t ∗ real p / (real m ∗ (1−δ′))))
      using δ′-gt-0 δ′-lt-1 by (intro mult-left-mono, auto simp add:v-def)
    also have ... = real t ∗ real p
      using δ′-gt-0 δ′-lt-1 t-gt-0 p-gt-0 m-ge-0 by auto
    finally have (1 − real-of-rat δ) ∗ m ∗ y′ < real t ∗ real p by simp
    hence f-2: estimate′ (sketch-rv′ ω) > (1 − real-of-rat δ) ∗ m
      using y′-pos by (simp add: h-3 pos-less-divide-eq)

     have abs (estimate′ (sketch-rv′ ω) − real-of-rat (F 0 as)) < real-of-rat δ ∗
(real-of-rat (F 0 as))
      using f-1 f-2 by (simp add:abs-less-iff algebra-simps m-eq-F-0)
    thus False using est by linarith
  qed
  also have ... ≤ 1/9 + (1/9 + 1/9)
    by (intro pmf-add-2[OF M-def] case-1 case-2 case-3)
  also have ... = 1/3 by simp
  finally show ?thesis by simp
next
  case False
  have prob {ω. real-of-rat δ ∗ of-rat (F 0 as) < |estimate′ (sketch-rv′ ω) − of-rat
(F 0 as)|} ≤
    prob {ω. ∃ x ∈ set as. ∃ y ∈ set as. x ≠ y ∧ tr-hash x ω ≤ real p ∧ tr-hash x ω
= tr-hash y ω}
  proof (rule pmf-mono[OF M-def])
    fix ω
     assume a:ω ∈ {ω. real-of-rat δ ∗ real-of-rat (F 0 as) < |estimate′ (sketch-rv′
ω) − real-of-rat (F 0 as)|}
    assume b:ω ∈ set-pmf (pmf-of-set space)
    have c: card (set as) < t using False by auto
    hence card ((λx. tr-hash x ω) ‘ set as) < t
      using card-image-le order-le-less-trans by blast
    hence d:card (sketch-rv′ ω) = card ((λx. tr-hash x ω) ‘ (set as))
      by (simp add:sketch-rv′-def card-least)
    have card (sketch-rv′ ω) < t
      by (metis List.finite-set  c d card-image-le  order-le-less-trans)
   hence estimate′ (sketch-rv′ ω) = card (sketch-rv′ ω) by (simp add:estimate′-def)
    hence card (sketch-rv′ ω) ≠ real-of-rat (F 0 as)
      using a δ-range by simp
        (metis abs-zero cancel-comm-monoid-add-class.diff-cancel of-nat-less-0-iff
```

pos-prod-lt zero-less-of-rat-iff)
hence *card* (*sketch-rv′ ω*) $\neq$ *card* (*set as*)
using *m-def m-eq-F-0* **by** *linarith*
hence $\neg$*inj-on* (λx. *tr-hash x ω*) (*set as*)
using *card-image d* **by** *auto*
moreover have *tr-hash x ω* $\leq$ *real p* **if** *a*:*x* $\in$ *set as* **for** *x*
proof −
have *hash x ω* $<$ *p*
using *hash-range as-lt-p a b* **by** (*simp add*:*mod-ring-carr M-def*)
thus *tr-hash x ω* $\leq$ *real p* **using** *truncate-down-le* **by** (*simp add*:*tr-hash-def*)
qed
ultimately show $\omega \in \{\omega.\ \exists x \in$ *set as*. $\exists y \in$ *set as*. $x \neq y \wedge$ *tr-hash x ω* $\leq$ *real p* $\wedge$ *tr-hash x ω* = *tr-hash y ω*$\}$
by (*simp add*:*inj-on-def*, *blast*)
qed
also have ... $\leq$ (*5/2*) $*$ (*real* (*card* (*set as*)))^2 $*$ (*real p*)^2 $*$ *2 powr* − *real r* / (*real p*)^2 + *1* / *real p*
using *p-gt-0* **by** (*intro collision-prob*, *auto*)
also have ... = (*5/2*) $*$ (*real* (*card* (*set as*)))^2 $*$ *2 powr* (− *real r*) + *1* / *real p*
using *p-gt-0* **by** (*simp add*:*ac-simps power2-eq-square*)
also have ... $\leq$ (*5/2*) $*$ (*real t*)^2 $*$ *2 powr* (−*real r*) + *1* / *real p*
using *False* **by** (*intro add-mono mult-right-mono mult-left-mono power-mono*, *auto*)
also have ... $\leq$ *1/6* + *1/6*
using *t-r-bound p-ge-18* **by** (*intro add-mono*, *simp-all*)
also have ... $\leq$ *1/3* **by** *simp*
finally show *?thesis* **by** *simp*
qed

private lemma *median-bounds*:
$\mathcal{P}$(*ω in measure-pmf* Ω_0. |*median s* (λi. *estimate* (*sketch-rv* (*ω i*))) − *F 0 as*| $\leq$ δ $*$ *F 0 as*) $\geq$ *1* − *real-of-rat ε*
proof −
have *strict-mono-on A real-of-float* **for** *A* **by** (*meson less-float.rep-eq strict-mono-onI*)
hence *real-g-2*: $\bigwedge\omega$. *sketch-rv′ ω* = *real-of-float* ‘ *sketch-rv ω*
by (*simp add*: *sketch-rv′-def sketch-rv-def tr-hash-def least-mono-commute image-comp*)

moreover have *inj-on real-of-float A* **for** *A*
using *real-of-float-inject* **by** (*simp add*:*inj-on-def*)
ultimately have *card-eq*: $\bigwedge\omega$. *card* (*sketch-rv ω*) = *card* (*sketch-rv′ ω*)
using *real-g-2* **by** (*auto intro*!: *card-image*[*symmetric*])

have *Max* (*sketch-rv′ ω*) = *real-of-float* (*Max* (*sketch-rv ω*)) **if** *a*:*card* (*sketch-rv′ ω*) $\geq$ *t* **for** *ω*
proof −
have *mono real-of-float*
using *less-eq-float.rep-eq mono-def* **by** *blast*
moreover have *finite* (*sketch-rv ω*)

by (*simp add*:*sketch-rv-def least-def*)
moreover have *sketch-rv* $\omega \neq \{\}$
using *card-eq*[*symmetric*] *card-gt-0-iff t-gt-0 a* **by** (*simp*, *force*)
ultimately show *?thesis*
by (*subst mono-Max-commute*[**where** *f*=*real-of-float*], *simp-all add*:*real-g-2*)
qed
hence *real-g*: $\bigwedge\omega$. *estimate'* (*sketch-rv'* ω) = *real-of-rat* (*estimate* (*sketch-rv* ω))
by (*simp add*:*estimate-def estimate'-def card-eq of-rat-divide of-rat-mult of-rat-add real-of-rat-of-float*)

have *indep*: *prob-space.indep-vars* (*measure-pmf* Ω_0) (λ-. *borel*) (λi ω. *estimate'* (*sketch-rv'* (ω *i*))) {*0..<s*}
unfolding Ω_0-*def*
by (*rule indep-vars-restrict-intro'*, *auto simp add*:*restrict-dfl-def lessThan-atLeast0*)

moreover have − (*18* ∗ *ln* (*real-of-rat* ε)) ≤ *real s*
using *of-nat-ceiling* **by** (*simp add*:*s-def*) *blast*

moreover have *i* < *s* ⟹ *measure* Ω_0 {ω. *of-rat* δ ∗ *of-rat* (*F 0 as*) < |*estimate'* (*sketch-rv'* (ω *i*)) − *of-rat* (*F 0 as*)|} ≤ *1/3*
for *i*
using *estimate'-bounds* **unfolding** Ω_0-*def M-def*
by (*subst prob-prod-pmf-slice*, *simp-all*)

ultimately have *1*−*real-of-rat* ε ≤ $\mathcal{P}$(ω *in measure-pmf* Ω_0.
|*median s* (λi. *estimate'* (*sketch-rv'* (ω *i*))) − *real-of-rat* (*F 0 as*)| ≤ *real-of-rat* δ ∗ *real-of-rat* (*F 0 as*))
using ε-*range prob-space-measure-pmf*
by (*intro prob-space.median-bound-2*) *auto*
also have ... = $\mathcal{P}$(ω *in measure-pmf* Ω_0.
|*median s* (λi. *estimate* (*sketch-rv* (ω *i*))) − *F 0 as*| ≤ δ ∗ *F 0 as*)
using *s-gt-0 median-rat*[*symmetric*] *real-g* **by** (*intro arg-cong2*[**where** *f*=*measure*])
(*simp-all add*:*of-rat-diff*[*symmetric*] *of-rat-mult*[*symmetric*] *of-rat-less-eq*)
finally show $\mathcal{P}$(ω *in measure-pmf* Ω_0. |*median s* (λi. *estimate* (*sketch-rv* (ω *i*))) − *F 0 as*| ≤ δ ∗ *F 0 as*) ≥ *1*−*real-of-rat* ε
by *blast*
qed

lemma *f0-alg-correct'*:
$\mathcal{P}$(ω *in measure-pmf result*. |ω − *F 0 as*| ≤ δ ∗ *F 0 as*) ≥ *1* − *of-rat* ε
proof −
have *f0-result-elim*: $\bigwedge x$. *f0-result* (*s*, *t*, *p*, *r*, *x*, λi∈{..<*s*}. *sketch-rv* (*x i*)) =
return-pmf (*median s* (λi. *estimate* (*sketch-rv* (*x i*))))
by (*simp add*:*estimate-def*, *rule median-cong*, *simp*)

have *result* = *map-pmf* (λx. (*s*, *t*, *p*, *r*, *x*, λi∈{..<*s*}. *sketch-rv* (*x i*))) Ω_0 ⤜ *f0-result*
by (*subst result-def*, *subst f0-alg-sketch*, *simp*)
also have ... = Ω_0 ⤜ (λx. *return-pmf* (*s*, *t*, *p*, *r*, *x*, λi∈{..<*s*}. *sketch-rv* (*x i*)))

$\gg=$ *f0-result*
  by (*simp add*:*t-def p-def r-def s-def map-pmf-def*)
 also have ... $= \Omega_0 \gg= (\lambda x.$ *return-pmf* (*median s* ($\lambda i.$ *estimate* (*sketch-rv* (*x i*)))))
  by (*subst bind-assoc-pmf*, *subst bind-return-pmf*, *subst f0-result-elim*) *simp*
 finally have *a*:*result* $= \Omega_0 \gg= (\lambda x.$ *return-pmf* (*median s* ($\lambda i.$ *estimate* (*sketch-rv* (*x i*)))))
  by *simp*

 show *?thesis*
  using *median-bounds* **by** (*simp add*: *a map-pmf-def*[*symmetric*])
qed

private lemma *f-subset*:
 assumes g ‘ $A \subseteq h$ ‘ B
 shows $(\lambda x.\ f\ (g\ x))$ ‘ $A \subseteq (\lambda x.\ f\ (h\ x))$ ‘ B
 using *assms* **by** *auto*

lemma *f0-exact-space-usage′*:
 defines $\Omega \equiv$ *fold* (λa *state*. *state* $\gg=$ *f0-update a*) *as* (*f0-init* δ ε *n*)
 shows *AE* ω *in* Ω. *bit-count* (*encode-f0-state* ω) $\leq$ *f0-space-usage* (n, ε, δ)
proof −

 have *log-2-4*: *log 2 4* = *2*
  by (*metis log2-of-power-eq mult-2 numeral-Bit0 of-nat-numeral power2-eq-square*)

 have *a*: *bit-count* (F_e (*float-of* (*truncate-down r y*))) $\leq$
  ereal (*12* + *4* $*$ *real r* + *2* $*$ *log 2* (*log 2* (*n*+*13*))) **if** *a-1*:$y \in \{..<p\}$ **for** y
 proof (*cases* $y \geq 1$)
  case *True*

  have *aux-1*: $0 < 2 +$ *log 2* (*real y*)
   using *True* **by** (*intro add-pos-nonneg*, *auto*)
  have *aux-2*: $0 < 2 +$ *log 2* (*real p*)
   using *p-gt-1* **by** (*intro add-pos-nonneg*, *auto*)

  have *bit-count* (F_e (*float-of* (*truncate-down r y*))) $\leq$
   ereal (*10* + *4* $*$ *real r* + *2* $*$ *log 2* (*2* + |*log 2* |*real y*||))
   by (*rule truncate-float-bit-count*)
  also have ... = *ereal* (*10* + *4* $*$ *real r* + *2* $*$ *log 2* (*2* + (*log 2* (*real y*))))
   using *True* **by** *simp*
  also have ... $\leq$ *ereal* (*10* + *4* $*$ *real r* + *2* $*$ *log 2* (*2* + *log 2 p*))
   using *aux-1 aux-2 True p-gt-0 a-1* **by** *simp*
  also have ... $\leq$ *ereal* (*10* + *4* $*$ *real r* + *2* $*$ *log 2* (*log 2 4* + *log 2* (*2* $*$ *n* + *40*)))
   using *log-2-4 p-le-n p-gt-0*
   by (*simp add*: *Transcendental.log-mono aux-2*)
  also have ... = *ereal* (*10* + *4* $*$ *real r* + *2* $*$ *log 2* (*log 2* (*8* $*$ *n* + *160*)))
   by (*simp flip*: *log-mult-pos*)

also have ... ≤ *ereal* (*10* + *4* ∗ *real r* + *2* ∗ *log 2* (*log 2* ((*n*+*13*) *powr 2*)))
by (*intro ereal-mono add-mono mult-left-mono Transcendental.log-mono of-nat-mono add-pos-nonneg*)
(*auto simp add*:*power2-eq-square algebra-simps*)
also have ... = *ereal* (*10* + *4* ∗ *real r* + *2* ∗ *log 2* (*log 2 4* ∗ *log 2* (*n* + *13*)))
using *log-2-4 log-powr* **by** *presburger*
also have ... = *ereal* (*12* + *4* ∗ *real r* + *2* ∗ *log 2* (*log 2* (*n* + *13*)))
by (*simp add*: *log-mult-pos log-2-4*)
finally show *?thesis* **by** *simp*
next
case *False*
hence $y = 0$ **using** *a-1* **by** *simp*
then show *?thesis* **by** (*simp add*:*float-bit-count-zero*)
qed

have *bit-count* (*encode-f0-state* (*s*, *t*, *p*, *r*, *x*, λ*i*∈{..<*s*}. *sketch-rv* (*x i*))) ≤
f0-space-usage (*n*, ε, δ) **if** *b*: *x* ∈ {..<*s*} $\to_E$ *space* **for** *x*
proof −
have *c*: *x* ∈ *extensional* {..<*s*} **using** *b* **by** (*simp add*:*PiE-def*)

have *d*: *sketch-rv* (*x y*) ⊆ (λ*k*. *float-of* (*truncate-down r k*)) ‘ {..<*p*}
if *d-1*: *y* < *s* **for** *y*
proof −
have *sketch-rv* (*x y*) ⊆ (λ*xa*. *float-of* (*truncate-down r* (*hash xa* (*x y*)))) ‘ *set as*
using *least-subset* **by** (*auto simp add*:*sketch-rv-def tr-hash-def*)
also have ... ⊆ (λ*k*. *float-of* (*truncate-down r* (*real k*))) ‘ {..<*p*}
using *b hash-range as-lt-p d-1*
by (*intro f-subset*[**where** *f*=λ*x*. *float-of* (*truncate-down r* (*real x*))] *image-subsetI*)
(*simp add*: *PiE-iff mod-ring-carr*)
finally show *?thesis*
by *simp*
qed

have ⋀*y*. *y* < *s* ⟹ *finite* (*sketch-rv* (*x y*))
unfolding *sketch-rv-def* **by** (*rule finite-subset*[*OF least-subset*], *simp*)
moreover have *card-sketch*: ⋀*y*. *y* < *s* ⟹ *card* (*sketch-rv* (*x y*)) ≤ *t*
by (*simp add*:*sketch-rv-def card-least*)
moreover have ⋀*y z*. *y* < *s* ⟹ *z* ∈ *sketch-rv* (*x y*) ⟹
bit-count (F_e *z*) ≤ *ereal* (*12* + *4* ∗ *real r* + *2* ∗ *log 2* (*log 2* (*real n* + *13*)))
using *a d* **by** *auto*
ultimately have *e*: ⋀*y*. *y* < *s* ⟹ *bit-count* (S_e F_e (*sketch-rv* (*x y*)))
≤ *ereal* (*real t*) ∗ (*ereal* (*12* + *4* ∗ *real r* + *2* ∗ *log 2* (*log 2* (*real* (*n* + *13*)))) + *1*) + *1*
using *float-encoding* **by** (*intro set-bit-count-est*, *auto*)

have *f*: ⋀*y*. *y* < *s* ⟹ *bit-count* (P_e *p 2* (*x y*)) ≤ *ereal* (*real 2* ∗ (*log 2* (*real p*) + *1*))

using *p-gt-1 b*
by (*intro bounded-degree-polynomial-bit-count*) (*simp-all add:space-def PiE-def Pi-def*)

have *bit-count* (*encode-f0-state* (*s*, *t*, *p*, *r*, *x*, $\lambda i \in \{..<s\}$. *sketch-rv* (*x i*))) =
bit-count (N_e *s*) + *bit-count* (N_e *t*) + *bit-count* (N_e *p*) + *bit-count* (N_e *r*) +
bit-count (([*0*..<*s*] $\rightarrow_e$ P_e *p 2*) *x*) +
bit-count (([*0*..<*s*] $\rightarrow_e$ S_e F_e) ($\lambda i \in \{..<s\}$. *sketch-rv* (*x i*)))
by (*simp add:encode-f0-state-def dependent-bit-count lessThan-atLeast0 s-def*[*symmetric*] *t-def*[*symmetric*] *p-def*[*symmetric*] *r-def*[*symmetric*] *ac-simps*)
also have ... $\leq$ *ereal* (*2*∗ *log 2* (*real s* + *1*) + *1*) + *ereal* (*2*∗ *log 2* (*real t* + *1*) + *1*)
+ *ereal* (*2*∗ *log 2* (*real p* + *1*) + *1*) + *ereal* (*2* ∗ *log 2* (*real r* + *1*) + *1*)
+ (*ereal* (*real s*) ∗ (*ereal* (*real 2* ∗ (*log 2* (*real p*) + *1*))))
+ (*ereal* (*real s*) ∗ ((*ereal* (*real t*) ∗
(*ereal* (*12* + *4* ∗ *real r* + *2* ∗ *log 2* (*log 2* (*real* (*n* + *13*)))) + *1*) + *1*)))
using *c e f*
by (*intro add-mono exp-golomb-bit-count fun-bit-count-est*[**where** *xs*=[*0*..<*s*], *simplified*])
(*simp-all add:lessThan-atLeast0*)
also have ... = *ereal* (*4* + *2* ∗ *log 2* (*real s* + *1*) + *2* ∗ *log 2* (*real t* + *1*) +
2 ∗ *log 2* (*real p* + *1*) + *2* ∗ *log 2* (*real r* + *1*) + *real s* ∗ (*3* + *2* ∗ *log 2* (*real p*) +
real t ∗ (*13* + (*4* ∗ *real r* + *2* ∗ *log 2* (*log 2* (*real n* + *13*))))))
by (*simp add:algebra-simps*)
also have ... $\leq$ *ereal* (*4* + *2* ∗ *log 2* (*real s* + *1*) + *2* ∗ *log 2* (*real t* + *1*) +
2 ∗ *log 2* (*2* ∗ (*21* + *real n*)) + *2* ∗ *log 2* (*real r* + *1*) + *real s* ∗ (*3* + *2* ∗ *log 2* (*2* ∗ (*21* + *real n*)) +
real t ∗ (*13* + (*4* ∗ *real r* + *2* ∗ *log 2* (*log 2* (*real n* + *13*))))))
using *p-le-n p-gt-0*
by (*intro ereal-mono add-mono mult-left-mono*, *auto*)
also have ... = *ereal* (*6* + *2* ∗ *log 2* (*real s* + *1*) + *2* ∗ *log 2* (*real t* + *1*) +
2 ∗ *log 2* (*21* + *real n*) + *2* ∗ *log 2* (*real r* + *1*) + *real s* ∗ (*5* + *2* ∗ *log 2* (*21* + *real n*) +
real t ∗ (*13* + (*4* ∗ *real r* + *2* ∗ *log 2* (*log 2* (*real n* + *13*))))))
by (*subst* (*1 2*) *log-mult*, *auto*)
also have ... $\leq$ *f0-space-usage* (*n*, ε, δ)
by (*simp add:s-def*[*symmetric*] *r-def*[*symmetric*] *t-def*[*symmetric*] *Let-def*)
(*simp add:algebra-simps*)
finally show *bit-count* (*encode-f0-state* (*s*, *t*, *p*, *r*, *x*, $\lambda i \in \{..<s\}$. *sketch-rv* (*x i*))) $\leq$
f0-space-usage (*n*, ε, δ) **by** *simp*
qed
hence $\bigwedge x.\ x \in$ *set-pmf* $\Omega_0 \Longrightarrow$
bit-count (*encode-f0-state* (*s*, *t*, *p*, *r*, *x*, $\lambda i \in \{..<s\}$. *sketch-rv* (*x i*))) $\leq$ *ereal* (*f0-space-usage* (*n*, ε, δ))
by (*simp add*:Ω_0*-def set-prod-pmf del:f0-space-usage.simps*)
hence $\bigwedge y.\ y \in$ *set-pmf* $\Omega \Longrightarrow$ *bit-count* (*encode-f0-state y*) $\leq$ *ereal* (*f0-space-usage* (*n*, ε, δ))

by (*simp add*: Ω*-def f0-alg-sketch del*:*f0-space-usage.simps f0-init.simps*)
(*metis* (*no-types*, *lifting*) *image-iff pmf.set-map*)
thus *?thesis*
by (*simp add*: *AE-measure-pmf-iff del*:*f0-space-usage.simps*)
qed

end

Main results of this section:

theorem *f0-alg-correct*:
assumes $\varepsilon \in \{0<..<1\}$
assumes $\delta \in \{0<..<1\}$
assumes *set as* $\subseteq \{..<n\}$
defines $\Omega \equiv$ *fold* (λa *state. state* $\gg=$ *f0-update a*) *as* (*f0-init* δ ε *n*) $\gg=$ *f0-result*
shows $\mathcal{P}$(ω *in measure-pmf* Ω. $|\omega - F\ 0\ as| \leq \delta * F\ 0\ as) \geq 1 -$ *of-rat* ε
using *f0-alg-correct′*[*OF assms*(*1*−*3*)] **unfolding** Ω*-def* **by** *blast*

theorem *f0-exact-space-usage*:
assumes $\varepsilon \in \{0<..<1\}$
assumes $\delta \in \{0<..<1\}$
assumes *set as* $\subseteq \{..<n\}$
defines $\Omega \equiv$ *fold* (λa *state. state* $\gg=$ *f0-update a*) *as* (*f0-init* δ ε *n*)
shows *AE* ω *in* Ω. *bit-count* (*encode-f0-state* ω) $\leq$ *f0-space-usage* (*n*, ε, δ)
using *f0-exact-space-usage′*[*OF assms*(*1*−*3*)] **unfolding** Ω*-def* **by** *blast*

theorem *f0-asymptotic-space-complexity*:
f0-space-usage $\in$ *O*[*at-top* $\times_F$ *at-right 0* $\times_F$ *at-right 0*](λ(*n*, ε, δ). *ln* (*1* / *of-rat* ε) $*$
(*ln* (*real n*) + *1* / (*of-rat* $\delta)^2 *$ (*ln* (*ln* (*real n*)) + *ln* (*1* / *of-rat* δ))))
(**is** - $\in$ *O*[*?F*](*?rhs*))
proof −
define *n-of* :: *nat* × *rat* × *rat* ⇒ *nat* **where** *n-of* = (λ(*n*, ε, δ). *n*)
define ε*-of* :: *nat* × *rat* × *rat* ⇒ *rat* **where** ε*-of* = (λ(*n*, ε, δ). ε)
define δ*-of* :: *nat* × *rat* × *rat* ⇒ *rat* **where** δ*-of* = (λ(*n*, ε, δ). δ)
define *t-of* **where** *t-of* = (λx. *nat* $\lceil 80$ / (*real-of-rat* (δ*-of x*))$^2\rceil$)
define *s-of* **where** *s-of* = (λx. *nat* $\lceil -(18 *$ *ln* (*real-of-rat* (ε*-of x*)))$\rceil$)
define *r-of* **where** *r-of* = (λx. *nat* (*4* $*$ $\lceil$*log 2* (*1* / *real-of-rat* (δ*-of x*))$\rceil$ + *23*))

define *g* **where** *g* = (λx. *ln* (*1* / *of-rat* (ε*-of x*)) $*$ (*ln* (*real* (*n-of x*)) +
1 / (*of-rat* (δ*-of x*))^2 $*$ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *of-rat* (δ*-of x*)))))

have *evt*: ($\bigwedge x$.
0 < *real-of-rat* (δ*-of x*) $\wedge$ *0* < *real-of-rat* (ε*-of x*) $\wedge$
1/*real-of-rat* (δ*-of x*) $\geq \delta$ $\wedge$ *1*/*real-of-rat* (ε*-of x*) $\geq \varepsilon$ $\wedge$
real (*n-of x*) $\geq n \Longrightarrow P\ x) \Longrightarrow$ *eventually P ?F* (**is** ($\bigwedge x$. *?prem x* $\Longrightarrow$ -) $\Longrightarrow$ -)
for δ ε *n P*
apply (*rule eventually-mono*[**where** *P*=*?prem* **and** *Q*=*P*])
apply (*simp add*:ε*-of-def case-prod-beta′* δ*-of-def n-of-def*)

apply (*intro eventually-conj eventually-prod1' eventually-prod2' sequentially-inf eventually-at-right-less inv-at-right-0-inf*)
by (*auto simp add*:*prod-filter-eq-bot*)

have *exp-pos*: *exp k* $\leq$ *real x* $\Longrightarrow$ *x* $>$ *0* **for** *k x*
using *exp-gt-zero gr0I* **by** *force*

have *exp-gt-1*: *exp 1* $\geq$ (*1*::*real*)
by *simp*

have *1*: (λ-. *1*) $\in$ *O*[*?F*](λ*x*. *ln* (*1* / *real-of-rat* (ε*-of x*)))
by (*auto intro*!:*landau-o.big-mono evt*[**where** ε=*exp 1*] *iffD2*[*OF ln-ge-iff*] *simp add*:*abs-ge-iff*)

have *2*: (λ-. *1*) $\in$ *O*[*?F*](λ*x*. *ln* (*1* / *real-of-rat* (δ*-of x*)))
by (*auto intro*!:*landau-o.big-mono evt*[**where** δ=*exp 1*] *iffD2*[*OF ln-ge-iff*] *simp add*:*abs-ge-iff*)

have *3*: (λ*x*. *1*) $\in$ *O*[*?F*](λ*x*. *ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (δ*-of x*)))
using *exp-pos*
by (*intro landau-sum-2 2 evt*[**where** *n*=*exp 1* **and** δ=*1*] *ln-ge-zero iffD2*[*OF ln-ge-iff*], *auto*)
have *4*: (λ-. *1*) $\in$ *O*[*?F*](λ*x*. *1* / (*real-of-rat* (δ*-of x*))^2)
using *one-le-power*
by (*intro landau-o.big-mono evt*[**where** δ=*1*], *auto simp add*:*power-one-over*[*symmetric*])

have (λ*x*. *80* $*$ (*1* / (*real-of-rat* (δ*-of x*))^2)) $\in$ *O*[*?F*](λ*x*. *1* / (*real-of-rat* (δ*-of x*))^2)
by (*subst landau-o.big.cmult-in-iff*, *auto*)
hence *5*: (λ*x*. *real* (*t-of x*)) $\in$ *O*[*?F*](λ*x*. *1* / (*real-of-rat* (δ*-of x*))^2)
unfolding *t-of-def*
by (*intro landau-real-nat landau-ceil 4*, *auto*)

have (λ*x*. *ln* (*real-of-rat* (ε*-of x*))) $\in$ *O*[*?F*](λ*x*. *ln* (*1* / *real-of-rat* (ε*-of x*)))
by (*intro landau-o.big-mono evt*[**where** ε=*1*], *auto simp add*:*ln-div*)
hence *6*: (λ*x*. *real* (*s-of x*)) $\in$ *O*[*?F*](λ*x*. *ln* (*1* / *real-of-rat* (ε*-of x*)))
unfolding *s-of-def* **by** (*intro landau-nat-ceil 1*, *simp*)

have *7*: (λ*x*. *1*) $\in$ *O*[*?F*](λ*x*. *ln* (*real* (*n-of x*)))
using *exp-pos* **by** (*auto intro*!: *landau-o.big-mono evt*[**where** *n*=*exp 1*] *iffD2*[*OF ln-ge-iff*] *simp*: *abs-ge-iff*)

have *8*: (λ-. *1*) $\in$
O[*?F*](λ*x*. *ln* (*real* (*n-of x*)) + *1* / (*real-of-rat* (δ*-of x*))^2 $*$ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (δ*-of x*))))
using *order-trans*[*OF exp-gt-1*] *exp-pos*
by(*intro landau-sum-1 7 evt*[**where** *n*=*exp 1* **and** δ=*1*] *ln-ge-zero iffD2*[*OF ln-ge-iff*]

mult-nonneg-nonneg add-nonneg-nonneg; *force*)

have (λ*x*. *ln* (*real* (*s-of x*) + *1*)) ∈ *O*[*?F*](λ*x*. *ln* (*1* / *real-of-rat* (ε*-of x*)))
by (*intro landau-ln-3 sum-in-bigo 6 1*, *simp*)

hence *9*: (λ*x*. *log 2* (*real* (*s-of x*) + *1*)) ∈ *O*[*?F*](*g*)
unfolding *g-def* **by** (*intro landau-o.big-mult-1 8*, *auto simp*:*log-def*)
have *10*: (λ*x*. *1*) ∈ *O*[*?F*](*g*)
unfolding *g-def* **by** (*intro landau-o.big-mult-1 8 1*)

have (λ*x*. *ln* (*real* (*t-of x*) + *1*)) ∈
O[*?F*](λ*x*. *1* / (*real-of-rat* (δ*-of x*))^2 ∗ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (δ*-of x*))))
using *5* **by** (*intro landau-o.big-mult-1 3 landau-ln-3 sum-in-bigo 4*, *simp-all*)
hence (λ*x*. *log 2* (*real* (*t-of x*) + *1*)) ∈
O[*?F*](λ*x*. *ln* (*real* (*n-of x*)) + *1* / (*real-of-rat* (δ*-of x*))^2 ∗ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (δ*-of x*))))
using *order-trans*[*OF exp-gt-1*] *exp-pos*
by (*intro landau-sum-2 evt*[**where** *n*=*exp 1* **and** δ=*1*] *ln-ge-zero iffD2*[*OF ln-ge-iff*]
mult-nonneg-nonneg add-nonneg-nonneg; *force simp add*:*log-def*)
hence *11*: (λ*x*. *log 2* (*real* (*t-of x*) + *1*)) ∈ *O*[*?F*](*g*)
unfolding *g-def* **by** (*intro landau-o.big-mult-1′ 1*, *auto*)
have (λ*x*. *1*) ∈ *O*[*?F*](λ*x*. *real* (*n-of x*))
by (*intro landau-o.big-mono evt*[**where** *n*=*1*], *auto*)
hence (λ*x*. *ln* (*real* (*n-of x*) + *21*)) ∈ *O*[*?F*](λ*x*. *ln* (*real* (*n-of x*)))
by (*intro landau-ln-2*[**where** *a*=*2*] *evt*[**where** *n*=*2*] *sum-in-bigo*, *auto*)

hence *12*: (λ*x*. *log 2* (*real* (*n-of x*) + *21*)) ∈ *O*[*?F*](*g*)
unfolding *g-def* **using** *exp-pos order-trans*[*OF exp-gt-1*]
by (*intro landau-o.big-mult-1′ 1 landau-sum-1 evt*[**where** *n*=*exp 1* **and** δ=*1*]
ln-ge-zero iffD2[*OF ln-ge-iff*] *mult-nonneg-nonneg add-nonneg-nonneg*; *force simp add*:*log-def*)

have (λ*x*. *ln* (*1* / *real-of-rat* (δ*-of x*))) ∈ *O*[*?F*](λ*x*. *1* / (*real-of-rat* (δ*-of x*))^2)
by (*intro landau-ln-3 evt*[**where** δ=*1*] *landau-o.big-mono*)
(*auto simp add*:*power-one-over*[*symmetric*] *self-le-power*)
hence (λ*x*. *real* (*nat* (*4*∗⌈*log 2* (*1* / *real-of-rat* (δ*-of x*))⌉+*23*))) ∈ *O*[*?F*](λ*x*. *1* / (*real-of-rat* (δ*-of x*))^2)
using *4* **by** (*auto intro*!: *landau-real-nat sum-in-bigo landau-ceil simp*:*log-def*)
hence (λ*x*. *ln* (*real* (*r-of x*) + *1*)) ∈ *O*[*?F*](λ*x*. *1* / (*real-of-rat* (δ*-of x*))^2)
unfolding *r-of-def*
by (*intro landau-ln-3 sum-in-bigo 4*, *auto*)
hence (λ*x*. *log 2* (*real* (*r-of x*) + *1*)) ∈
O[*?F*](λ*x*. (*1* / (*real-of-rat* (δ*-of x*))^2) ∗ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (δ*-of x*))))
by (*intro landau-o.big-mult-1 3*, *simp add*:*log-def*)
hence (λ*x*. *log 2* (*real* (*r-of x*) + *1*)) ∈
O[*?F*](λ*x*. *ln* (*real* (*n-of x*)) + *1* / (*real-of-rat* (δ*-of x*))^2 ∗ (*ln* (*ln* (*real* (*n-of*

x))) + *ln* (*1* / *real-of-rat* (*δ-of x*))))
  using *exp-pos order-trans*[*OF exp-gt-1*]
  by (*intro landau-sum-2 evt*[**where** *n*=*exp 1* **and** *δ*=*1*] *ln-ge-zero*
    iffD2[*OF ln-ge-iff*] *add-nonneg-nonneg mult-nonneg-nonneg*; *force*)
 hence *13*: (*λx*. *log 2* (*real* (*r-of x*) + *1*)) ∈ *O*[*?F*](*g*)
  unfolding *g-def* **by** (*intro landau-o.big-mult-1′ 1*, *auto*)
 have *14*: (*λx*. *1*) ∈ *O*[*?F*](*λx*. *real* (*n-of x*))
  by (*intro landau-o.big-mono evt*[**where** *n*=*1*], *auto*)

 have (*λx*. *ln* (*real* (*n-of x*) + *13*)) ∈ *O*[*?F*](*λx*. *ln* (*real* (*n-of x*)))
  using *14* **by** (*intro landau-ln-2*[**where** *a*=*2*] *evt*[**where** *n*=*2*] *sum-in-bigo*, *auto*)

 hence (*λx*. *ln* (*log 2* (*real* (*n-of x*) + *13*))) ∈ *O*[*?F*](*λx*. *ln* (*ln* (*real* (*n-of x*))))
  using *exp-pos* **by** (*intro landau-ln-2*[**where** *a*=*2*] *iffD2*[*OF ln-ge-iff*] *evt*[**where** *n*=*exp 2*])
   (*auto simp add*:*log-def*)

 hence (*λx*. *log 2* (*log 2* (*real* (*n-of x*) + *13*))) ∈ *O*[*?F*](*λx*. *ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (*δ-of x*)))
  using *exp-pos* **by** (*intro landau-sum-1 evt*[**where** *n*=*exp 1* **and** *δ*=*1*] *ln-ge-zero* *iffD2*[*OF ln-ge-iff*])
   (*auto simp add*:*log-def*)

 moreover have (*λx*. *real* (*r-of x*)) ∈ *O*[*?F*](*λx*. *ln* (*1* / *real-of-rat* (*δ-of x*)))
  unfolding *r-of-def* **using** *2*
  by (*auto intro*!: *landau-real-nat sum-in-bigo landau-ceil simp*:*log-def*)
 hence (*λx*. *real* (*r-of x*)) ∈ *O*[*?F*](*λx*. *ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (*δ-of x*)))
  using *exp-pos*
  by (*intro landau-sum-2 evt*[**where** *n*=*exp 1* **and** *δ*=*1*] *ln-ge-zero* *iffD2*[*OF ln-ge-iff*], *auto*)

 ultimately have *15*: (*λx*. *real* (*t-of x*) ∗ (*13* + *4* ∗ *real* (*r-of x*) + *2* ∗ *log 2* (*log 2* (*real* (*n-of x*) + *13*))))
   ∈ *O*[*?F*](*λx*. *1* / (*real-of-rat* (*δ-of x*))^2 ∗ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (*δ-of x*))))
  using *5 3*
  by (*intro landau-o.mult sum-in-bigo*, *auto*)

 have (*λx*. *5* + *2* ∗ *log 2* (*21* + *real* (*n-of x*)) + *real* (*t-of x*) ∗ (*13* + *4* ∗ *real* (*r-of x*) + *2* ∗ *log 2* (*log 2* (*real* (*n-of x*) + *13*))))
   ∈ *O*[*?F*](*λx*. *ln* (*real* (*n-of x*)) + *1* / (*real-of-rat* (*δ-of x*))^2 ∗ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (*δ-of x*))))
 proof −
  have ∀$_F$ *x in ?F*. *0* ≤ *ln* (*real* (*n-of x*))
   by (*intro evt*[**where** *n*=*1*] *ln-ge-zero*, *auto*)
  moreover have ∀$_F$ *x in ?F*. *0* ≤ *1* / (*real-of-rat* (*δ-of x*))^2 ∗ (*ln* (*ln* (*real* (*n-of x*))) + *ln* (*1* / *real-of-rat* (*δ-of x*)))

using *exp-pos*
by (*intro evt*[**where** *n=exp 1* **and** *δ=1*] *mult-nonneg-nonneg add-nonneg-nonneg ln-ge-zero iffD2*[*OF ln-ge-iff*]) *auto*
moreover have (*λx. ln* (*21* + *real* (*n-of x*))) ∈ *O*[*?F*](*λx. ln* (*real* (*n-of x*)))
using *14* **by** (*intro landau-ln-2*[**where** *a=2*] *sum-in-bigo evt*[**where** *n=2*], *auto*)
hence (*λx. 5* + *2* ∗ *log 2* (*21* + *real* (*n-of x*))) ∈ *O*[*?F*](*λx. ln* (*real* (*n-of x*)))
using *7* **by** (*intro sum-in-bigo*, *auto simp add*:*log-def*)
ultimately show *?thesis*
using *15* **by** (*rule landau-sum*)
qed

hence *16*: (*λx. real* (*s-of x*) ∗ (*5* + *2* ∗ *log 2* (*21* + *real* (*n-of x*)) + *real* (*t-of x*) ∗
(*13* + *4* ∗ *real* (*r-of x*) + *2* ∗ *log 2* (*log 2* (*real* (*n-of x*) + *13*))))) ∈ *O*[*?F*](*g*)
unfolding *g-def*
by (*intro landau-o.mult 6*, *auto*)

have *f0-space-usage* = (*λx. f0-space-usage* (*n-of x*, *ε-of x*, *δ-of x*))
by (*simp add*:*case-prod-beta′ n-of-def ε-of-def δ-of-def*)
also have ... ∈ *O*[*?F*](*g*)
using *9 10 11 12 13 16*
by (*simp add*:*fun-cong*[*OF s-of-def*[*symmetric*]] *fun-cong*[*OF t-of-def*[*symmetric*]] *fun-cong*[*OF r-of-def*[*symmetric*]] *Let-def*) (*intro sum-in-bigo*, *auto*)
also have ... = *O*[*?F*](*?rhs*)
by (*simp add*:*case-prod-beta′ g-def n-of-def ε-of-def δ-of-def*)
finally show *?thesis*
by *simp*
qed

end

7 Frequency Moment 2

theory *Frequency-Moment-2*
imports
Universal-Hash-Families.Carter-Wegman-Hash-Family
Equivalence-Relation-Enumeration.Equivalence-Relation-Enumeration
Landau-Ext
Median-Method.Median
Probability-Ext
Universal-Hash-Families.Universal-Hash-Families-More-Product-PMF
Frequency-Moments
begin

hide-const (**open**) *Discrete-Topology.discrete*
hide-const (**open**) *Isolated.discrete*

This section contains a formalization of the algorithm for the second frequency moment. It is based on the algorithm described in [1, §2.2]. The only difference is that the algorithm is adapted to work with prime field of odd order, which greatly reduces the implementation complexity.

fun *f2-hash* **where**
f2-hash p h k = (*if even* (*ring.hash* (*ring-of* (*mod-ring p*)) *k h*) *then int p* − *1 else* − *int p* − *1*)

type-synonym *f2-state* = *nat* × *nat* × *nat* × (*nat* × *nat* ⇒ *nat list*) × (*nat* × *nat* ⇒ *int*)

fun *f2-init* :: *rat* ⇒ *rat* ⇒ *nat* ⇒ *f2-state pmf* **where**
f2-init δ ε *n* =
do {
let s_1 = *nat* $\lceil 6 / \delta^2 \rceil$;
let s_2 = *nat* $\lceil -(18 * ln\ (real\text{-}of\text{-}rat\ \varepsilon)) \rceil$;
let p = *prime-above* (*max n 3*);
h ← *prod-pmf* ({..<s_1}×{..<s_2}) (λ-. *pmf-of-set* (*bounded-degree-polynomials* (*ring-of* (*mod-ring p*)) *4*));
return-pmf (s_1, s_2, *p*, *h*, (λ- ∈ {..<s_1} × {..<s_2}. (*0* :: *int*)))
}

fun *f2-update* :: *nat* ⇒ *f2-state* ⇒ *f2-state pmf* **where**
f2-update x (s_1, s_2, *p*, *h*, *sketch*) =
return-pmf (s_1, s_2, *p*, *h*, λ*i* ∈ {..<s_1} × {..<s_2}. *f2-hash p* (*h i*) *x* + *sketch i*)

fun *f2-result* :: *f2-state* ⇒ *rat pmf* **where**
f2-result (s_1, s_2, *p*, *h*, *sketch*) =
return-pmf (*median* s_2 (λi_2 ∈ {..<s_2}.
($\sum i_1 \in \{..<s_1\}$. (*rat-of-int* (*sketch* (i_1, i_2)))^2) / (((*rat-of-nat p*)^2−*1*) ∗ *rat-of-nat* s_1)))

fun *f2-space-usage* :: (*nat* × *nat* × *rat* × *rat*) ⇒ *real* **where**
f2-space-usage (*n*, *m*, ε, δ) = (
let s_1 = *nat* $\lceil 6 / \delta^2 \rceil$ *in*
let s_2 = *nat* $\lceil -(18 * ln\ (real\text{-}of\text{-}rat\ \varepsilon)) \rceil$ *in*
3 +
2 ∗ *log 2* (s_1 + *1*) +
2 ∗ *log 2* (s_2 + *1*) +
2 ∗ *log 2* (*9* + *2* ∗ *real n*) +
s_1 ∗ s_2 ∗ (*5* + *4*∗*log 2* (*8* + *2* ∗ *real n*) + *2* ∗ *log 2* (*real m* ∗ (*18* + *4* ∗ *real n*) + *1*)))

definition *encode-f2-state* :: *f2-state* ⇒ *bool list option* **where**
encode-f2-state =
N_e $\bowtie_e$ (λs_1.
N_e $\bowtie_e$ (λs_2.
N_e $\bowtie_e$ (λ*p*.
(*List.product* [*0*..<s_1] [*0*..<s_2] $\rightarrow_e$ P_e *p 4*) $\times_e$
(*List.product* [*0*..<s_1] [*0*..<s_2] $\rightarrow_e$ I_e))))

lemma *inj-on encode-f2-state* (*dom encode-f2-state*)
proof −
  have *is-encoding encode-f2-state*
    unfolding *encode-f2-state-def*
     by (*intro dependent-encoding exp-golomb-encoding fun-encoding list-encoding int-encoding poly-encoding*)

  thus *?thesis*
    by (*rule encoding-imp-inj*)
qed

context
  fixes ε δ :: *rat*
  fixes *n* :: *nat*
  fixes *as* :: *nat list*
  fixes *result*
  assumes *ε-range*: $\varepsilon \in \{0<..<1\}$
  assumes *δ-range*: $\delta > 0$
  assumes *as-range*: *set as* $\subseteq \{..<n\}$
  defines *result* $\equiv$ *fold* (λa *state. state* $\ggg$ *f2-update a*) *as* (*f2-init* δ ε *n*) $\ggg$ *f2-result*
begin

private definition s_1 **where** $s_1 = nat \lceil 6 \ / \ \delta^2 \rceil$

lemma *s1-gt-0*: $s_1 > 0$
    using *δ-range* **by** (*simp add*:s_1*-def*)

private definition s_2 **where** $s_2 = nat \lceil -(18 * \ ln \ (real\text{-}of\text{-}rat \ \varepsilon)) \rceil$

lemma *s2-gt-0*: $s_2 > 0$
    using *ε-range* **by** (*simp add*:s_2*-def*)

private definition *p* **where** *p* = *prime-above* (*max n 3*)

lemma *p-prime*: *Factorial-Ring.prime p*
  unfolding *p-def* **using** *prime-above-prime* **by** *blast*

lemma *p-ge-3*: $p \geq 3$
    unfolding *p-def* **by** (*meson max.boundedE prime-above-lower-bound*)

lemma *p-gt-0*: $p > 0$ **using** *p-ge-3* **by** *linarith*

lemma *p-gt-1*: $p > 1$ **using** *p-ge-3* **by** *simp*

lemma *p-ge-n*: $p \geq n$ **unfolding** *p-def*
  by (*meson max.boundedE prime-above-lower-bound*)

interpretation *carter-wegman-hash-family ring-of* (*mod-ring p*) *4*
using *carter-wegman-hash-familyI*[*OF mod-ring-is-field mod-ring-finite*]
using *p-prime* **by** *auto*

definition *sketch* **where** *sketch* = *fold* (λ*a state. state* $\gg=$ *f2-update a*) *as* (*f2-init* δ ε *n*)
private definition Ω **where**Ω = *prod-pmf* ($\{..<s_1\} \times \{..<s_2\}$) ($\lambda$-. *pmf-of-set space*)
private definition Ω_p **where**Ω_p = *measure-pmf* Ω
private definition *sketch-rv* **where** *sketch-rv* ω = *of-int* (*sum-list* (*map* (*f2-hash p* ω) *as*))^*2*
private definition *mean-rv* **where** *mean-rv* ω = (λi_2. ($\sum i_1 = 0..<s_1$. *sketch-rv* (ω (i_1, i_2))) / (((*of-nat p*)^2 − *1*) ∗ *of-nat* s_1))
private definition *result-rv* **where** *result-rv* ω = *median* s_2 ($\lambda i_2 \in \{..<s_2\}$. *mean-rv* ω i_2)

lemma *mean-rv-alg-sketch*:
sketch = Ω $\gg=$ ($\lambda\omega$. *return-pmf* (s_1, s_2, *p*, ω, $\lambda i \in \{..<s_1\} \times \{..<s_2\}$. *sum-list* (*map* (*f2-hash p* (ω *i*)) *as*)))
proof −
have *sketch* = *fold* (λ*a state. state* $\gg=$ *f2-update a*) *as* (*f2-init* δ ε *n*)
by (*simp add*:*sketch-def*)
also have ... = Ω $\gg=$ ($\lambda\omega$. *return-pmf* (s_1, s_2, *p*, ω,
$\lambda i \in \{..<s_1\} \times \{..<s_2\}$. *sum-list* (*map* (*f2-hash p* (ω *i*)) *as*)))
proof (*induction as rule*:*rev-induct*)
case *Nil*
then show *?case*
by (*simp add*:s_1*-def* s_2*-def space-def p-def*[*symmetric*] Ω*-def restrict-def Let-def*)
next
case (*snoc a as*)
have *fold* (λ*a state. state* $\gg=$ *f2-update a*) (*as* @ [*a*]) (*f2-init* δ ε *n*) = Ω $\gg=$
($\lambda\omega$. *return-pmf* (s_1, s_2, *p*, ω, $\lambda s \in \{..<s_1\} \times \{..<s_2\}$. ($\sum x \leftarrow as$. *f2-hash p* ($\omega$ *s*) *x*)) $\gg=$ *f2-update a*)
using *snoc* **by** (*simp add*: *bind-assoc-pmf restrict-def del*:*f2-hash.simps f2-init.simps*)
also have ... = Ω $\gg=$ ($\lambda\omega$. *return-pmf* (s_1, s_2, *p*, ω, $\lambda i \in \{..<s_1\} \times \{..<s_2\}$. ($\sum x \leftarrow as@[a]$. *f2-hash p* ($\omega$ *i*) *x*)))
by (*subst bind-return-pmf*) (*simp add*: *add.commute del*:*f2-hash.simps cong*:*restrict-cong*)
finally show *?case* **by** *blast*
qed
finally show *?thesis* **by** *auto*
qed

lemma *distr*: *result* = *map-pmf result-rv* Ω
proof −
have *result* = *sketch* $\gg=$ *f2-result*
by (*simp add*:*result-def sketch-def*)
also have ... = Ω $\gg=$ (λx. *f2-result* (s_1, s_2, *p*, *x*, $\lambda i \in \{..<s_1\} \times \{..<s_2\}$. *sum-list* (*map* (*f2-hash p* (*x i*)) *as*)))

by (*simp add*: *mean-rv-alg-sketch bind-assoc-pmf bind-return-pmf*)
also have ... = *map-pmf result-rv* Ω
by (*simp add*:*map-pmf-def result-rv-def mean-rv-def sketch-rv-def lessThan-atLeast0 cong*:*restrict-cong*)
finally show *?thesis* **by** *simp*
qed

private lemma *f2-hash-pow-exp*:
assumes $k < p$
shows
expectation (λω. *real-of-int* (*f2-hash p ω k*) ^*m*) =
((*real p* − *1*) ^ *m* ∗ (*real p* + *1*) + (− *real p* − *1*) ^ *m* ∗ (*real p* − *1*)) / (*2* ∗ *real p*)
proof −

have *odd p* **using** *p-prime p-ge-3 prime-odd-nat assms* **by** *simp*
then obtain *t* **where** *t-def*: *p=2∗t+1*
using *oddE* **by** *blast*

have *Collect even* ∩ {..<*2* ∗ *t* + *1*} ⊆ (∗) *2* ‘ {..<*t* + *1*}
by (*rule in-image-by-witness*[**where** *g*=λ*x*. *x div 2*], *simp*, *linarith*)
moreover have (∗) *2* ‘ {..<*t* + *1*} ⊆ *Collect even* ∩ {..<*2* ∗ *t* + *1*}
by (*rule image-subsetI*, *simp*)
ultimately have *card* ({*k*. *even k*} ∩ {..<*p*}) = *card* ((λ*x*. *2∗x*) ‘ {..<*t+1*})
unfolding *t-def* **using** *order-antisym* **by** *metis*
also have ... = *card* {..<*t+1*}
by (*rule card-image*, *simp add*: *inj-on-mult*)
also have ... = *t+1* **by** *simp*
finally have *card-even*: *card* ({*k*. *even k*} ∩ {..<*p*}) = *t+1* **by** *simp*
hence *card* ({*k*. *even k*} ∩ {..<*p*}) ∗ *2* = (*p+1*) **by** (*simp add*:*t-def*)
hence *prob-even*: *prob* {ω. *hash k* ω ∈ *Collect even*} = (*real p* + *1*)/(*2∗real p*)
using *assms*
by (*subst prob-range*, *auto simp*:*frac-eq-eq p-gt-0 mod-ring-def ring-of-def lessThan-def*)

have *p* = *card* {..<*p*} **by** *simp*
also have ... = *card* (({*k*. *odd k*} ∩ {..<*p*}) ∪ ({*k*. *even k*} ∩ {..<*p*}))
by (*rule arg-cong*[**where** *f*=*card*], *auto*)
also have ... = *card* ({*k*. *odd k*} ∩ {..<*p*}) + *card* ({*k*. *even k*} ∩ {..<*p*})
by (*rule card-Un-disjoint*, *simp*, *simp*, *blast*)
also have ... = *card* ({*k*. *odd k*} ∩ {..<*p*}) + *t+1*
by (*simp add*:*card-even*)
finally have *p* = *card* ({*k*. *odd k*} ∩ {..<*p*}) + *t+1*
by *simp*
hence *card* ({*k*. *odd k*} ∩ {..<*p*}) ∗ *2* = (*p−1*)
by (*simp add*:*t-def*)
hence *prob-odd*: *prob* {ω. *hash k* ω ∈ *Collect odd*} = (*real p* − *1*)/(*2∗real p*)
using *assms*
by (*subst prob-range*, *auto simp add*: *frac-eq-eq mod-ring-def ring-of-def lessThan-def*)

have *expectation* $(\lambda x.\ real\text{-}of\text{-}int\ (f2\text{-}hash\ p\ x\ k)\ \hat{}\ m) =$
$expectation\ (\lambda\omega.\ indicator\ \{\omega.\ even\ (hash\ k\ \omega)\}\ \omega * (real\ p - 1)\hat{}m +$
$indicator\ \{\omega.\ odd\ (hash\ k\ \omega)\}\ \omega * (-real\ p - 1)\hat{}m)$
by (*rule Bochner-Integration.integral-cong*, *simp*, *simp*)
also have ... =
$prob\ \{\omega.\ hash\ \ k\ \omega \in Collect\ even\}\ * (real\ p - 1)\ \hat{}\ m\ +$
$prob\ \{\omega.\ hash\ \ k\ \omega \in Collect\ odd\}\ * (-real\ p - 1)\ \hat{}\ m$
by (*simp*, *simp add*:*M-def*)
also have ... = $(real\ p + 1) * (real\ p - 1)\ \hat{}\ m\ /\ (2 * real\ p) + (real\ p - 1) * (-\ real\ p - 1)\ \hat{}\ m\ /\ (2 * real\ p)$
by (*subst prob-even*, *subst prob-odd*, *simp*)
also have ... =
$((real\ p - 1)\ \hat{}\ m * (real\ p + 1) + (-\ real\ p - 1)\ \hat{}\ m * (real\ p - 1))\ /\ (2 * real\ p)$
by (*simp add*:*add-divide-distrib ac-simps*)
finally show *expectation* $(\lambda x.\ real\text{-}of\text{-}int\ (f2\text{-}hash\ p\ x\ k)\ \hat{}\ m) =$
$((real\ p - 1)\ \hat{}\ m * (real\ p + 1) + (-\ real\ p - 1)\ \hat{}\ m * (real\ p - 1))\ /\ (2 * real\ p)$ **by** *simp*
qed

lemma
shows *var-sketch-rv*:*variance sketch-rv* $\leq 2*(real\text{-}of\text{-}rat\ (F\ 2\ as)\hat{}2) * ((real\ p)^2 - 1)^2$ (**is** *?A*)
and *exp-sketch-rv*:*expectation sketch-rv* $= real\text{-}of\text{-}rat\ (F\ 2\ as) * ((real\ p)^2 - 1)$ (**is** *?B*)
proof −
define *h* **where** $h = (\lambda\omega\ x.\ real\text{-}of\text{-}int\ (f2\text{-}hash\ p\ \omega\ x))$
define *c* **where** $c = (\lambda x.\ real\ (count\text{-}list\ as\ x))$
define *r* **where** $r = (\lambda(m::nat).\ ((real\ p - 1)\ \hat{}\ m * (real\ p + 1) + (-\ real\ p - 1)\ \hat{}\ m * (real\ p - 1))\ /\ (2 * real\ p))$
define *h-prod* **where** $h\text{-}prod = (\lambda as\ \omega.\ prod\text{-}list\ (map\ (h\ \omega)\ as))$

define *exp-h-prod* :: *nat list* $\Rightarrow$ *real* **where** $exp\text{-}h\text{-}prod = (\lambda as.\ (\prod i \in set\ as.\ r\ (count\text{-}list\ as\ i)))$

have *f-eq*: $sketch\text{-}rv = (\lambda\omega.\ (\sum x \in set\ as.\ c\ x * h\ \omega\ x)\hat{}2)$
by (*rule ext*, *simp add*:*sketch-rv-def c-def h-def sum-list-eval del*:*f2-hash.simps*)

have *r-one*: $r\ (Suc\ 0) = 0$
by (*simp add*:*r-def algebra-simps*)

have *r-two*: $r\ 2 = (real\ p\hat{}2 - 1)$
using *p-gt-0* **unfolding** *r-def power2-eq-square*
by (*simp add*:*nonzero-divide-eq-eq*, *simp add*:*algebra-simps*)

have$(real\ p)\hat{}2 \geq 2\hat{}2$
by (*rule power-mono*, *use p-gt-1* **in** *linarith*, *simp*)
hence *p-square-ge-4*: $(real\ p)^2 \geq 4$ **by** *simp*

have *r 4* = (*real p*)^*4*+*2*∗(*real p*)^2 − *3*
using *p-gt-0* **unfolding** *r-def*
by (*subst nonzero-divide-eq-eq, auto simp*:*power4-eq-xxxx power2-eq-square algebra-simps*)
also have ... ≤ (*real p*)^*4*+*2*∗(*real p*)^2 + *3*
by *simp*
also have ... ≤ *3* ∗ *r 2* ∗ *r 2*
using *p-square-ge-4*
by (*simp add*:*r-two power4-eq-xxxx power2-eq-square algebra-simps mult-left-mono*)
finally have *r-four-est*: *r 4* ≤ *3* ∗ *r 2* ∗ *r 2* **by** *simp*

have *exp-h-prod-elim*: *exp-h-prod* = (λ*as. prod-list* (*map* (*r* ∘ *count-list as*) (*remdups as*)))
by (*simp add*:*exp-h-prod-def prod.set-conv-list*[*symmetric*])

have *exp-h-prod*: ⋀*x. set x* ⊆ *set as* ⟹ *length x* ≤ *4* ⟹ *expectation* (*h-prod x*) = *exp-h-prod x*
proof −
fix *x*
assume *set x* ⊆ *set as*
hence *x-sub-p*: *set x* ⊆ {..<*p*} **using** *as-range p-ge-n* **by** *auto*
hence *x-le-p*: ⋀*k. k* ∈ *set x* ⟹ *k* < *p* **by** *auto*
assume *length x* ≤ *4*
hence *card-x*: *card* (*set x*) ≤ *4* **using** *card-length dual-order.trans* **by** *blast*

have *set x* ⊆ *carrier* (*ring-of* (*mod-ring p*))
using *x-sub-p* **by** (*simp add*:*mod-ring-def ring-of-def lessThan-def*)

hence *h-indep*: *indep-vars* (λ-. *borel*) (λ*i ω. h ω i* ^ *count-list x i*) (*set x*)
using *k-wise-indep-vars-subset*[*OF k-wise-indep*] *card-x as-range h-def*
by (*auto intro*:*indep-vars-compose2*[**where** *X*=*hash* **and** *M′*= (λ-. *discrete*)])

have *expectation* (*h-prod x*) = *expectation* (λ*ω*. ∏ *i* ∈ *set x. h ω i*^(*count-list x i*))
by (*simp add*:*h-prod-def prod-list-eval*)
also have ... = (∏ *i* ∈ *set x. expectation* (λ*ω. h ω i*^(*count-list x i*)))
by (*simp add*: *indep-vars-lebesgue-integral*[*OF - h-indep*])
also have ... = (∏ *i* ∈ *set x. r* (*count-list x i*))
using *f2-hash-pow-exp x-le-p*
by (*simp add*:*h-def r-def M-def*[*symmetric*] *del*:*f2-hash.simps*)
also have ... = *exp-h-prod x*
by (*simp add*:*exp-h-prod-def*)
finally show *expectation* (*h-prod x*) = *exp-h-prod x* **by** *simp*
qed

have ⋀*x y. kernel-of x* = *kernel-of y* ⟹ *exp-h-prod x* = *exp-h-prod y*
proof −
fix *x y* :: *nat list*
assume *a*:*kernel-of x* = *kernel-of y*

then obtain *f* **where** *b*:*bij-betw f* (*set x*) (*set y*) **and** *c*:$\bigwedge z.\ z \in set\ x \Longrightarrow$ *count-list x z* = *count-list y* (*f z*)
using *kernel-of-eq-imp-bij* **by** *blast*
have *exp-h-prod x* = *prod* ((λi. *r*(*count-list y i*)) $\circ$ *f*) (*set x*)
by (*simp add*:*exp-h-prod-def c*)
also have ... = ($\prod i \in f$ ‘ (*set x*). *r*(*count-list y i*))
by (*metis b bij-betw-def prod.reindex*)
also have ... = *exp-h-prod y*
unfolding *exp-h-prod-def*
by (*rule prod.cong*, *metis b bij-betw-def*) *simp*
finally show *exp-h-prod x* = *exp-h-prod y* **by** *simp*
qed

hence *exp-h-prod-cong*: $\bigwedge p\ x$. *of-bool* (*kernel-of x* = *kernel-of p*) $*$ *exp-h-prod p* =
of-bool (*kernel-of x* = *kernel-of p*) $*$ *exp-h-prod x*
by (*metis* (*full-types*) *of-bool-eq-0-iff vector-space-over-itself.scale-zero-left*)

have *c*:($\sum p \leftarrow$ *enum-rgfs n*. *of-bool* (*kernel-of xs* = *kernel-of p*) $*$ *r*) = *r*
if *a*:*length xs* = *n* **for** *xs* :: *nat list* **and** *n* **and** *r* :: *real*
proof −
have ($\sum p \leftarrow$ *enum-rgfs n*. *of-bool* (*kernel-of xs* = *kernel-of p*) $*$ *1*) = (*1*::*real*)
using *equiv-rels-2*[*OF a*[*symmetric*]] **by** (*simp add*:*equiv-rels-def comp-def*)
thus ($\sum p \leftarrow$ *enum-rgfs n*. *of-bool* (*kernel-of xs* = *kernel-of p*) $*$ *r*) = (*r*::*real*)
by (*simp add*:*sum-list-mult-const*)
qed

have *expectation sketch-rv* = ($\sum i \in$*set as*. ($\sum j \in$*set as*. *c i* $*$ *c j* $*$ *expectation* (*h-prod* [*i*,*j*])))
by (*simp add*:*f-eq h-prod-def power2-eq-square sum-distrib-left sum-distrib-right Bochner-Integration.integral-sum algebra-simps*)
also have ... = ($\sum i \in$*set as*. ($\sum j \in$*set as*. *c i* $*$ *c j* $*$ *exp-h-prod* [*i*,*j*]))
by (*simp add*:*exp-h-prod*)
also have ... = ($\sum i \in$ *set as*. ($\sum j \in$ *set as*.
c i $*$ *c j* $*$ (*sum-list* (*map* (λp. *of-bool* (*kernel-of* [*i*,*j*] = *kernel-of p*) $*$ *exp-h-prod p*) (*enum-rgfs 2*)))))
by (*subst exp-h-prod-cong*, *simp add*:*c*)
also have ... = ($\sum i \in$*set as*. *c i* $*$ *c i* $*$ *r 2*)
by (*simp add*: *numeral-eq-Suc kernel-of-eq All-less-Suc exp-h-prod-elim r-one distrib-left sum.distrib sum-collapse*)
also have ... = *real-of-rat* (*F 2 as*) $*$ ((*real p*)^*2*−*1*)
by (*simp add*: *sum-distrib-right*[*symmetric*] *c-def F-def power2-eq-square of-rat-sum of-rat-mult r-two*)
finally show *b*:*?B* **by** *simp*

have *expectation* (λx. (*sketch-rv x*)^2) = ($\sum$ *i1* $\in$ *set as*. ($\sum$ *i2* $\in$ *set as*. ($\sum$ *i3* $\in$ *set as*. ($\sum$ *i4* $\in$ *set as*.
c i1 $*$ *c i2* $*$ *c i3* $*$ *c i4* $*$ *expectation* (*h-prod* [*i1*, *i2*, *i3*, *i4*])))))
by (*simp add*:*f-eq h-prod-def power4-eq-xxxx sum-distrib-left sum-distrib-right*

Bochner-Integration.integral-sum algebra-simps)
also have ... = ($\sum$ *i1* ∈ *set as*. ($\sum$ *i2* ∈ *set as*. ($\sum$ *i3* ∈ *set as*. ($\sum$ *i4* ∈ *set as*. *c i1* ∗ *c i2* ∗ *c i3* ∗ *c i4* ∗ *exp-h-prod* [*i1*,*i2*,*i3*,*i4*]))))
by (*simp add*:*exp-h-prod*)
also have ... = ($\sum$ *i1* ∈ *set as*. ($\sum$ *i2* ∈ *set as*. ($\sum$ *i3* ∈ *set as*. ($\sum$ *i4* ∈ *set as*. *c i1* ∗ *c i2* ∗ *c i3* ∗ *c i4* ∗
(*sum-list* (*map* (λ*p*. *of-bool* (*kernel-of* [*i1*,*i2*,*i3*,*i4*] = *kernel-of p*) ∗ *exp-h-prod p*) (*enum-rgfs 4*)))))))
by (*subst exp-h-prod-cong*, *simp add*:*c*)
also have ... =
3 ∗ ($\sum$ *i* ∈ *set as*. ($\sum$ *j* ∈ *set as*. *c i*^*2* ∗ *c j*^*2* ∗ *r 2* ∗ *r 2*)) + (($\sum$ *i* ∈ *set as*. *c i*^*4* ∗ *r 4*) − *3* ∗ ($\sum$ *i* ∈ *set as*. *c i* ^ *4* ∗ *r 2* ∗ *r 2*))
apply (*simp add*: *numeral-eq-Suc exp-h-prod-elim r-one*)
apply (*simp add*: *kernel-of-eq All-less-Suc numeral-eq-Suc distrib-left sum.distrib sum-collapse neq-commute of-bool-not-iff*)
apply (*simp add*: *algebra-simps sum-subtractf sum-collapse*)
apply (*simp add*: *sum-distrib-left algebra-simps*)
done
also have ... = *3* ∗ ($\sum$ *i* ∈ *set as*. *c i*^*2* ∗ *r 2*)^*2* + ($\sum$ *i* ∈ *set as*. *c i* ^ *4* ∗ (*r 4* − *3* ∗ *r 2* ∗ *r 2*))
by (*simp add*:*power2-eq-square sum-distrib-left algebra-simps sum-subtractf*)
also have ... = *3* ∗ ($\sum$ *i* ∈ *set as*. *c i*^*2*)^*2* ∗ (*r 2*)^*2* + ($\sum$ *i* ∈ *set as*. *c i* ^ *4* ∗ (*r 4* − *3* ∗ *r 2* ∗ *r 2*))
by (*simp add*:*power-mult-distrib sum-distrib-right*[*symmetric*])
also have ... ≤ *3* ∗ ($\sum$ *i* ∈ *set as*. *c i*^*2*)^*2* ∗ (*r 2*)^*2* + ($\sum$ *i* ∈ *set as*. *c i* ^ *4* ∗ *0*)
using *r-four-est*
by (*auto intro*!: *sum-nonpos simp add*:*mult-nonneg-nonpos*)
also have ... = *3* ∗ (*real-of-rat* (*F 2 as*)^*2*) ∗ $((real\ p)^2-1)^2$
by (*simp add*:*c-def r-two F-def of-rat-sum of-rat-power*)
finally have *expectation* (λ*x*. $(sketch\text{-}rv\ x)^2$) ≤ *3* ∗ (*real-of-rat* (*F 2 as*)^*2*) ∗ $((real\ p)^2-1)^2$
by *simp*

thus *variance sketch-rv* ≤ *2*∗(*real-of-rat* (*F 2 as*)^*2*) ∗ $((real\ p)^2-1)^2$
by (*simp add*: *variance-eq*, *simp add*:*power-mult-distrib b*)
qed

lemma *space-omega-1* [*simp*]: *Sigma-Algebra.space* Ω_p = *UNIV*
by (*simp add*:Ω_p-*def*)

interpretation Ω: *prob-space* Ω_p
by (*simp add*:Ω_p-*def prob-space-measure-pmf*)

lemma *integrable*-Ω:
fixes *f* :: ((*nat* × *nat*) ⇒ (*nat list*)) ⇒ *real*
shows *integrable* Ω_p *f*
unfolding Ω_p-*def* Ω-*def*
by (*rule integrable-measure-pmf-finite*, *auto intro*:*finite-PiE simp*:*set-prod-pmf*)

lemma *sketch-rv-exp*:
  assumes $i_2 < s_2$
  assumes $i_1 \in \{0..<s_1\}$
  shows $\Omega.expectation\ (\lambda\omega.\ sketch\text{-}rv\ (\omega\ (i_1,\ i_2))) = real\text{-}of\text{-}rat\ (F\ 2\ as) * ((real\ p)^2 - 1)$
proof −
  have $\Omega.expectation\ (\lambda\omega.\ (sketch\text{-}rv\ (\omega\ (i_1,\ i_2))) :: real) = expectation\ sketch\text{-}rv$
    using *integrable-Ω integrable-M assms*
    unfolding *Ω-def* Ω_p*-def M-def*
    by (*subst expectation-Pi-pmf-slice*, *auto*)
  also have ... $= (real\text{-}of\text{-}rat\ (F\ 2\ as)) * ((real\ p)^2 - 1)$
    using *exp-sketch-rv* **by** *simp*
  finally show *?thesis* **by** *simp*
qed

lemma *sketch-rv-var*:
  assumes $i_2 < s_2$
  assumes $i_1 \in \{0..<s_1\}$
  shows $\Omega.variance\ (\lambda\omega.\ sketch\text{-}rv\ (\omega\ (i_1,\ i_2))) \le 2 * (real\text{-}of\text{-}rat\ (F\ 2\ as))^2 * ((real\ p)^2 - 1)^2$
proof −
  have $\Omega.variance\ (\lambda\omega.\ (sketch\text{-}rv\ (\omega\ (i_1,\ i_2)) :: real)) = variance\ sketch\text{-}rv$
    using *integrable-Ω integrable-M assms*
    unfolding *Ω-def* Ω_p*-def M-def*
    by (*subst variance-prod-pmf-slice*, *auto*)
  also have ... $\le 2 * (real\text{-}of\text{-}rat\ (F\ 2\ as))^2 * ((real\ p)^2 - 1)^2$
    using *var-sketch-rv* **by** *simp*
  finally show *?thesis* **by** *simp*
qed

lemma *mean-rv-exp*:
  assumes $i < s_2$
  shows $\Omega.expectation\ (\lambda\omega.\ mean\text{-}rv\ \omega\ i) = real\text{-}of\text{-}rat\ (F\ 2\ as)$
proof −
  have $a{:}(real\ p)^2 > 1$ **using** *p-gt-1* **by** *simp*

  have $\Omega.expectation\ (\lambda\omega.\ mean\text{-}rv\ \omega\ i) = (\sum i_1 = 0..<s_1.\ \Omega.expectation\ (\lambda\omega.\ sketch\text{-}rv\ (\omega\ (i_1,\ i)))) / (((real\ p)^2 - 1) * real\ s_1)$
    using *assms integrable-Ω* **by** (*simp add:mean-rv-def*)
  also have ... $= (\sum i_1 = 0..<s_1.\ real\text{-}of\text{-}rat\ (F\ 2\ as) * ((real\ p)^2 - 1)) / (((real\ p)^2 - 1) * real\ s_1)$
    using *sketch-rv-exp*[*OF assms*] **by** *simp*
  also have ... $= real\text{-}of\text{-}rat\ (F\ 2\ as)$
    using *s1-gt-0 a* **by** *simp*
  finally show *?thesis* **by** *simp*
qed

lemma *mean-rv-var*:

assumes $i < s_2$
shows Ω.*variance* ($\lambda\omega$. *mean-rv* ω *i*) $\leq$ (*real-of-rat* (δ $*$ *F 2 as*))^2 / *3*
proof −
have *a*: Ω.*indep-vars* (λ-. *borel*) (λi_1 *x*. *sketch-rv* (*x* (i_1, *i*))) {*0*..<s_1}
using *assms*
unfolding Ω_p-*def* Ω-*def*
by (*intro indep-vars-restrict-intro'*[**where** *f*=*fst*])
(*auto simp add*: *restrict-dfl-def case-prod-beta lessThan-atLeast0*)

have *p-sq-ne-1*: (*real p*)^*2* $\neq$ *1*
by (*metis p-gt-1 less-numeral-extra*(*4*) *of-nat-power one-less-power pos2 semiring-char-0-class.of-nat-eq-1-iff*)

have *s1-bound*: *6* / (*real-of-rat* δ)^2 $\leq$ *real* s_1
unfolding s_1-*def*
by (*metis* (*mono-tags*, *opaque-lifting*) *of-rat-ceiling of-rat-divide of-rat-numeral-eq of-rat-power real-nat-ceiling-ge*)

have Ω.*variance* ($\lambda\omega$. *mean-rv* ω *i*) = Ω.*variance* ($\lambda\omega$. $\sum$ i_1 = *0*..<s_1. *sketch-rv* (ω (i_1, *i*))) / (((*real p*)^2 − *1*) $*$ *real* s_1)^2
unfolding *mean-rv-def* **by** (*subst* Ω.*variance-divide*[*OF integrable*-Ω], *simp*)
also have ... = ($\sum$ i_1 = *0*..<s_1. Ω.*variance* ($\lambda\omega$. *sketch-rv* (ω (i_1, *i*)))) / (((*real p*)^2 − *1*) $*$ *real* s_1)^2
by (*subst* Ω.*bienaymes-identity-full-indep*[*OF - - integrable*-Ω *a*]) (*auto simp*: Ω-*def* Ω_p-*def*)
also have ... $\leq$ ($\sum$ i_1 = *0*..<s_1. *2*$*$(*real-of-rat* (*F 2 as*)^*2*) $*$ ((*real p*)^2−*1*)^2) / (((*real p*)^2 − *1*) $*$ *real* s_1)^2
by (*rule divide-right-mono*, *rule sum-mono*[*OF sketch-rv-var*[*OF assms*]], *auto*)
also have ... = *2* $*$ (*real-of-rat* (*F 2 as*)^*2*) / *real* s_1
using *p-sq-ne-1 s1-gt-0* **by** (*subst frac-eq-eq*, *auto simp*:*power2-eq-square*)
also have ... $\leq$ *2* $*$ (*real-of-rat* (*F 2 as*)^*2*) / (*6* / (*real-of-rat* δ)^2)
using *s1-gt-0* δ-*range* **by** (*intro divide-left-mono mult-pos-pos s1-bound*) *auto*
also have ... = (*real-of-rat* (δ $*$ *F 2 as*))^2 / *3*
by (*simp add*:*of-rat-mult algebra-simps*)
finally show *?thesis* **by** *simp*
qed

lemma *mean-rv-bounds*:
assumes $i < s_2$
shows Ω.*prob* {ω. *real-of-rat* δ $*$ *real-of-rat* (*F 2 as*) < |*mean-rv* ω *i* − *real-of-rat* (*F 2 as*)|} $\leq$ *1*/*3*
proof (*cases as* = [])
case *True*
then show *?thesis*
using *assms* **by** (*subst mean-rv-def*, *subst sketch-rv-def*, *simp add*:*F-def*)
next
case *False*
hence *F 2 as* > *0* **using** *F-gr-0* **by** *auto*

hence *a*: $0 < \textit{real-of-rat}\ (\delta * F\ 2\ as)$
using *δ-range* **by** *simp*
have [*simp*]: $(\lambda\omega.\ \textit{mean-rv}\ \omega\ i) \in \textit{borel-measurable}\ \Omega_p$
by (*simp add*:*Ω-def* Ω_p*-def*)
have $\Omega.\textit{prob}\ \{\omega.\ \textit{real-of-rat}\ \delta * \textit{real-of-rat}\ (F\ 2\ as) < |\textit{mean-rv}\ \omega\ i - \textit{real-of-rat}\ (F\ 2\ as)|\} \le$
$\Omega.\textit{prob}\ \{\omega.\ \textit{real-of-rat}\ (\delta * F\ 2\ as) \le |\textit{mean-rv}\ \omega\ i - \textit{real-of-rat}\ (F\ 2\ as)|\}$
by (*rule Ω.pmf-mono*[*OF* Ω_p*-def*], *simp add*:*of-rat-mult*)
also have ... $\le$ $\Omega.\textit{variance}\ (\lambda\omega.\ \textit{mean-rv}\ \omega\ i)\ /\ (\textit{real-of-rat}\ (\delta * F\ 2\ as))^2$
using *Ω.Chebyshev-inequality*[**where** $a=\textit{real-of-rat}\ (\delta * F\ 2\ as)$ **and** $f=\lambda\omega.\ \textit{mean-rv}\ \omega\ i$,*simplified*]
a prob-space-measure-pmf[**where** $p=\Omega$] *mean-rv-exp*[*OF assms*] *integrable-Ω* **by** *simp*
also have ... $\le ((\textit{real-of-rat}\ (\delta * F\ 2\ as))^2/3)\ /\ (\textit{real-of-rat}\ (\delta * F\ 2\ as))^2$
by (*rule divide-right-mono*, *rule mean-rv-var*[*OF assms*], *simp*)
also have ... $= 1/3$ **using** *a* **by** *force*
finally show *?thesis* **by** *blast*
qed

lemma *f2-alg-correct′*:
$\mathcal{P}(\omega\ \textit{in measure-pmf result}.\ |\omega - F\ 2\ as| \le \delta * F\ 2\ as) \ge 1-\textit{of-rat}\ \varepsilon$
proof −
have *a*: $\Omega.\textit{indep-vars}\ (\lambda\text{-}.\ \textit{borel})\ (\lambda i\ \omega.\ \textit{mean-rv}\ \omega\ i)\ \{0..<s_2\}$
using *s1-gt-0* **unfolding** Ω_p*-def* *Ω-def*
by (*intro indep-vars-restrict-intro′*[**where** *f*=*snd*])
(*auto simp*: Ω_p*-def* *Ω-def mean-rv-def restrict-dfl-def*)

have *b*: $-\ 18 * \ln\ (\textit{real-of-rat}\ \varepsilon) \le \textit{real}\ s_2$
unfolding s_2*-def* **using** *of-nat-ceiling* **by** *auto*

have $1 - \textit{of-rat}\ \varepsilon \le \Omega.\textit{prob}\ \{\omega.\ |\textit{median}\ s_2\ (\textit{mean-rv}\ \omega) - \textit{real-of-rat}\ (F\ 2\ as)| \le \textit{of-rat}\ \delta * \textit{of-rat}\ (F\ 2\ as)\}$
using *ε-range Ω.median-bound-2*[*OF - a b*, **where** $\delta=\textit{real-of-rat}\ \delta * \textit{real-of-rat}\ (F\ 2\ as)$
and $\mu=\textit{real-of-rat}\ (F\ 2\ as)$] *mean-rv-bounds*
by *simp*
also have ... $= \Omega.\textit{prob}\ \{\omega.\ |\textit{real-of-rat}\ (\textit{result-rv}\ \omega) - \textit{of-rat}\ (F\ 2\ as)| \le \textit{of-rat}\ \delta * \textit{of-rat}\ (F\ 2\ as)\}$
by (*simp add*:*result-rv-def median-restrict lessThan-atLeast0 median-rat*[*OF s2-gt-0*]
mean-rv-def sketch-rv-def of-rat-divide of-rat-sum of-rat-mult of-rat-diff of-rat-power)
also have ... $= \Omega.\textit{prob}\ \{\omega.\ |\textit{result-rv}\ \omega - F\ 2\ as| \le \delta * F\ 2\ as\}$
by (*simp add*:*of-rat-less-eq of-rat-mult*[*symmetric*] *of-rat-diff*[*symmetric*] *set-eq-iff*)
finally have $\Omega.\textit{prob}\ \{y.\ |\textit{result-rv}\ y - F\ 2\ as| \le \delta * F\ 2\ as\} \ge 1-\textit{of-rat}\ \varepsilon$ **by** *simp*
thus *?thesis* **by** (*simp add*: *distr* Ω_p*-def*)
qed

lemma *f2-exact-space-usage'*:
  AE ω *in sketch . bit-count* (*encode-f2-state* ω) $\leq$ *f2-space-usage* (*n*, *length as*, ε, δ)
proof –
  have $p \leq 2 * max\ n\ 3 + 2$
    by (*subst p-def*, *rule prime-above-upper-bound*)
  also have ... $\leq 2 * n + 8$
    by (*cases* $n \leq 2$, *simp-all*)
  finally have *p-bound*: $p \leq 2 * n + 8$
    by *simp*
  have *bit-count* (N_e *p*) $\leq$ *ereal* (*2* $*$ *log 2* (*real p* + *1*) + *1*)
    by (*rule exp-golomb-bit-count*)
  also have ... $\leq$ *ereal* (*2* $*$ *log 2* (*2* $*$ *real n* + *9*) + *1*)
    using *p-bound* **by** *simp*
  finally have *p-bit-count*: *bit-count* (N_e *p*) $\leq$ *ereal* (*2* $*$ *log 2* (*2* $*$ *real n* + *9*) + *1*)
    by *simp*

  have *a*: *bit-count* (*encode-f2-state* (s_1, s_2, *p*, *y*, $\lambda i \in \{..<s_1\} \times \{..<s_2\}$.
      sum-list (*map* (*f2-hash p* (*y i*)) *as*))) $\leq$ *ereal* (*f2-space-usage* (*n*, *length as*, ε, δ))
     if *a*:$y \in \{..<s_1\} \times \{..<s_2\} \rightarrow_E$ *bounded-degree-polynomials* (*ring-of* (*mod-ring p*)) *4* **for** *y*
  proof –
    have $y \in$ *extensional* ($\{..<s_1\} \times \{..<s_2\}$) **using** *a PiE-iff* **by** *blast*
    hence *y-ext*: $y \in$ *extensional* (*set* (*List.product* [$0..<s_1$] [$0..<s_2$]))
      by (*simp add*:*lessThan-atLeast0*)

    have *h-bit-count-aux*: *bit-count* (P_e *p 4* (*y x*)) $\leq$ *ereal* (*4* + *4* $*$ *log 2* (*8* + *2* $*$ *real n*))
       if *b*:$x \in$ *set* (*List.product* [$0..<s_1$] [$0..<s_2$]) **for** *x*
    proof –
      have *y x* $\in$ *bounded-degree-polynomials* (*ring-of* (*mod-ring p*)) *4*
        using *b a* **by** *force*
      hence *bit-count* (P_e *p 4* (*y x*)) $\leq$ *ereal* (*real 4* $*$ (*log 2* (*real p*) + *1*))
        by (*rule bounded-degree-polynomial-bit-count*[*OF p-gt-1*])
      also have ... $\leq$ *ereal* (*real 4* $*$ (*log 2* (*8* + *2* $*$ *real n*) + *1*))
        using *p-gt-0 p-bound* **by** *simp*
      also have ... $\leq$ *ereal* (*4* + *4* $*$ *log 2* (*8* + *2* $*$ *real n*))
        by *simp*
      finally show *?thesis*
        by *blast*
    qed

    have *h-bit-count*:
     bit-count ((*List.product* [$0..<s_1$] [$0..<s_2$] $\rightarrow_e$ P_e *p 4*) *y*) $\leq$ *ereal* (*real* s_1 $*$ *real* s_2 $*$ (*4* + *4* $*$ *log 2* (*8* + *2* $*$ *real n*)))
      using *fun-bit-count-est*[**where** *e*=P_e *p 4*, *OF y-ext h-bit-count-aux*]
      by *simp*

have *sketch-bit-count-aux*:
bit-count (I_e (*sum-list* (*map* (*f2-hash p* (*y x*)) *as*))) $\leq$ *ereal* (*1* + *2* $*$ *log 2* (*real* (*length as*) $*$ (*18* + *4* $*$ *real n*) + *1*)) (**is** *?lhs* $\leq$ *?rhs*)
if $x \in \{0..<s_1\} \times \{0..<s_2\}$ **for** *x*
proof −
have |*sum-list* (*map* (*f2-hash p* (*y x*)) *as*)| $\leq$ *sum-list* (*map* (*abs* $\circ$ (*f2-hash p* (*y x*))) *as*)
by (*subst map-map*[*symmetric*]) (*rule sum-list-abs*)
also have ... $\leq$ *sum-list* (*map* (λ-. (*int p*+*1*)) *as*)
by (*rule sum-list-mono*) (*simp add*:*p-gt-0*)
also have ... = *int* (*length as*) $*$ (*int p*+*1*)
by (*simp add*: *sum-list-triv*)
also have ... $\leq$ *int* (*length as*) $*$ (*9*+*2*$*$(*int n*))
using *p-bound* **by** (*intro mult-mono*, *auto*)
finally have |*sum-list* (*map* (*f2-hash p* (*y x*)) *as*)| $\leq$ *int* (*length as*) $*$ (*9* + *2* $*$ *int n*) **by** *simp*
hence *?lhs* $\leq$ *ereal* (*2* $*$ *log 2* (*real-of-int* (*2*$*$ (*int* (*length as*) $*$ (*9* + *2* $*$ *int n*)) + *1*)) + *1*)
by (*rule int-bit-count-est*)
also have ... = *?rhs* **by** (*simp add*:*algebra-simps*)
finally show *?thesis* **by** *simp*
qed

have
bit-count ((*List.product* [*0*..<s_1] [*0*..<s_2] $\rightarrow_e$ I_e) ($\lambda i \in \{..<s_1\} \times \{..<s_2\}$. *sum-list* (*map* (*f2-hash p* (*y i*)) *as*)))
$\leq$ *ereal* (*real* (*length* (*List.product* [*0*..<s_1] [*0*..<s_2]))) $*$ (*ereal* (*1* + *2* $*$ *log 2* (*real* (*length as*) $*$ (*18* + *4* $*$ *real n*) + *1*)))
by (*intro fun-bit-count-est*)
(*simp-all add*:*extensional-def lessThan-atLeast0 sketch-bit-count-aux del*:*f2-hash.simps*)
also have ... = *ereal* (*real* s_1 $*$ *real* s_2 $*$ (*1* + *2* $*$ *log 2* (*real* (*length as*) $*$ (*18* + *4* $*$ *real n*) + *1*)))
by *simp*
finally have *sketch-bit-count*:
bit-count ((*List.product* [*0*..<s_1] [*0*..<s_2] $\rightarrow_e$ I_e) ($\lambda i \in \{..<s_1\} \times \{..<s_2\}$. *sum-list* (*map* (*f2-hash p* (*y i*)) *as*))) $\leq$
ereal (*real* s_1 $*$ *real* s_2 $*$ (*1* + *2* $*$ *log 2* (*real* (*length as*) $*$ (*18* + *4* $*$ *real n*) + *1*))) **by** *simp*

have *bit-count* (*encode-f2-state* (s_1, s_2, *p*, *y*, $\lambda i \in \{..<s_1\} \times \{..<s_2\}$. *sum-list* (*map* (*f2-hash p* (*y i*)) *as*))) $\leq$
bit-count (N_e s_1) + *bit-count* (N_e s_2) +*bit-count* (N_e *p*) +
bit-count ((*List.product* [*0*..<s_1] [*0*..<s_2] $\rightarrow_e$ P_e *p 4*) *y*) +
bit-count ((*List.product* [*0*..<s_1] [*0*..<s_2] $\rightarrow_e$ I_e) ($\lambda i \in \{..<s_1\} \times \{..<s_2\}$. *sum-list* (*map* (*f2-hash p* (*y i*)) *as*)))
by (*simp add*:*Let-def* s_1*-def* s_2*-def encode-f2-state-def dependent-bit-count add.assoc*)
also have ... $\leq$ *ereal* (*2* $*$ *log 2* (*real* s_1 + *1*) + *1*) + *ereal* (*2* $*$ *log 2* (*real* s_2

+ 1) + 1) + ereal (2 ∗ log 2 (2 ∗ real n + 9) + 1) +
(ereal (real s_1 ∗ real s_2) ∗ (4 + 4 ∗ log 2 (8 + 2 ∗ real n))) +
(ereal (real s_1 ∗ real s_2) ∗ (1 + 2 ∗ log 2 (real (length as) ∗ (18 + 4 ∗ real n) + 1)))
by (intro add-mono exp-golomb-bit-count p-bit-count, auto intro: h-bit-count sketch-bit-count)
also have ... = ereal (f2-space-usage (n, length as, ε, δ))
by (simp add:distrib-left add.commute s_1-def[symmetric] s_2-def[symmetric] Let-def)
finally show bit-count (encode-f2-state (s_1, s_2, p, y, λi∈{..<s_1} × {..<s_2}. sum-list (map (f2-hash p (y i)) as))) ≤
ereal (f2-space-usage (n, length as, ε, δ))
by simp
qed

have set-pmf Ω = {..<s_1} × {..<s_2} $\rightarrow_E$ bounded-degree-polynomials (ring-of (mod-ring p)) 4
by (simp add: Ω-def set-prod-pmf) (simp add: space-def)
thus ?thesis
by (simp add:mean-rv-alg-sketch AE-measure-pmf-iff del:f2-space-usage.simps, metis a)
qed

end

Main results of this section:

theorem f2-alg-correct:
assumes ε ∈ {0<..<1}
assumes δ > 0
assumes set as ⊆ {..<n}
defines Ω ≡ fold (λa state. state ⋙= f2-update a) as (f2-init δ ε n) ⋙= f2-result
shows $\mathcal{P}$(ω in measure-pmf Ω. |ω − F 2 as| ≤ δ ∗ F 2 as) ≥ 1−of-rat ε
using f2-alg-correct'[OF assms(1,2,3)] Ω-def **by** auto

theorem f2-exact-space-usage:
assumes ε ∈ {0<..<1}
assumes δ > 0
assumes set as ⊆ {..<n}
defines M ≡ fold (λa state. state ⋙= f2-update a) as (f2-init δ ε n)
shows AE ω in M. bit-count (encode-f2-state ω) ≤ f2-space-usage (n, length as, ε, δ)
using f2-exact-space-usage'[OF assms(1,2,3)]
by (subst (asm) sketch-def[OF assms(1,2,3)], subst M-def, simp)

theorem f2-asymptotic-space-complexity:
f2-space-usage ∈ O[at-top $\times_F$ at-top $\times_F$ at-right 0 $\times_F$ at-right 0](λ (n, m, ε, δ).
(ln (1 / of-rat ε)) / (of-rat δ)^2 ∗ (ln (real n) + ln (real m)))
(**is** - ∈ O[?F](?rhs))
proof −

```
define n-of :: nat × nat × rat × rat ⇒ nat where n-of = (λ(n, m, ε, δ). n)
define m-of :: nat × nat × rat × rat ⇒ nat where m-of = (λ(n, m, ε, δ). m)
define ε-of :: nat × nat × rat × rat ⇒ rat where ε-of = (λ(n, m, ε, δ). ε)
define δ-of :: nat × nat × rat × rat ⇒ rat where δ-of = (λ(n, m, ε, δ). δ)

define g where g = (λx. (1/ (of-rat (δ-of x))^2) ∗ (ln (1 / of-rat (ε-of x))) ∗ (ln
(real (n-of x)) + ln (real (m-of x))))

have evt: (⋀x.
  0 < real-of-rat (δ-of x) ∧ 0 < real-of-rat (ε-of x) ∧
  1/real-of-rat (δ-of x) ≥ δ ∧ 1/real-of-rat (ε-of x) ≥ ε ∧
  real (n-of x) ≥ n ∧ real (m-of x) ≥ m⟹ P x)
  ⟹ eventually P ?F  (is (⋀x. ?prem x ⟹ -) ⟹ -)
  for δ ε n m P
  apply (rule eventually-mono[where P=?prem and Q=P])
  apply (simp add:ε-of-def case-prod-beta' δ-of-def n-of-def m-of-def)
   apply (intro eventually-conj eventually-prod1' eventually-prod2'
      sequentially-inf eventually-at-right-less inv-at-right-0-inf)
  by (auto simp add:prod-filter-eq-bot)

have unit-1: (λ-. 1) ∈ O[?F](λx. 1 / (real-of-rat (δ-of x))^2)
  using one-le-power
 by (intro landau-o.big-mono evt[where δ=1], auto simp add:power-one-over[symmetric])

have unit-2: (λ-. 1) ∈ O[?F](λx. ln (1 / real-of-rat (ε-of x)))
  by (intro landau-o.big-mono  evt[where ε=exp 1])
   (auto intro!:iffD2[OF ln-ge-iff] simp add:abs-ge-iff)

have unit-3: (λ-. 1) ∈ O[?F](λx. real (n-of x))
  using of-nat-le-iff by (intro landau-o.big-mono evt; fastforce)

have unit-4: (λ-. 1) ∈ O[?F](λx. real (m-of x))
  using of-nat-le-iff by (intro landau-o.big-mono evt; fastforce)

have unit-5: (λ-. 1) ∈ O[?F](λx. ln (real (n-of x)))
  by (auto intro!: landau-o.big-mono evt[where n=exp 1])
    (metis abs-ge-self linorder-not-le ln-ge-iff not-exp-le-zero order.trans)

have unit-6: (λ-. 1) ∈ O[?F](λx. ln (real (n-of x)) + ln (real (m-of x)))
  by (intro landau-sum-1 evt[where m=1 and n=1] unit-5 iffD2[OF ln-ge-iff])
auto

have unit-7: (λ-. 1) ∈ O[?F](λx. 1 / real-of-rat (ε-of x))
  by (intro landau-o.big-mono  evt[where ε=1], auto)

have unit-8: (λ-. 1) ∈ O[?F](g)
  unfolding g-def by (intro landau-o.big-mult-1 unit-1 unit-2 unit-6)

have unit-9: (λ-. 1) ∈ O[?F](λx. real (n-of x) ∗ real (m-of x))
```

by (*intro landau-o.big-mult-1 unit-3 unit-4*)

have $(\lambda x.\ 6 * (1\ /\ (\textit{real-of-rat}\ (\delta\textit{-of}\ x))^2)) \in O[?F](\lambda x.\ 1\ /\ (\textit{real-of-rat}\ (\delta\textit{-of}\ x))^2)$
by (*subst landau-o.big.cmult-in-iff*, *simp-all*)
hence *l1*: $(\lambda x.\ \textit{real}\ (\textit{nat}\ \lceil 6\ /\ (\delta\textit{-of}\ x)^2 \rceil)) \in O[?F](\lambda x.\ 1\ /\ (\textit{real-of-rat}\ (\delta\textit{-of}\ x))^2)$
by (*intro landau-real-nat landau-rat-ceil*[*OF unit-1*]) (*simp-all add*:*of-rat-divide of-rat-power*)

have $(\lambda x.\ -\ (\ \ln\ (\textit{real-of-rat}\ (\varepsilon\textit{-of}\ x)))) \in O[?F](\lambda x.\ \ln\ (1\ /\ \textit{real-of-rat}\ (\varepsilon\textit{-of}\ x)))$
by (*intro landau-o.big-mono evt*) (*subst ln-div*, *auto*)
hence *l2*: $(\lambda x.\ \textit{real}\ (\textit{nat}\ \lceil -\ (18 * \ln\ (\textit{real-of-rat}\ (\varepsilon\textit{-of}\ x)))\rceil)) \in O[?F](\lambda x.\ \ln\ (1\ /\ \textit{real-of-rat}\ (\varepsilon\textit{-of}\ x)))$
by (*intro landau-real-nat landau-ceil*[*OF unit-2*], *simp*)

have *l3-aux*: $(\lambda x.\ \textit{real}\ (m\textit{-of}\ x) * (18 + 4 * \textit{real}\ (n\textit{-of}\ x)) + 1) \in O[?F](\lambda x.\ \textit{real}\ (n\textit{-of}\ x) * \textit{real}\ (m\textit{-of}\ x))$
by (*rule sum-in-bigo*[*OF -unit-9*], *subst mult.commute*)
(*intro landau-o.mult sum-in-bigo*, *auto simp*:*unit-3*)

note *of-nat-int-ceiling* [*simp del*]

have $(\lambda x.\ \ln\ (\textit{real}\ (m\textit{-of}\ x) * (18 + 4 * \textit{real}\ (n\textit{-of}\ x)) + 1)) \in O[?F](\lambda x.\ \ln\ (\textit{real}\ (n\textit{-of}\ x) * \textit{real}\ (m\textit{-of}\ x)))$
apply (*rule landau-ln-2*[**where** *a=2*], *simp*, *simp*)
apply (*rule evt*[**where** *m=2* **and** *n=1*])
apply (*metis dual-order.trans mult-left-mono mult-of-nat-commute of-nat-0-le-iff verit-prod-simplify*(*1*))
using *l3-aux* **by** *simp*
also have $(\lambda x.\ \ln\ (\textit{real}\ (n\textit{-of}\ x) * \textit{real}\ (m\textit{-of}\ x))) \in O[?F](\lambda x.\ \ln\ (\textit{real}\ (n\textit{-of}\ x)) + \ln(\textit{real}\ (m\textit{-of}\ x)))$
by (*intro landau-o.big-mono evt*[**where** *m=1* **and** *n=1*], *auto simp add*:*ln-mult*)
finally have *l3*: $(\lambda x.\ \ln\ (\textit{real}\ (m\textit{-of}\ x) * (18 + 4 * \textit{real}\ (n\textit{-of}\ x)) + 1)) \in O[?F](\lambda x.\ \ln\ (\textit{real}\ (n\textit{-of}\ x)) + \ln\ (\textit{real}\ (m\textit{-of}\ x)))$
using *landau-o.big-trans* **by** *simp*

have §: $(\lambda x.\ q + 2 * \textit{real}\ (n\textit{-of}\ x))$
$\in O[\textit{sequentially} \times_F \textit{sequentially} \times_F \textit{at-right}\ 0 \times_F \textit{at-right}\ 0](\lambda x.\ \textit{real}\ (n\textit{-of}\ x))$
if $q>0$ **for** *q*
using *that*
by (*auto intro*!: *sum-in-bigo simp add*:*unit-3*)
have *l4*: $(\lambda x.\ \ln\ (8 + 2 * \textit{real}\ (n\textit{-of}\ x))) \in O[?F](\lambda x.\ \ln\ (\textit{real}\ (n\textit{-of}\ x)) + \ln\ (\textit{real}\ (m\textit{-of}\ x)))$
by (*intro* § *landau-sum-1 evt*[**where** *m=1* **and** *n=2*] *landau-ln-2*[**where** *a=2*] *iffD2*[*OF ln-ge-iff*]) *auto*
have *l5*: $(\lambda x.\ \ln\ (9 + 2 * \textit{real}\ (n\textit{-of}\ x))) \in O[?F](\lambda x.\ \ln\ (\textit{real}\ (n\textit{-of}\ x)) + \ln\ (\textit{real}\ (m\textit{-of}\ x)))$
by (*intro* § *landau-sum-1 evt*[**where** *m=1* **and** *n=2*] *landau-ln-2*[**where** *a=2*]

iffD2[*OF ln-ge-iff*]) *auto*

have *l6*: $(\lambda x.\ ln\ (real\ (nat\ \lceil 6\ /\ (\delta\text{-}of\ x)^2\rceil) + 1)) \in O[?F](g)$
unfolding *g-def*
by (*intro landau-o.big-mult-1 landau-ln-3 sum-in-bigo unit-6 unit-2 l1 unit-1*, *simp*)

have *l7*: $(\lambda x.\ ln\ (9 + 2 * real\ (n\text{-}of\ x))) \in O[?F](g)$
unfolding *g-def*
by (*intro landau-o.big-mult-1′ unit-1 unit-2 l5*)

have *l8*: $(\lambda x.\ ln\ (real\ (nat\ \lceil -\ (18 * ln\ (real\text{-}of\text{-}rat\ (\varepsilon\text{-}of\ x)))\rceil) + 1)\ ) \in O[?F](g)$
unfolding *g-def*
by (*intro landau-o.big-mult-1 unit-6 landau-o.big-mult-1′ unit-1 landau-ln-3 sum-in-bigo l2 unit-2*) *simp*

have *l9*: $(\lambda x.\ 5 + 4 * ln\ (8 + 2 * real\ (n\text{-}of\ x))\ /\ ln\ 2 + 2 * ln\ (real\ (m\text{-}of\ x) * (18 + 4 * real\ (n\text{-}of\ x)) + 1)\ /\ ln\ 2)$
$\in O[?F](\lambda x.\ ln\ (real\ (n\text{-}of\ x)) + ln\ (real\ (m\text{-}of\ x)))$
by (*intro sum-in-bigo*, *auto simp*: *l3 l4 unit-6*)

have *l10*: $(\lambda x.\ real\ (nat\ \lceil 6\ /\ (\delta\text{-}of\ x)^2\rceil) * real\ (nat\ \lceil -\ (18 * ln\ (real\text{-}of\text{-}rat\ (\varepsilon\text{-}of\ x)))\rceil) *$
$(5 + 4 * ln\ (8 + 2 * real\ (n\text{-}of\ x))\ /\ ln\ 2 + 2 * ln(real\ (m\text{-}of\ x) * (18 + 4 * real\ (n\text{-}of\ x)) + 1)\ /\ ln\ 2))$
$\in O[?F](g)$
unfolding *g-def* **by** (*intro landau-o.mult*, *auto simp*: *l1 l2 l9*)

have *f2-space-usage* = $(\lambda x.\ f2\text{-}space\text{-}usage\ (n\text{-}of\ x,\ m\text{-}of\ x,\ \varepsilon\text{-}of\ x,\ \delta\text{-}of\ x))$
by (*simp add*:*case-prod-beta′ n-of-def ε-of-def δ-of-def m-of-def*)
also have ... $\in O[?F](g)$
by (*auto intro*!:*sum-in-bigo simp*:*Let-def log-def l6 l7 l8 l10 unit-8*)
also have ... = $O[?F](?rhs)$
by (*simp add*:*case-prod-beta′ g-def n-of-def ε-of-def δ-of-def m-of-def*)
finally show *?thesis* **by** *simp*
qed

end

8 Frequency Moment *k*

theory *Frequency-Moment-k*
imports
Frequency-Moments
Landau-Ext
Lp.Lp
Median-Method.Median
Probability-Ext
Universal-Hash-Families.Universal-Hash-Families-More-Product-PMF

begin

This section contains a formalization of the algorithm for the k-th frequency moment. It is based on the algorithm described in [1, §2.1].

```
type-synonym fk-state = nat × nat × nat × nat × (nat × nat ⇒ (nat × nat))

fun fk-init :: nat ⇒ rat ⇒ rat ⇒ nat ⇒ fk-state pmf where
  fk-init k δ ε n =
    do {
      let s₁ = nat ⌈3 ∗ real k ∗ n powr (1−1/real k) / (real-of-rat δ)²⌉;
      let s₂ = nat ⌈−18 ∗ ln (real-of-rat ε)⌉;
      return-pmf (s₁, s₂, k, 0, (λ- ∈ {0..<s₁} × {0..<s₂}. (0,0)))
    }

fun fk-update :: nat ⇒ fk-state ⇒ fk-state pmf where
  fk-update a (s₁, s₂, k, m, r) =
    do {
     coins ← prod-pmf ({0..<s₁} × {0..<s₂}) (λ-. bernoulli-pmf (1/(real m+1)));
      return-pmf (s₁, s₂, k, m+1, λi ∈ {0..<s₁} × {0..<s₂}.
        if coins i then
          (a,0)
        else (
          let (x,l) = r i in (x, l + of-bool (x=a))
        )
      )
    }

fun fk-result :: fk-state ⇒ rat pmf where
  fk-result (s₁, s₂, k, m, r) =
    return-pmf (median s₂ (λi₂ ∈ {0..<s₂}.
      (∑ i₁∈{0..<s₁}. rat-of-nat (let t = snd (r (i₁, i₂)) + 1 in m ∗ (t^k − (t −
1)^k))) / (rat-of-nat s₁))
    )

lemma bernoulli-pmf-1: bernoulli-pmf 1 = return-pmf True
  by (rule pmf-eqI, simp add:indicator-def)

fun fk-space-usage :: (nat × nat × nat × rat × rat) ⇒ real where
  fk-space-usage (k, n, m, ε, δ) = (
    let s₁ = nat ⌈3∗real k∗ (real n) powr (1−1/ real k) / (real-of-rat δ)² ⌉ in
    let s₂ = nat ⌈−(18 ∗ ln (real-of-rat ε))⌉ in
    4 +
    2 ∗ log 2 (s₁ + 1) +
    2 ∗ log 2 (s₂ + 1) +
    2 ∗ log 2 (real k + 1) +
    2 ∗ log 2 (real m + 1) +
    s₁ ∗ s₂ ∗ (2 + 2 ∗ log 2 (real n+1) + 2 ∗ log 2 (real m+1)))

definition encode-fk-state :: fk-state ⇒ bool list option where
```

encode-fk-state =
  $N_e \bowtie_e (\lambda s_1.$
  $N_e \bowtie_e (\lambda s_2.$
  $N_e \times_e$
  $N_e \times_e$
  $(List.product\ [0..<s_1]\ [0..<s_2] \rightarrow_e (N_e \times_e N_e))))$

lemma *inj-on encode-fk-state* (*dom encode-fk-state*)
proof −
  have *is-encoding encode-fk-state*
    by (*simp add*:*encode-fk-state-def*)
     (*intro dependent-encoding exp-golomb-encoding fun-encoding*)

  thus *?thesis* **by** (*rule encoding-imp-inj*)
qed

This is an intermediate non-parallel form *fk-update* used only in the correctness proof.

fun *fk-update-2* :: *'a* ⇒ (*nat* × *'a* × *nat*) ⇒ (*nat* × *'a* × *nat*) *pmf* **where**
  fk-update-2 a (*m*,*x*,*l*) =
    do {
      coin ← *bernoulli-pmf* (*1*/(*real m*+*1*));
      return-pmf (*m*+*1*,*if coin then* (*a*,*0*) *else* (*x*, *l* + *of-bool* (*x*=*a*)))
    }

definition *sketch* **where** *sketch as i* = (*as* ! *i*, *count-list* (*drop* (*i*+*1*) *as*) (*as* ! *i*))

lemma *fk-update-2-distr*:
  assumes *as* ≠ []
  shows *fold* (λ*x s. s* ≫= *fk-update-2 x*) *as* (*return-pmf* (*0*,*0*,*0*)) =
  pmf-of-set {..<*length as*} ≫= (λ*k. return-pmf* (*length as*, *sketch as k*))
  using *assms*
proof (*induction as rule*:*rev-nonempty-induct*)
  case (*single x*)
  show *?case* **using** *single*
    by (*simp add:bind-return-pmf pmf-of-set-singleton bernoulli-pmf-1 lessThan-def sketch-def*)
next
  case (*snoc x xs*)
  let *?h* = (λ*xs k. count-list* (*drop* (*Suc k*) *xs*) (*xs* ! *k*))
  let *?q* = (λ*xs k.* (*length xs*, *sketch xs k*))

  have *non-empty*: {..<*Suc* (*length xs*)} ≠ {} {..<*length xs*} ≠ {} **using** *snoc* **by** *auto*

  have *fk-update-2-eta*:*fk-update-2 x* = (λ*a. fk-update-2 x* (*fst a*, *fst* (*snd a*), *snd* (*snd a*)))
    by *auto*

have *pmf-of-set* {..<*length xs*} ≫= (λ*k*. *bernoulli-pmf* (*1* / (*real* (*length xs*) + *1*)) ≫=
(λ*coin*. *return-pmf* (*if coin then length xs else k*))) =
bernoulli-pmf (*1* / (*real* (*length xs*) + *1*)) ≫= (λ*y*. *pmf-of-set* {..<*length xs*} ≫=
(λ*k*. *return-pmf* (*if y then length xs else k*)))
by (*subst bind-commute-pmf*, *simp*)
also have ... = *pmf-of-set* {..<*length xs* + *1*}
using *snoc*(*1*) *non-empty*
by (*intro pmf-eqI*, *simp add*: *pmf-bind measure-pmf-of-set*)
(*simp add*:*indicator-def algebra-simps frac-eq-eq*)
finally have *b*: *pmf-of-set* {..<*length xs*} ≫= (λ*k*. *bernoulli-pmf* (*1* / (*real* (*length xs*) + *1*)) ≫=
(λ*coin*. *return-pmf* (*if coin then length xs else k*))) = *pmf-of-set* {..<*length xs* +*1*} **by** *simp*

have *fold* (λ*x s*. (*s* ≫= *fk-update-2 x*)) (*xs*@[*x*]) (*return-pmf* (*0*,*0*,*0*)) =
(*pmf-of-set* {..<*length xs*} ≫= (λ*k*. *return-pmf* (*length xs*, *sketch xs k*))) ≫= *fk-update-2 x*
using *snoc* **by** (*simp add*:*case-prod-beta′*)
also have ... = (*pmf-of-set* {..<*length xs*} ≫= (λ*k*. *return-pmf* (*length xs*, *sketch xs k*))) ≫=
(λ(*m*,*a*,*l*). *bernoulli-pmf* (*1* / (*real m* + *1*)) ≫= (λ*coin*.
return-pmf (*m* + *1*, *if coin then* (*x*, *0*) *else* (*a*, (*l* + *of-bool* (*a* = *x*))))))
by (*subst fk-update-2-eta*, *subst fk-update-2.simps*, *simp add*:*case-prod-beta′*)
also have ... = *pmf-of-set* {..<*length xs*} ≫= (λ*k*. *bernoulli-pmf* (*1* / (*real* (*length xs*) + *1*)) ≫=
(λ*coin*. *return-pmf* (*length xs* + *1*, *if coin then* (*x*, *0*) *else* (*xs* ! *k*, *?h xs k* + *of-bool* (*xs* ! *k* = *x*)))))
by (*subst bind-assoc-pmf*, *simp add*: *bind-return-pmf sketch-def*)
also have ... = *pmf-of-set* {..<*length xs*} ≫= (λ*k*. *bernoulli-pmf* (*1* / (*real* (*length xs*) + *1*)) ≫=
(λ*coin*. *return-pmf* (*if coin then length xs else k*) ≫= (λ*k′*. *return-pmf* (*?q* (*xs*@[*x*]) *k′*))))
using *non-empty*
by (*intro bind-pmf-cong*, *auto simp add*:*bind-return-pmf nth-append count-list-append sketch-def*)
also have ... = *pmf-of-set* {..<*length xs*} ≫= (λ*k*. *bernoulli-pmf* (*1* / (*real* (*length xs*) + *1*)) ≫=
(λ*coin*. *return-pmf* (*if coin then length xs else k*))) ≫= (λ*k′*. *return-pmf* (*?q* (*xs*@[*x*]) *k′*))
by (*subst bind-assoc-pmf*, *subst bind-assoc-pmf*, *simp*)
also have ... = *pmf-of-set* {..<*length* (*xs*@[*x*])} ≫= (λ*k′*. *return-pmf* (*?q* (*xs*@[*x*]) *k′*))
by (*subst b*, *simp*)
finally show *?case* **by** *simp*
qed

context

fixes ε δ :: *rat*
fixes *n k* :: *nat*
fixes *as*
assumes *k-ge-1*: $k \geq 1$
assumes *ε-range*: $\varepsilon \in \{0<..<1\}$
assumes *δ-range*: $\delta > 0$
assumes *as-range*: *set as* $\subseteq$ $\{..<n\}$
begin

definition s_1 **where** $s_1 = nat \lceil 3 * real\ k * (real\ n)\ powr\ (1-1/real\ k)\ /\ (real\text{-}of\text{-}rat\ \delta)^2 \rceil$
definition s_2 **where** $s_2 = nat \lceil -(18 * ln\ (real\text{-}of\text{-}rat\ \varepsilon)) \rceil$

definition $M_1 = \{(u, v).\ v < count\text{-}list\ as\ u\}$
definition $\Omega_1 = measure\text{-}pmf\ (pmf\text{-}of\text{-}set\ M_1)$

definition $M_2 = prod\text{-}pmf\ (\{0..<s_1\} \times \{0..<s_2\})\ (\lambda\text{-}.\ pmf\text{-}of\text{-}set\ M_1)$
definition $\Omega_2 = measure\text{-}pmf\ M_2$

interpretation *prob-space* Ω_1
unfolding Ω_1*-def* **by** (*simp add:prob-space-measure-pmf*)

interpretation Ω_2:*prob-space* Ω_2
unfolding Ω_2*-def* **by** (*simp add:prob-space-measure-pmf*)

lemma *split-space*: $(\sum a{\in}M_1.\ f\ (snd\ a)) = (\sum u \in set\ as.\ (\sum v \in\{0..<count\text{-}list\ as\ u\}.\ f\ v))$
proof −
define *A* **where** $A = (\lambda u.\ \{u\} \times \{v.\ v < count\text{-}list\ as\ u\})$

have *a*: *inj-on snd* (*A x*) **for** *x*
by (*simp add:A-def inj-on-def*)

have $\bigwedge u\ v.\ u < count\text{-}list\ as\ v \Longrightarrow v \in set\ as$
by (*subst count-list-gr-1*, *force*)
hence $M_1 = \bigcup\ (A\ `\ set\ as)$
by (*auto simp add:set-eq-iff A-def* M_1*-def*)
hence $(\sum a{\in}M_1.\ f\ (snd\ a)) = sum\ (f \circ snd)\ \ (\bigcup\ (A\ `\ set\ as))$
by (*intro sum.cong*, *auto*)
also have ... $= sum\ (\lambda x.\ sum\ (f \circ snd)\ (A\ x))\ (set\ as)$
by (*rule sum.UNION-disjoint*, *simp*, *simp add:A-def*, *simp add:A-def*, *blast*)
also have ... $= sum\ (\lambda x.\ sum\ f\ (snd\ `\ A\ x))\ (set\ as)$
by (*intro sum.cong*, *auto simp add:sum.reindex*[*OF a*])
also have ... $= (\sum u \in set\ as.\ (\sum v \in\{0..<count\text{-}list\ as\ u\}.\ f\ v))$
unfolding *A-def* **by** (*intro sum.cong*, *auto*)
finally show *?thesis* **by** *blast*
qed

lemma

```
  assumes as ≠ []
  shows fin-space: finite M₁
    and non-empty-space: M₁ ≠ {}
    and card-space: card M₁ = length as
proof −
  have M₁ ⊆ set as × {k. k < length as}
  proof (rule subsetI)
    fix x
    assume a:x ∈ M₁
    have fst x ∈ set as
      using a by (simp add:case-prod-beta count-list-gr-1 M₁-def)
    moreover have snd x < length as
      using a count-le-length order-less-le-trans
      by (simp add:case-prod-beta M₁-def) fast
    ultimately show x ∈ set as × {k. k < length as}
      by (simp add:mem-Times-iff)
  qed
  thus fin-space: finite M₁
    using finite-subset by blast

  have (as ! 0, 0) ∈ M₁
    using assms(1) unfolding M₁-def
   by (simp, metis count-list-gr-1 gr0I length-greater-0-conv not-one-le-zero nth-mem)
  thus M₁ ≠ {} by blast

  show card M₁ = length as
    using fin-space split-space[where f=λ-. (1::nat)]
    by (simp add:sum-count-set[where X=set as and xs=as, simplified])
qed

lemma
  assumes as ≠ []
  shows integrable-1: integrable Ω₁ (f :: - ⇒ real) and
    integrable-2: integrable Ω₂ (g :: - ⇒ real)
proof −
  have fin-omega: finite (set-pmf (pmf-of-set M₁))
    using fin-space[OF assms] non-empty-space[OF assms] by auto
  thus integrable Ω₁ f
    unfolding Ω₁-def
    by (rule integrable-measure-pmf-finite)

  have finite (set-pmf M₂)
    unfolding M₂-def using fin-omega
    by (subst set-prod-pmf) (auto intro:finite-PiE)

  thus integrable Ω₂ g
    unfolding Ω₂-def by (intro integrable-measure-pmf-finite)
qed
```

lemma *sketch-distr*:
  assumes *as* $\neq$ []
  shows *pmf-of-set* {..<*length as*} $\gg=$ (λk. *return-pmf* (*sketch as k*)) = *pmf-of-set* M_1
proof –
  have $x < y \Longrightarrow y <$ *length as* $\Longrightarrow$
    count-list (*drop* (*y+1*) *as*) (*as* ! *y*) < *count-list* (*drop* (*x+1*) *as*) (*as* ! *y*) **for** *x y*
    by (*intro count-list-lt-suffix suffix-drop-drop*, *simp-all*)
     (*metis Suc-diff-Suc diff-Suc-Suc diff-add-inverse lessI less-natE*)
  hence *a1*: *inj-on* (*sketch as*) {*k*. *k* < *length as*}
     unfolding *sketch-def* **by** (*intro inj-onI*) (*metis Pair-inject mem-Collect-eq nat-neq-iff*)

  have *x* < *length as* $\Longrightarrow$ *count-list* (*drop* (*x+1*) *as*) (*as* ! *x*) < *count-list as* (*as* ! *x*) **for** *x*
    by (*rule count-list-lt-suffix*, *auto simp add*:*suffix-drop*)
  hence *sketch as* ‘ {*k*. *k* < *length as*} $\subseteq M_1$
    by (*intro image-subsetI*, *simp add*:*sketch-def* M_1-*def*)
  moreover have *card* $M_1 \leq$ *card* (*sketch as* ‘ {*k*. *k* < *length as*})
    by (*simp add*: *card-space*[*OF assms*(*1*)] *card-image*[*OF a1*])
  ultimately have *sketch as* ‘ {*k*. *k* < *length as*} = M_1
    using *fin-space*[*OF assms*(*1*)] **by** (*intro card-seteq*, *simp-all*)
  hence *bij-betw* (*sketch as*) {*k*. *k* < *length as*} M_1
    using *a1* **by** (*simp add*:*bij-betw-def*)
  hence *map-pmf* (*sketch as*) (*pmf-of-set* {*k*. *k* < *length as*}) = *pmf-of-set* M_1
    using *assms* **by** (*intro map-pmf-of-set-bij-betw*, *auto*)
  thus *?thesis* **by** (*simp add*: *sketch-def map-pmf-def lessThan-def*)
qed

lemma *fk-update-distr*:
  fold (λx *s*. *s* $\gg=$ *fk-update x*) *as* (*fk-init k* δ ε *n*) =
  prod-pmf ({*0*..<s_1} × {*0*..<s_2}) (λ-. *fold* (λx *s*. *s* $\gg=$ *fk-update-2 x*) *as* (*return-pmf* (*0*,*0*,*0*)))
    $\gg=$ (λx. *return-pmf* (s_1,s_2,*k*, *length as*, $\lambda i \in$ {*0*..<s_1}×{*0*..<s_2}. *snd* (*x i*)))
proof (*induction as rule*:*rev-induct*)
  case *Nil*
  then show *?case*
    by (*auto simp*:*Let-def* s_1-*def*[*symmetric*] s_2-*def*[*symmetric*] *bind-return-pmf*)
next
  case (*snoc x xs*)

  have *fk-update-2-eta*:*fk-update-2 x* = (λa. *fk-update-2 x* (*fst a*, *fst* (*snd a*), *snd* (*snd a*)))
    by *auto*

  have *a*: *fk-update x* (s_1, s_2, *k*, *length xs*, $\lambda i \in$ {*0*..<s_1} × {*0*..<s_2}. *snd* (*f i*)) =
    prod-pmf ({*0*..<s_1} × {*0*..<s_2}) (λi. *fk-update-2 x* (*f i*)) $\gg=$
    (λa. *return-pmf* (s_1,s_2, *k*, *Suc* (*length xs*), $\lambda i \in$ {*0*..<s_1} × {*0*..<s_2}. *snd* (*a i*)))
    if *b*: *f* $\in$ *set-pmf* (*prod-pmf* ({*0*..<s_1} × {*0*..<s_2})

$(\lambda$-. *fold* $(\lambda a\ s.\ s \gg= \mathit{fk\text{-}update\text{-}2}\ a)$ *xs* (*return-pmf* $(0,\ 0,\ 0))))$ **for** f
proof −
have c:*fst* $(f\ i)$ = *length xs* **if** d:$i \in \{0..<s_1\} \times \{0..<s_2\}$ **for** i
proof (*cases xs* = [])
case *True*
then show *?thesis* **using** $b\ d$ **by** (*simp add*: *set-Pi-pmf*)
next
case *False*
hence $\{..<\mathit{length}\ xs\} \neq \{\}$ **by** *force*
thus *?thesis* **using** $b\ d$
by (*simp add*:*set-Pi-pmf fk-update-2-distr*[*OF False*] *PiE-dflt-def*) *force*
qed
show *?thesis*
apply (*subst fk-update-2-eta*, *subst fk-update-2.simps*, *simp*)
apply (*simp add*: *Pi-pmf-bind-return*[**where** d'=*undefined*] *bind-assoc-pmf*)
apply (*rule bind-pmf-cong*, *simp add*:*c cong*:*Pi-pmf-cong*)
by (*auto simp add*:*bind-return-pmf case-prod-beta*)
qed

have *fold* $(\lambda x\ s.\ s \gg= \mathit{fk\text{-}update}\ x)$ $(xs\ @\ [x])$ $(\mathit{fk\text{-}init}\ k\ \delta\ \varepsilon\ n)$ =
prod-pmf $(\{0..<s_1\} \times \{0..<s_2\})$ $(\lambda$-. *fold* $(\lambda x\ s.\ s \gg= \mathit{fk\text{-}update\text{-}2}\ x)$ *xs* (*return-pmf* $(0,0,0)))$
$\gg=$ $(\lambda\omega.$ *return-pmf* $(s_1,s_2,k,$ *length xs*, $\lambda i\in\{0..<s_1\}\times\{0..<s_2\}.\ \mathit{snd}\ (\omega\ i)) \gg=$ *fk-update* $x)$
using *snoc*
by (*simp add*:*restrict-def bind-assoc-pmf del*:*fk-init.simps*)
also have ... = *prod-pmf* $(\{0..<s_1\} \times \{0..<s_2\})$
$(\lambda$-. *fold* $(\lambda a\ s.\ s \gg= \mathit{fk\text{-}update\text{-}2}\ a)$ *xs* (*return-pmf* $(0,\ 0,\ 0)))$ $\gg=$
$(\lambda f.$ *prod-pmf* $(\{0..<s_1\} \times \{0..<s_2\})$ $(\lambda i.\ \mathit{fk\text{-}update\text{-}2}\ x\ (f\ i))$ $\gg=$
$(\lambda a.$ *return-pmf* $(s_1,\ s_2,\ k,$ *Suc* (*length xs*), $\lambda i\in\{0..<s_1\} \times \{0..<s_2\}.\ \mathit{snd}\ (a\ i))))$
using a
by (*intro bind-pmf-cong*, *simp-all add*:*bind-return-pmf del*:*fk-update.simps*)
also have ... = *prod-pmf* $(\{0..<s_1\} \times \{0..<s_2\})$
$(\lambda$-. *fold* $(\lambda a\ s.\ s \gg= \mathit{fk\text{-}update\text{-}2}\ a)$ *xs* (*return-pmf* $(0,\ 0,\ 0)))$ $\gg=$
$(\lambda f.$ *prod-pmf* $(\{0..<s_1\} \times \{0..<s_2\})$ $(\lambda i.\ \mathit{fk\text{-}update\text{-}2}\ x\ (f\ i)))$ $\gg=$
$(\lambda a.$ *return-pmf* $(s_1,\ s_2,\ k,$ *Suc* (*length xs*), $\lambda i\in\{0..<s_1\} \times \{0..<s_2\}.\ \mathit{snd}\ (a\ i)))$
by (*simp add*:*bind-assoc-pmf*)
also have ... = (*prod-pmf* $(\{0..<s_1\} \times \{0..<s_2\})$
$(\lambda$-. *fold* $(\lambda a\ s.\ s \gg= \mathit{fk\text{-}update\text{-}2}\ a)$ $(xs@[x])$ (*return-pmf* $(0,0,0)))$
$\gg=$ $(\lambda a.$ *return-pmf* $(s_1,s_2,k,$ *length* $(xs@[x])$, $\lambda i\in\{0..<s_1\}\times\{0..<s_2\}.\ \mathit{snd}\ (a\ i))))$
by (*simp*, *subst Pi-pmf-bind*, *auto*)

finally show *?case* **by** *blast*
qed

lemma *power-diff-sum*:

fixes $a\ b :: {}'a :: \{comm\text{-}ring\text{-}1, power\}$
assumes $k > 0$
shows $a\hat{}k - b\hat{}k = (a-b) * (\sum i = 0..<k.\ a\ \hat{}\ i * b\ \hat{}\ (k - 1 - i))$ (**is** $?lhs = ?rhs$)
proof $-$
have *insert-lb*: $m < n \Longrightarrow insert\ m\ \{Suc\ m..<n\} = \{m..<n\}$ **for** $m\ n :: nat$
by *auto*

have $?rhs = sum\ (\lambda i.\ a * (a\hat{}i * b\hat{}(k-1-i)))\ \{0..<k\} -$
$sum\ (\lambda i.\ b * (a\hat{}i * b\hat{}(k-1-i)))\ \{0..<k\}$
by (*simp add*: *sum-distrib-left*[*symmetric*] *algebra-simps*)
also have ... $= sum\ ((\lambda i.\ (a\hat{}i * b\hat{}(k-i))) \circ (\lambda i.\ i+1))\ \{0..<k\} -$
$sum\ (\lambda i.\ (a\hat{}i * (b\hat{}(1+(k-1-i)))))\ \{0..<k\}$
by (*simp add*:*algebra-simps*)
also have ... $= sum\ ((\lambda i.\ (a\hat{}i * b\hat{}(k-i))) \circ (\lambda i.\ i+1))\ \{0..<k\} -$
$sum\ (\lambda i.\ (a\hat{}i * b\hat{}(k-i)))\ \{0..<k\}$
by (*intro arg-cong2*[**where** $f=(-)$] *sum.cong arg-cong2*[**where** $f=(*)$]
arg-cong2[**where** $f=(\lambda x\ y.\ x\ \hat{}\ y)$]) *auto*
also have ... $= sum\ (\lambda i.\ (a\hat{}i * b\hat{}(k-i)))\ (insert\ k\ \{1..<k\}) -$
$sum\ (\lambda i.\ (a\hat{}i * b\hat{}(k-i)))\ (insert\ 0\ \{Suc\ 0..<k\})$
using *assms*
by (*subst sum.reindex*[*symmetric*], *simp*, *subst insert-lb*, *auto*)
also have ... $= ?lhs$
by *simp*
finally show *?thesis* **by** *presburger*
qed

lemma *power-diff-est*:
assumes $k > 0$
assumes $(a :: real) \geq b$
assumes $b \geq 0$
shows $a\hat{}k - b\hat{}k \leq (a-b) * k * a\hat{}(k-1)$
proof $-$
have $\bigwedge i.\ i < k \Longrightarrow a\ \hat{}\ i * b\ \hat{}\ (k - 1 - i) \leq a\ \hat{}\ i * a\ \hat{}\ (k-1-i)$
using *assms* **by** (*intro mult-left-mono power-mono*) *auto*
also have $\bigwedge i.\ i < k \Longrightarrow a\ \hat{}\ i * a\ \hat{}\ (k - 1 - i) = a\ \hat{}\ (k - Suc\ 0)$
using *assms*(*1*) **by** (*subst power-add*[*symmetric*], *simp*)
finally have *a*: $\bigwedge i.\ i < k \Longrightarrow a\ \hat{}\ i * b\ \hat{}\ (k - 1 - i) \leq a\ \hat{}\ (k - Suc\ 0)$
by *blast*
have $a\hat{}k - b\hat{}k = (a-b) * (\sum i = 0..<k.\ a\ \hat{}\ i * b\ \hat{}\ (k - 1 - i))$
by (*rule power-diff-sum*[*OF assms*(*1*)])
also have ... $\leq (a-b) * (\sum i = 0..<k.\ \ a\hat{}(k-1))$
using *a assms* **by** (*intro mult-left-mono sum-mono*, *auto*)
also have ... $= (a-b) * (k * a\hat{}(k-Suc\ 0))$
by *simp*
finally show *?thesis* **by** *simp*
qed

Specialization of the Hoelder inquality for sums.

lemma *Holder-inequality-sum*:
  assumes $p > (0::real)$ $q > 0$ $1/p + 1/q = 1$
  assumes *finite A*
  shows $|\sum x{\in}A.\ f\ x * g\ x| \le (\sum x{\in}A.\ |f\ x|\ powr\ p)\ powr\ (1/p) * (\sum x{\in}A.\ |g\ x|\ powr\ q)\ powr\ (1/q)$
proof −
  have $|LINT\ x|count\text{-}space\ A.\ f\ x * g\ x| \le$
    $(LINT\ x|count\text{-}space\ A.\ |f\ x|\ powr\ p)\ powr\ (1\ /\ p) *$
    $(LINT\ x|count\text{-}space\ A.\ |g\ x|\ powr\ q)\ powr\ (1\ /\ q)$
    using *assms integrable-count-space*
    by (*intro Lp.Holder-inequality*, *auto*)
  thus *?thesis*
    using *assms* **by** (*simp add*: *lebesgue-integral-count-space-finite*[*symmetric*])
qed

lemma *real-count-list-pos*:
  assumes $x \in set\ as$
  shows $real\ (count\text{-}list\ as\ x) > 0$
  using *count-list-gr-1 assms* **by** *force*

lemma *fk-estimate*:
  assumes $as \neq []$
  shows $length\ as * of\text{-}rat\ (F\ (2{*}k{-}1)\ as) \le n\ powr\ (1 - 1\ /\ real\ k) * (of\text{-}rat\ (F\ k\ as))\hat{}2$
  (**is** $?lhs \le ?rhs$)
proof ($cases\ k \ge 2$)
  case *True*
  define *M* **where** $M = Max\ (count\text{-}list\ as\ `\ set\ as)$
  have $M \in count\text{-}list\ as\ `\ set\ as$
    unfolding *M-def* **using** *assms* **by** (*intro Max-in*, *auto*)
  then obtain *m* **where** *m-in*: $m \in set\ as$ **and** *m-def*: $M = count\text{-}list\ as\ m$
    by *blast*

  have *a*: $real\ M > 0$ **using** *m-in count-list-gr-1* **by** (*simp add*:*m-def*, *force*)
  have *b*: $2{*}k{-}1 = (k{-}1) + k$ **by** *simp*

  have $0 < real\ (count\text{-}list\ as\ m)$
    using *m-in count-list-gr-1* **by** *force*
  hence $M\ powr\ k = real\ (count\text{-}list\ as\ m)\ \hat{}\ k$
    by (*simp add*: *powr-realpow m-def*)
  also have ... $\le (\sum x{\in}set\ as.\ real\ (count\text{-}list\ as\ x)\ \hat{}\ k)$
    using *m-in* **by** (*intro member-le-sum*, *simp-all*)
  also have ... $\le real\text{-}of\text{-}rat\ (F\ k\ as)$
    by (*simp add*:*F-def of-rat-sum of-rat-power*)
  finally have *d*: $M\ powr\ k \le real\text{-}of\text{-}rat\ (F\ k\ as)$ **by** *simp*

  have *e*: $0 \le real\text{-}of\text{-}rat\ (F\ k\ as)$
    using *F-gr-0*[*OF assms*(*1*)] **by** (*simp add*: *order-le-less*)

have *real* (*k* − *1*) / *real k* + *1* = *real* (*k* − *1*) / *real k* + *real k* / *real k*
using *assms True* **by** *simp*
also have ... = *real* (*2* ∗ *k* − *1*) / *real k*
using *b* **by** (*subst add-divide-distrib*[*symmetric*], *force*)
finally have *f*: *real* (*k* − *1*) / *real k* + *1* = *real* (*2* ∗ *k* − *1*) / *real k*
by *blast*

have *real-of-rat* (*F* (*2*∗*k*−*1*) *as*) =
($\sum$ *x*∈*set as*. *real* (*count-list as x*) ^ (*k* − *1*) ∗ *real* (*count-list as x*) ^ *k*)
using *b* **by** (*simp add*:*F-def of-rat-sum sum-distrib-left of-rat-mult power-add of-rat-power*)
also have ... ≤ ($\sum$ *x*∈*set as*. *real M* ^ (*k* − *1*) ∗ *real* (*count-list as x*) ^ *k*)
by (*intro sum-mono mult-right-mono power-mono of-nat-mono*) (*auto simp*:*M-def*)
also have ... = *M powr* (*k*−*1*) ∗ *of-rat* (*F k as*) **using** *a*
by (*simp add*:*sum-distrib-left F-def of-rat-mult of-rat-sum of-rat-power powr-realpow*)
also have ... = (*M powr k*) *powr* (*real* (*k* − *1*) / *real k*) ∗ *of-rat* (*F k as*) *powr 1*
using *e* **by** (*simp add*:*powr-powr*)
also have ... ≤ (*real-of-rat* (*F k as*)) *powr* ((*k*−*1*)/*k*) ∗ (*real-of-rat* (*F k as*) *powr 1*)
using *d* **by** (*intro mult-right-mono powr-mono2*, *auto*)
also have ... = (*real-of-rat* (*F k as*)) *powr* ((*2*∗*k*−*1*) / *k*)
by (*subst powr-add*[*symmetric*], *subst f*, *simp*)
finally have *a*: *real-of-rat* (*F* (*2*∗*k*−*1*) *as*) ≤ (*real-of-rat* (*F k as*)) *powr* ((*2*∗*k*−*1*) / *k*)
by *blast*

have *g*: *card* (*set as*) ≤ *n*
using *card-mono*[*OF - as-range*] **by** *simp*

have *length as* = *abs* (*sum* (λ*x*. *real* (*count-list as x*)) (*set as*))
by (*subst of-nat-sum*[*symmetric*], *simp add*: *sum-count-set*)
also have ... ≤ *card* (*set as*) *powr* ((*real k* − *1*)/*k*) ∗
(*sum* (λ*x*. |*real* (*count-list as x*)| *powr k*) (*set as*)) *powr* (*1*/*k*)
using *assms True*
by (*intro Holder-inequality-sum*[**where** *p*=*k*/(*k*−*1*) **and** *q*=*k* **and** *f*=λ-.*1*, *simplified*])
(*auto simp add*:*algebra-simps add-divide-distrib*[*symmetric*])
also have ... = (*card* (*set as*)) *powr* ((*real k* − *1*) / *real k*) ∗ *of-rat* (*F k as*) *powr* (*1*/ *k*)
using *real-count-list-pos*
by (*simp add*:*F-def of-rat-sum of-rat-power powr-realpow*)
also have ... = (*card* (*set as*)) *powr* (*1* − *1* / *real k*) ∗ *of-rat* (*F k as*) *powr* (*1*/ *k*)
using *k-ge-1 assms True* **by** (*simp add*: *divide-simps*)
also have ... ≤ *n powr* (*1* − *1* / *real k*) ∗ *of-rat* (*F k as*) *powr* (*1*/ *k*)
using *k-ge-1 g*
by (*intro mult-right-mono powr-mono2*, *auto*)
finally have *h*: *length as* ≤ *n powr* (*1* − *1* / *real k*) ∗ *of-rat* (*F k as*) *powr* (*1*/*real k*)

by *blast*

have i:*1* / *real k* + *real* (*2* ∗ *k* − *1*) / *real k* = *real 2*
using *True* **by** (*subst add-divide-distrib*[*symmetric*], *simp-all add*:*of-nat-diff*)

have *?lhs* ≤ *n powr* (*1* − *1*/*k*) ∗ *of-rat* (*F k as*) *powr* (*1*/*k*) ∗ (*of-rat* (*F k as*)) *powr* ((*2*∗*k*−*1*) / *k*)
using *a h F-ge-0* **by** (*intro mult-mono mult-nonneg-nonneg*, *auto*)
also have ... = *?rhs*
using *i F-gr-0*[*OF assms*] **by** (*simp add*:*powr-add*[*symmetric*] *powr-realpow*[*symmetric*])
finally show *?thesis*
by *blast*
next
case *False*
have *n* = *0* ⟹ *False*
using *as-range assms* **by** *auto*
hence *n* > *0*
by *auto*
moreover have *k* = *1*
using *assms k-ge-1 False* **by** *linarith*
moreover have *length as* = *real-of-rat* (*F* (*Suc 0*) *as*)
by (*simp add*:*F-def sum-count-set of-nat-sum*[*symmetric*] *del*:*of-nat-sum*)
ultimately show *?thesis*
by (*simp add*:*power2-eq-square*)
qed

definition *result*
where *result a* = *of-nat* (*length as*) ∗ *of-nat* (*Suc* (*snd a*) ^ *k* − *snd a* ^ *k*)

lemma *result-exp-1*:
assumes *as* ≠ []
shows *expectation result* = *real-of-rat* (*F k as*)
proof −
have *expectation result* = ($\sum$ *a*∈M_1. *result a* ∗ *pmf* (*pmf-of-set* M_1) *a*)
unfolding Ω_1-*def* **using** *non-empty-space assms fin-space*
by (*subst integral-measure-pmf-real*) *auto*
also have ... = ($\sum$ *a*∈M_1. *result a* / *real* (*length as*))
using *non-empty-space assms fin-space card-space* **by** *simp*
also have ... = ($\sum$ *a*∈M_1. *real* (*Suc* (*snd a*) ^ *k* − *snd a* ^ *k*))
using *assms* **by** (*simp add*:*result-def*)
also have ... = ($\sum$ *u*∈*set as*. $\sum$ *v* = *0*..<*count-list as u*. *real* (*Suc v* ^ *k*) − *real* (*v* ^ *k*))
using *k-ge-1* **by** (*subst split-space*, *simp add*:*of-nat-diff*)
also have ... = ($\sum$ *u*∈*set as*. *real* (*count-list as u*)^*k*)
using *k-ge-1* **by** (*subst sum-Suc-diff'*) (*auto simp add*:*zero-power*)
also have ... = *of-rat* (*F k as*)
by (*simp add*:*F-def of-rat-sum of-rat-power*)
finally show *?thesis* **by** *simp*
qed

```
lemma result-var-1:
  assumes as ≠ []
  shows variance result ≤ (of-rat (F k as))^2 ∗ k ∗ n powr (1 − 1 / real k)
proof −
  have k-gt-0: k > 0 using k-ge-1 by linarith

  have c:real (Suc v ^ k) − real (v ^ k) ≤ k ∗ real (count-list as a) ^(k − Suc 0)
    if c-1: v < count-list as a for a v
  proof −
    have real (Suc v ^ k) − real (v ^ k) ≤ (real (v+1) − real v) ∗ k ∗ (1 + real
v) ^(k − Suc 0)
      using k-gt-0 power-diff-est[where a=Suc v and b=v] by simp
    moreover have (real (v+1) − real v) = 1 by auto
    ultimately have real (Suc v ^ k) − real (v ^ k) ≤ k ∗ (1 + real v) ^(k −
Suc 0)
      by auto
    also have ... ≤ k ∗ real (count-list as a) ^(k− Suc 0)
      using c-1 by (intro mult-left-mono power-mono, auto)
    finally show ?thesis by blast
  qed

  have length as ∗ (∑ a∈ M_1. (real (Suc (snd a)  ^ k − (snd a) ^ k))^2) =
    length as ∗ (∑ a∈ set as. (∑ v ∈ {0..<count-list as a}.
    real (Suc v ^ k − v ^ k) ∗ real (Suc v ^ k − v^k)))
    by (subst split-space, simp add:power2-eq-square)
  also have ... ≤ length as ∗ (∑ a∈ set as. (∑ v ∈ {0..<count-list as a}.
    k ∗ real (count-list as a) ^(k−1) ∗ real (Suc v ^ k − v ^ k)))
   using c by (intro mult-left-mono sum-mono mult-right-mono) (auto simp:power-mono
of-nat-diff)
  also have ... = length as ∗ k ∗ (∑ a∈ set as. real (count-list as a) ^(k−1) ∗
    (∑ v ∈ {0..<count-list as a}. real (Suc v ^ k) − real (v ^ k)))
    by (simp add:sum-distrib-left ac-simps of-nat-diff power-mono)
  also have ... = length as ∗ k ∗ (∑ a∈ set as. real (count-list as a ^(2∗k−1)))
    using assms k-ge-1
   by (subst sum-Suc-diff', auto simp: zero-power[OF k-gt-0] mult-2 power-add[symmetric])
  also have ... = k ∗ (length as ∗ of-rat (F (2∗k−1) as))
    by (simp add:sum-distrib-left[symmetric] F-def of-rat-sum of-rat-power)
  also have ... ≤ k ∗ (of-rat (F k as)^2 ∗ n powr (1 − 1 / real k))
   using fk-estimate[OF assms] by (intro mult-left-mono) (auto simp: mult.commute)
  finally have b: real (length as) ∗ (∑ a∈ M_1. (real (Suc (snd a) ^ k − (snd a) ^
k))^2) ≤
    k ∗ ((of-rat (F k as))^2 ∗ n powr (1 − 1 / real k))
    by blast

 have expectation (λω. (result ω :: real)^2) − (expectation result)^2 ≤ expectation
(λω. result ω^2)
    by simp
  also have ... = (∑ a∈M_1. (length as ∗ real (Suc (snd a) ^ k − snd a ^ k))^2 ∗
```

pmf (*pmf-of-set* M_1) *a*)
  using *fin-space non-empty-space assms* **unfolding** Ω_1*-def result-def*
  by (*subst integral-measure-pmf-real*[**where** A=M_1], *auto*)
 also have ... = ($\sum$ $a \in M_1$. *length as* $*$ (*real* (*Suc* (*snd a*) ^ *k* − *snd a* ^ *k*))^2)
  using *assms non-empty-space fin-space* **by** (*subst pmf-of-set*)
   (*simp-all add*:*card-space power-mult-distrib power2-eq-square ac-simps*)
 also have ... $\leq$ *k* $*$ ((*of-rat* (*F k as*))^2 $*$ *n powr* (*1* − *1* / *real k*))
  using *b* **by** (*simp add*:*sum-distrib-left*[*symmetric*])
 also have ... = *of-rat* (*F k as*)^*2* $*$ *k* $*$ *n powr* (*1* − *1* / *real k*)
  by (*simp add*:*ac-simps*)
 finally have *expectation* ($\lambda\omega$. *result* ω^*2*) − (*expectation result*)^*2* $\leq$
  of-rat (*F k as*)^*2* $*$ *k* $*$ *n powr* (*1* − *1* / *real k*)
  by *blast*

 thus *?thesis*
  using *integrable-1*[*OF assms*] **by** (*simp add*:*variance-eq*)
qed

theorem *fk-alg-sketch*:
 assumes *as* $\neq$ []
 shows *fold* (λ*a state*. *state* $\gg=$ *fk-update a*) *as* (*fk-init k* δ ε *n*) =
  map-pmf (λ*x*. (s_1,s_2,*k*,*length as*, *x*)) M_2 (**is** *?lhs* = *?rhs*)
proof −
 have *?lhs* = *prod-pmf* ({*0*..<s_1} × {*0*..<s_2})
  (λ-. *fold* (λ*x s*. *s* $\gg=$ *fk-update-2 x*) *as* (*return-pmf* (*0*, *0*, *0*))) $\gg=$
  (λ*x*. *return-pmf* (s_1, s_2, *k*, *length as*, $\lambda i \in$ {*0*..<s_1} × {*0*..<s_2}. *snd* (*x i*)))
  by (*subst fk-update-distr*, *simp*)
 also have ... = *prod-pmf* ({*0*..<s_1} × {*0*..<s_2}) (λ-. *pmf-of-set* {..<*length as*}
$\gg=$
  (λ*k*. *return-pmf* (*length as*, *sketch as k*))) $\gg=$
  (λ*x*. *return-pmf* (s_1, s_2, *k*, *length as*, $\lambda i \in$ {*0*..<s_1} × {*0*..<s_2}. *snd* (*x i*)))
  by (*subst fk-update-2-distr*[*OF assms*], *simp*)
 also have ... = *prod-pmf* ({*0*..<s_1} × {*0*..<s_2}) (λ-. *pmf-of-set* {..<*length as*}
$\gg=$
  (λ*k*. *return-pmf* (*sketch as k*)) $\gg=$ (λ*s*. *return-pmf* (*length as*, *s*))) $\gg=$
  (λ*x*. *return-pmf* (s_1, s_2, *k*, *length as*, $\lambda i \in$ {*0*..<s_1} × {*0*..<s_2}. *snd* (*x i*)))
  by (*subst bind-assoc-pmf*, *subst bind-return-pmf*, *simp*)
 also have ... = *prod-pmf* ({*0*..<s_1} × {*0*..<s_2}) (λ-. *pmf-of-set* {..<*length as*}
$\gg=$
  (λ*k*. *return-pmf* (*sketch as k*))) $\gg=$
  (λ*x*. *return-pmf* ($\lambda i \in$ {*0*..<s_1} × {*0*..<s_2}. (*length as*, *x i*))) $\gg=$
  (λ*x*. *return-pmf* (s_1, s_2, *k*, *length as*, $\lambda i \in$ {*0*..<s_1} × {*0*..<s_2}. *snd* (*x i*)))
  by (*subst Pi-pmf-bind-return*[**where** *d'*=*undefined*], *simp*, *simp add*:*restrict-def*)
 also have ... = *prod-pmf* ({*0*..<s_1} × {*0*..<s_2}) (λ-. *pmf-of-set* {..<*length as*}
$\gg=$
  (λ*k*. *return-pmf* (*sketch as k*))) $\gg=$
  (λ*x*. *return-pmf* (s_1, s_2, *k*, *length as*, *restrict x* ({*0*..<s_1} × {*0*..<s_2})))
  by (*subst bind-assoc-pmf*, *simp add*:*bind-return-pmf cong*:*restrict-cong*)
 also have ... = M_2 $\gg=$

$(\lambda x.\ \textit{return-pmf}\ (s_1,\ s_2,\ k,\ \textit{length as},\ \textit{restrict}\ x\ (\{0..<s_1\} \times \{0..<s_2\})))$
by $(\textit{subst sketch-distr}[\textit{OF assms}],\ \textit{simp add}{:}M_2\textit{-def})$
also have ... $= M_2 \gg\!= (\lambda x.\ \textit{return-pmf}\ (s_1,\ s_2,\ k,\ \textit{length as},\ x))$
by $(\textit{rule bind-pmf-cong},\ \textit{auto simp add}{:}\textit{PiE-dflt-def}\ M_2\textit{-def set-Pi-pmf})$
also have ... $=$ *?rhs*
by (*simp add*:*map-pmf-def*)
finally show *?thesis* **by** *simp*
qed

definition *mean-rv*
where $\textit{mean-rv}\ \omega\ i_2 = (\sum i_1 = 0..<s_1.\ \textit{result}\ (\omega\ (i_1,\ i_2)))\ /\ \textit{of-nat}\ s_1$

definition *median-rv*
where $\textit{median-rv}\ \omega = \textit{median}\ s_2\ (\lambda i_2.\ \textit{mean-rv}\ \omega\ i_2)$

lemma *fk-alg-correct′*:
defines $M \equiv \textit{fold}\ (\lambda a\ \textit{state}.\ \textit{state} \gg\!= \textit{fk-update}\ a)\ \textit{as}\ (\textit{fk-init}\ k\ \delta\ \varepsilon\ n) \gg\!= \textit{fk-result}$
shows $\mathcal{P}(\omega\ \textit{in measure-pmf}\ M.\ |\omega - F\ k\ \textit{as}| \le \delta * F\ k\ \textit{as}) \ge 1 - \textit{of-rat}\ \varepsilon$
proof $(\textit{cases as} = [])$
case *True*
have a: $\textit{nat}\ \lceil -\ (18 * \ln\ (\textit{real-of-rat}\ \varepsilon))\rceil > 0$ **using** ε*-range* **by** *simp*
show *?thesis* **using** *True* ε*-range*
by (*simp add*:*F-def M-def bind-return-pmf median-const*[*OF a*] *Let-def*)
next
case *False*

have $\textit{set as} \neq \{\}$ **using** *assms False* **by** *blast*
hence *n-nonzero*: $n > 0$ **using** *as-range* **by** *fastforce*

have *fk-nonzero*: $F\ k\ \textit{as} > 0$
using *F-gr-0*[*OF False*] **by** *simp*

have *s1-nonzero*: $s_1 > 0$
using δ*-range k-ge-1 n-nonzero* **by** $(\textit{simp add}{:}s_1\textit{-def})$
have *s2-nonzero*: $s_2 > 0$
using ε*-range* **by** $(\textit{simp add}{:}s_2\textit{-def})$

have *real-of-rat-mean-rv*: $\bigwedge x\ i.\ \textit{mean-rv}\ x = (\lambda i.\ \textit{real-of-rat}\ (\textit{mean-rv}\ x\ i))$
by (*rule ext*, *simp add*:*of-rat-divide of-rat-sum of-rat-mult result-def mean-rv-def*)
have *real-of-rat-median-rv*: $\bigwedge x.\ \textit{median-rv}\ x = \textit{real-of-rat}\ (\textit{median-rv}\ x)$
unfolding *median-rv-def* **using** *s2-nonzero*
by (*subst real-of-rat-mean-rv*, *simp add*: *median-rat median-restrict*)

have *space*-Ω_2: $\textit{space}\ \Omega_2 = \textit{UNIV}$ **by** $(\textit{simp add}{:}\Omega_2\textit{-def})$

have *fk-result-eta*: $\textit{fk-result} = (\lambda(x,y,z,u,v).\ \textit{fk-result}\ (x,y,z,u,v))$
by *auto*

have *a*:*fold* (λx *state*. *state* $\gg=$ *fk-update* x) *as* (*fk-init* k δ ε n) =
map-pmf (λx. (s_1,s_2,k,*length as*, x)) M_2
by (*subst fk-alg-sketch*[*OF False*]) (*simp add*:s_1-*def*[*symmetric*] s_2-*def*[*symmetric*])

have M = *map-pmf* (λx. (s_1, s_2, k, *length as*, x)) M_2 $\gg=$ *fk-result*
by (*subst M-def*, *subst a*, *simp*)
also have ... = M_2 $\gg=$ *return-pmf* $\circ$ *median-rv*
by (*subst fk-result-eta*)
(*auto simp add*:*map-pmf-def bind-assoc-pmf bind-return-pmf median-rv-def mean-rv-def comp-def*
M_1-*def result-def median-restrict*)
finally have *b*: M = M_2 $\gg=$ *return-pmf* $\circ$ *median-rv*
by *simp*

have *result-exp*:
$i_1 < s_1 \Longrightarrow i_2 < s_2 \Longrightarrow \Omega_2$.*expectation* ($\lambda x$. *result* ($x$ (i_1, i_2))) = *real-of-rat* (*F* *k as*)
for i_1 i_2
unfolding Ω_2-*def* M_2-*def*
using *integrable-1*[*OF False*] *result-exp-1*[*OF False*]
by (*subst expectation-Pi-pmf-slice*, *auto simp*:Ω_1-*def*)

have *result-var*: Ω_2.*variance* ($\lambda\omega$. *result* (ω (i_1, i_2))) $\le$ *of-rat* (δ $*$ *F k as*)^*2* $*$ *real* s_1 / *3*
if *result-var-assms*: $i_1 < s_1$ $i_2 < s_2$ **for** i_1 i_2
proof −
have *3* $*$ *real k* $*$ *n powr* (*1* − *1* / *real k*) =
(*of-rat* δ)^2 $*$ (*3* $*$ *real k* $*$ *n powr* (*1* − *1* / *real k*) / (*of-rat* δ)^2)
using δ-*range* **by** *simp*
also have ... $\le$ (*real-of-rat* δ)^2 $*$ (*real* s_1)
unfolding s_1-*def*
by (*intro mult-mono of-nat-ceiling*, *simp-all*)
finally have *f2-var-2*: *3* $*$ *real k* $*$ *n powr* (*1* − *1* / *real k*) $\le$ (*of-rat* δ)^2 $*$ (*real* s_1)
by *blast*

have Ω_2.*variance* ($\lambda\omega$. *result* (ω (i_1, i_2)) :: *real*) = *variance result*
using *result-var-assms integrable-1*[*OF False*]
unfolding Ω_2-*def* M_2-*def* Ω_1-*def*
by (*subst variance-prod-pmf-slice*, *auto*)
also have ... $\le$ *of-rat* (*F k as*)^*2* $*$ *real k* $*$ *n powr* (*1* − *1* / *real k*)
using *assms False result-var-1* Ω_1-*def* **by** *simp*
also have ... =
of-rat (*F k as*)^*2* $*$ (*real k* $*$ *n powr* (*1* − *1* / *real k*))
by (*simp add*:*ac-simps*)
also have ... $\le$ *of-rat* (*F k as*)^*2* $*$ (*of-rat* δ^*2* $*$ (*real* s_1 / *3*))
using *f2-var-2* **by** (*intro mult-left-mono*, *auto*)
also have ... = *of-rat* (*F k as* $*$ δ)^*2* $*$ (*real* s_1 / *3*)

by (*simp add*: *of-rat-mult power-mult-distrib*)
also have ... = *of-rat* (δ ∗ *F k as*)^*2* ∗ *real* s_1 / *3*
by (*simp add*:*ac-simps*)
finally show *?thesis*
by *simp*
qed

have *mean-rv-exp*: Ω_2.*expectation* ($\lambda\omega$. *mean-rv* ω *i*) = *real-of-rat* (*F k as*)
if *mean-rv-exp-assms*: *i* < s_2 **for** *i*
proof –
have Ω_2.*expectation* ($\lambda\omega$. *mean-rv* ω *i*) = Ω_2.*expectation* ($\lambda\omega$. $\sum$ *n* = *0*..<s_1. *result* (ω (*n*, *i*)) / *real* s_1)
by (*simp add*:*mean-rv-def sum-divide-distrib*)
also have ... = ($\sum$ *n* = *0*..<s_1. Ω_2.*expectation* ($\lambda\omega$. *result* (ω (*n*, *i*))) / *real* s_1)
using *integrable-2*[*OF False*]
by (*subst Bochner-Integration.integral-sum*, *auto*)
also have ... = *of-rat* (*F k as*)
using *s1-nonzero mean-rv-exp-assms*
by (*simp add*:*result-exp*)
finally show *?thesis* **by** *simp*
qed

have *mean-rv-var*: Ω_2.*variance* ($\lambda\omega$. *mean-rv* ω *i*) ≤ *real-of-rat* (δ ∗ *F k as*)^*2*/*3*
if *mean-rv-var-assms*: *i* < s_2 **for** *i*
proof –
have *a*:Ω_2.*indep-vars* (λ-. *borel*) (λ*n x*. *result* (*x* (*n*, *i*)) / *real* s_1) {*0*..<s_1}
unfolding Ω_2-*def* M_2-*def* **using** *mean-rv-var-assms*
by (*intro indep-vars-restrict-intro′*[**where** *f*=*fst*], *simp*, *simp add*:*restrict-dfl-def*, *simp*, *simp*)
have Ω_2.*variance* ($\lambda\omega$. *mean-rv* ω *i*) = Ω_2.*variance* ($\lambda\omega$. $\sum$ *j* = *0*..<s_1. *result* (ω (*j*, *i*)) / *real* s_1)
by (*simp add*:*mean-rv-def sum-divide-distrib*)
also have ... = ($\sum$ *j* = *0*..<s_1. Ω_2.*variance* ($\lambda\omega$. *result* (ω (*j*, *i*)) / *real* s_1))
using *a integrable-2*[*OF False*]
by (*subst* Ω_2.*bienaymes-identity-full-indep*, *auto simp add*:Ω_2-*def*)
also have ... = ($\sum$ *j* = *0*..<s_1. Ω_2.*variance* ($\lambda\omega$. *result* (ω (*j*, *i*))) / *real* s_1^*2*)
using *integrable-2*[*OF False*]
by (*subst* Ω_2.*variance-divide*, *auto*)
also have ... ≤ ($\sum$ *j* = *0*..<s_1. ((*real-of-rat* (δ ∗ *F k as*))^2 ∗ *real* s_1 / *3*) / (*real* s_1^*2*))
using *result-var*[*OF - mean-rv-var-assms*]
by (*intro sum-mono divide-right-mono*, *auto*)
also have ... = *real-of-rat* (δ ∗ *F k as*)^*2*/*3*
using *s1-nonzero*
by (*simp add*:*algebra-simps power2-eq-square*)
finally show *?thesis* **by** *simp*
qed

have Ω_2.*prob* {*y*. *of-rat* (δ ∗ *F k as*) < |*mean-rv y i* − *real-of-rat* (*F k as*)|} ≤

1/3
  (**is** *?lhs* ≤ -) **if** *c-assms*: $i < s_2$ **for** *i*
proof −
  define *a* **where** *a* = *real-of-rat* (δ ∗ *F k as*)
  have *c*: *0* < *a* **unfolding** *a-def*
    using *assms* δ*-range fk-nonzero*
    by (*metis zero-less-of-rat-iff mult-pos-pos*)
  have *?lhs* ≤ Ω_2*.prob* {*y* ∈ *space* Ω_2. *a* ≤ |*mean-rv y i* − Ω_2*.expectation* (λω. *mean-rv* ω *i*)|}
    by (*intro* Ω_2*.pmf-mono*[*OF* Ω_2*-def*], *simp add*:*a-def mean-rv-exp*[*OF c-assms*] *space-*Ω_2)
  also have ... ≤ Ω_2*.variance* (λω. *mean-rv* ω *i*)/*a*^*2*
    by (*intro* Ω_2*.Chebyshev-inequality integrable-2 c False*) (*simp add*:Ω_2*-def*)
  also have ... ≤ *1/3* **using** *c*
    using *mean-rv-var*[*OF c-assms*]
    by (*simp add*:*algebra-simps*, *simp add*:*a-def*)
  finally show *?thesis*
    by *blast*
qed

moreover have Ω_2*.indep-vars* (λ-. *borel*) (λ*i* ω. *mean-rv* ω *i*) {*0*..<s_2}
  using *s1-nonzero* **unfolding** Ω_2*-def* M_2*-def*
  by (*intro indep-vars-restrict-intro'*[**where** *f*=*snd*] *finite-cartesian-product*)
    (*simp-all add*:*mean-rv-def restrict-dfl-def space-*Ω_2)
moreover have − (*18* ∗ *ln* (*real-of-rat* ε)) ≤ *real* s_2
  by (*simp add*:s_2*-def*, *linarith*)
ultimately have *1* − *of-rat* ε ≤
  Ω_2*.prob* {*y* ∈ *space* Ω_2. |*median* s_2 (*mean-rv y*) − *real-of-rat* (*F k as*)| ≤ *of-rat* (δ ∗ *F k as*)}
  using ε*-range*
  by (*intro* Ω_2*.median-bound-2*, *simp-all add*:*space-*Ω_2)
also have ... = Ω_2*.prob* {*y*. |*median-rv y* − *real-of-rat* (*F k as*)| ≤ *real-of-rat* (δ ∗ *F k as*)}
  by (*simp add*:*median-rv-def space-*Ω_2)
also have ... = Ω_2*.prob* {*y*. |*median-rv y* − *F k as*| ≤ δ ∗ *F k as*}
  by (*simp add*:*real-of-rat-median-rv of-rat-less-eq flip*: *of-rat-diff*)
also have ... = $\mathcal{P}$(ω *in measure-pmf M*. |ω − *F k as*| ≤ δ ∗ *F k as*)
  by (*simp add*: *b comp-def map-pmf-def*[*symmetric*] Ω_2*-def*)
finally show *?thesis* **by** *simp*
qed

lemma *fk-exact-space-usage'*:
  defines *M* ≡ *fold* (λ*a state*. *state* ⤜ *fk-update a*) *as* (*fk-init k* δ ε *n*)
  shows *AE* ω *in M*. *bit-count* (*encode-fk-state* ω) ≤ *fk-space-usage* (*k*, *n*, *length as*, ε, δ)
    (**is** *AE* ω *in M*. (- ≤ *?rhs*))
proof −
  define *H* **where** *H* = (*if as* = [] *then return-pmf* (λ*i*∈ {*0*..<s_1}×{*0*..<s_2}. (*0*,*0*)) *else* M_2)

have a:$M = map\text{-}pmf\ (\lambda x.(s_1,s_2,k,length\ as,\ x))\ H$
proof $(cases\ as \neq [])$
case *True*
then show *?thesis*
unfolding *M-def fk-alg-sketch*[*OF True*] *H-def*
by $(simp\ add{:}M_2\text{-}def)$
next
case *False*
then show *?thesis*
by $(simp\ add{:}H\text{-}def\ M\text{-}def\ s_1\text{-}def[symmetric]\ Let\text{-}def\ s_2\text{-}def[symmetric]\ map\text{-}pmf\text{-}def\ bind\text{-}return\text{-}pmf)$
qed

have $bit\text{-}count\ (encode\text{-}fk\text{-}state\ (s_1,\ s_2,\ k,\ length\ as,\ y)) \leq\ ?rhs$
if b:$y \in set\text{-}pmf\ H$ **for** y
proof −
have $b0$: $as \neq [] \Longrightarrow y \in \{0..<s_1\} \times \{0..<s_2\} \rightarrow_E M_1$
using *b non-empty-space fin-space* **by** $(simp\ add{:}H\text{-}def\ M_2\text{-}def\ set\text{-}prod\text{-}pmf)$

have $bit\text{-}count\ ((N_e \times_e N_e)\ (y\ x)) \leq$
$ereal\ (2 * log\ 2\ (real\ n + 1) + 1) + ereal\ (2 * log\ 2\ (real\ (length\ as) + 1) + 1)$
(**is** - $\leq$ *?rhs1*)
if $b1\text{-}assms$: $x \in \{0..<s_1\}\times\{0..<s_2\}$ **for** x
proof −
have $fst\ (y\ x) \leq n$
proof $(cases\ as = [])$
case *True*
then show *?thesis* **using** *b b1-assms* **by** $(simp\ add{:}H\text{-}def)$
next
case *False*
hence $1 \leq count\text{-}list\ as\ (fst\ (y\ x))$
using *b0 b1-assms* **by** $(simp\ add{:}PiE\text{-}iff\ case\text{-}prod\text{-}beta\ M_1\text{-}def,\ fastforce)$
hence $fst\ (y\ x) \in set\ as$
using *count-list-gr-1* **by** *metis*
then show *?thesis*
by (*meson lessThan-iff less-imp-le-nat subsetD as-range*)
qed
moreover have $snd\ (y\ x) \leq length\ as$
proof $(cases\ as = [])$
case *True*
then show *?thesis* **using** *b b1-assms* **by** $(simp\ add{:}H\text{-}def)$
next
case *False*
hence $(y\ x) \in M_1$
using *b0 b1-assms* **by** *auto*
hence $snd\ (y\ x) \leq count\text{-}list\ as\ (fst\ (y\ x))$
by $(simp\ add{:}M_1\text{-}def\ case\text{-}prod\text{-}beta)$

then show *?thesis* **using** *count-le-length* **by** (*metis order-trans*)
qed
ultimately have *bit-count* (N_e (*fst* (*y x*))) + *bit-count* (N_e (*snd* (*y x*))) ≤ *?rhs1*
using *exp-golomb-bit-count-est* **by** (*intro add-mono*, *auto*)
thus *?thesis*
by (*subst dependent-bit-count-2*, *simp*)
qed

moreover have *y* ∈ *extensional* ({*0..<*s_1} × {*0..<*s_2})
using *b0 b PiE-iff* **by** (*cases as* = [], *auto simp*:*H-def PiE-iff*)

ultimately have *bit-count* ((*List.product* [*0..<*s_1] [*0..<*s_2] $\rightarrow_e$ N_e $\times_e$ N_e) *y*) ≤
ereal (*real* s_1 ∗ *real* s_2) ∗ (*ereal* (*2* ∗ *log 2* (*real n* + *1*) + *1*) +
ereal (*2* ∗ *log 2* (*real* (*length as*) + *1*) + *1*))
by (*intro fun-bit-count-est*[**where** *xs*=(*List.product* [*0..<*s_1] [*0..<*s_2]), *simplified*], *auto*)
hence *bit-count* (*encode-fk-state* (s_1, s_2, *k*, *length as*, *y*)) ≤
ereal (*2* ∗ *log 2* (*real* s_1 + *1*) + *1*) +
(*ereal* (*2* ∗ *log 2* (*real* s_2 + *1*) + *1*) +
(*ereal* (*2* ∗ *log 2* (*real k* + *1*) + *1*) +
(*ereal* (*2* ∗ *log 2* (*real* (*length as*) + *1*) + *1*) +
(*ereal* (*real* s_1 ∗ *real* s_2) ∗ (*ereal* (*2* ∗ *log 2* (*real n*+*1*) + *1*) +
ereal (*2* ∗ *log 2* (*real* (*length as*)+*1*) + *1*))))))
unfolding *encode-fk-state-def dependent-bit-count*
by (*intro add-mono exp-golomb-bit-count*, *auto*)
also have ... ≤ *?rhs*
by (*simp add*: s_1-*def*[*symmetric*] s_2-*def*[*symmetric*] *Let-def*) (*simp add*:*ac-simps*)
finally show *bit-count* (*encode-fk-state* (s_1, s_2, *k*, *length as*, *y*)) ≤ *?rhs*
by *blast*
qed
thus *?thesis*
by (*simp add*: *a AE-measure-pmf-iff del*:*fk-space-usage.simps*)
qed

end

Main results of this section:

theorem *fk-alg-correct*:
assumes *k* ≥ *1*
assumes ε ∈ {*0<..<1*}
assumes δ > *0*
assumes *set as* ⊆ {*..<n*}
defines *M* ≡ *fold* (λ*a state*. *state* ≫= *fk-update a*) *as* (*fk-init k* δ ε *n*) ≫= *fk-result*
shows $\mathcal{P}$(ω *in measure-pmf M*. |ω − *F k as*| ≤ δ ∗ *F k as*) ≥ *1* − *of-rat* ε
unfolding *M-def* **using** *fk-alg-correct*′[*OF assms*(*1*−*4*)] **by** *blast*

theorem *fk-exact-space-usage*:

assumes $k \geq 1$
assumes $\varepsilon \in \{0<..<1\}$
assumes $\delta > 0$
assumes *set as* $\subseteq$ {..<*n*}
defines *M* ≡ *fold* (λ*a state. state* ⤜ *fk-update a*) *as* (*fk-init k* δ ε *n*)
shows *AE* ω *in M. bit-count* (*encode-fk-state* ω) ≤ *fk-space-usage* (*k*, *n*, *length as*, ε, δ)
unfolding *M-def* **using** *fk-exact-space-usage'*[*OF assms*(*1*−*4*)] **by** *blast*

theorem *fk-asymptotic-space-complexity*:
fk-space-usage ∈
O[*at-top* $\times_F$ *at-top* $\times_F$ *at-top* $\times_F$ *at-right* (*0*::*rat*) $\times_F$ *at-right* (*0*::*rat*)](λ (*k*, *n*, *m*, ε, δ).
real k ∗ *real n powr* (*1*−*1*/ *real k*) / (*of-rat* δ)^2 ∗ (*ln* (*1* / *of-rat* ε)) ∗ (*ln* (*real n*) + *ln* (*real m*)))
(**is** - ∈ *O*[*?F*](*?rhs*))
proof −
define *k-of* :: *nat* × *nat* × *nat* × *rat* × *rat* ⇒ *nat* **where** *k-of* = (λ(*k*, *n*, *m*, ε, δ). *k*)
define *n-of* :: *nat* × *nat* × *nat* × *rat* × *rat* ⇒ *nat* **where** *n-of* = (λ(*k*, *n*, *m*, ε, δ). *n*)
define *m-of* :: *nat* × *nat* × *nat* × *rat* × *rat* ⇒ *nat* **where** *m-of* = (λ(*k*, *n*, *m*, ε, δ). *m*)
define ε*-of* :: *nat* × *nat* × *nat* × *rat* × *rat* ⇒ *rat* **where** ε*-of* = (λ(*k*, *n*, *m*, ε, δ). ε)
define δ*-of* :: *nat* × *nat* × *nat* × *rat* × *rat* ⇒ *rat* **where** δ*-of* = (λ(*k*, *n*, *m*, ε, δ). δ)

define *g1* **where**
g1 = (λ*x*. *real* (*k-of x*)∗(*real* (*n-of x*)) *powr* (*1*−*1*/ *real* (*k-of x*)) ∗ (*1* / *of-rat* (δ*-of x*)^*2*))

define *g* **where**
g = (λ*x*. *g1 x* ∗ (*ln* (*1* / *of-rat* (ε*-of x*))) ∗ (*ln* (*real* (*n-of x*)) + *ln* (*real* (*m-of x*))))

define *s1-of* **where** *s1-of* = (λ*x*.
nat ⌈*3* ∗ *real* (*k-of x*) ∗ *real* (*n-of x*) *powr* (*1* − *1* / *real* (*k-of x*)) / (*real-of-rat* (δ*-of x*))^2⌉)
define *s2-of* **where** *s2-of* = (λ*x*. *nat* ⌈− (*18* ∗ *ln* (*real-of-rat* (ε*-of x*)))⌉)

have *evt*: (⋀*x*.
0 < *real-of-rat* (δ*-of x*) ∧ *0* < *real-of-rat* (ε*-of x*) ∧
1/*real-of-rat* (δ*-of x*) ≥ δ ∧ *1*/*real-of-rat* (ε*-of x*) ≥ ε ∧
real (*n-of x*) ≥ *n* ∧ *real* (*k-of x*) ≥ *k* ∧ *real* (*m-of x*) ≥ *m*⟹ *P x*)
⟹ *eventually P ?F* (**is** (⋀*x*. *?prem x* ⟹ -) ⟹ -)
for δ ε *n k m P*
apply (*rule eventually-mono*[**where** *P*=*?prem* **and** *Q*=*P*])
apply (*simp add*:ε*-of-def case-prod-beta'* δ*-of-def n-of-def k-of-def m-of-def*)

apply (*intro eventually-conj eventually-prod1′ eventually-prod2′ sequentially-inf eventually-at-right-less inv-at-right-0-inf*)
by (*auto simp add:prod-filter-eq-bot*)

have *1*:
$(\lambda\text{-}.\ 1) \in O[?F](\lambda x.\ real\ (n\text{-}of\ x))$
$(\lambda\text{-}.\ 1) \in O[?F](\lambda x.\ real\ (m\text{-}of\ x))$
$(\lambda\text{-}.\ 1) \in O[?F](\lambda x.\ real\ (k\text{-}of\ x))$
using *landau-o.big-mono eventually-mono*[*OF evt*]
by (*smt* (*verit*, *del-insts*) *real-norm-def*)+
have $(\lambda x.\ ln\ (real\ (m\text{-}of\ x) + 1)) \in O[?F](\lambda x.\ ln\ (real\ (m\text{-}of\ x)))$
by (*intro landau-ln-2*[**where** *a=2*] *evt*[**where** *m=2*] *sum-in-bigo 1*, *auto*)
hence *2*: $(\lambda x.\ log\ 2\ (real\ (m\text{-}of\ x) + 1)) \in O[?F](\lambda x.\ ln\ (real\ (n\text{-}of\ x)) + ln\ (real\ (m\text{-}of\ x)))$
by (*intro landau-sum-2 eventually-mono*[*OF evt*[**where** *n=1* **and** *m=1*]])
(*auto simp add:log-def*)

have *3*: $(\lambda\text{-}.\ 1) \in O[?F](\lambda x.\ ln\ (1\ /\ real\text{-}of\text{-}rat\ (\varepsilon\text{-}of\ x)))$
using *order-less-le-trans*[*OF exp-gt-zero*] *ln-ge-iff*
by (*intro landau-o.big-mono evt*[**where** ε*=exp 1*])
(*simp add*: *abs-ge-iff*, *blast*)

have *4*: $(\lambda\text{-}.\ 1) \in O[?F](\lambda x.\ 1\ /\ (real\text{-}of\text{-}rat\ (\delta\text{-}of\ x))^2)$
using *one-le-power*
by (*intro landau-o.big-mono evt*[**where** δ*=1*])
(*simp add:power-one-over*[*symmetric*], *blast*)

have $(\lambda x.\ 1) \in O[?F](\lambda x.\ ln\ (real\ (n\text{-}of\ x)))$
using *order-less-le-trans*[*OF exp-gt-zero*] *ln-ge-iff*
by (*intro landau-o.big-mono evt*[**where** *n=exp 1*])
(*simp add*: *abs-ge-iff*, *blast*)

hence *5*: $(\lambda x.\ 1) \in O[?F](\lambda x.\ ln\ (real\ (n\text{-}of\ x)) + ln\ (real\ (m\text{-}of\ x)))$
by (*intro landau-sum-1 evt*[**where** *n=1* **and** *m=1*], *auto*)

have $(\lambda x.\ -ln(of\text{-}rat\ (\varepsilon\text{-}of\ x))) \in O[?F](\lambda x.\ ln\ (1\ /\ real\text{-}of\text{-}rat\ (\varepsilon\text{-}of\ x)))$
by (*intro landau-o.big-mono evt*) (*auto simp add:ln-div*)
hence *6*: $(\lambda x.\ real\ (s2\text{-}of\ x)) \in O[?F](\lambda x.\ ln\ (1\ /\ real\text{-}of\text{-}rat\ (\varepsilon\text{-}of\ x)))$
unfolding *s2-of-def*
by (*intro landau-nat-ceil 3*, *simp*)

have *7*: $(\lambda\text{-}.\ 1) \in O[?F](\lambda x.\ real\ (n\text{-}of\ x)\ powr\ (1 - 1\ /\ real\ (k\text{-}of\ x)))$
by (*intro landau-o.big-mono evt*[**where** *n=1* **and** *k=1*])
(*auto simp add*: *ge-one-powr-ge-zero*)

have *8*: $(\lambda\text{-}.\ 1) \in O[?F](g1)$
unfolding *g1-def* **by** (*intro landau-o.big-mult-1 1 7 4*)

have $(\lambda x.\ 3 * (real\ (k\text{-}of\ x) * (n\text{-}of\ x)\ powr\ (1 - 1\ /\ real\ (k\text{-}of\ x))\ /\ (of\text{-}rat$

$(\delta$-*of* $x))^2))$
  $\in O[?F](g1)$
  by (*subst landau-o.big.cmult-in-iff*, *simp*, *simp add*:*g1-def*)
 hence *9*: (λx. *real* (*s1-of x*)) $\in O[?F](g1)$
  unfolding *s1-of-def* **by** (*intro landau-nat-ceil 8*, *auto simp*:*ac-simps*)

 have *10*: (λ-. *1*) $\in O[?F](g)$
  unfolding *g-def* **by** (*intro landau-o.big-mult-1 8 3 5*)

 have (λx. *real* (*s1-of x*)) $\in O[?F](g)$
  unfolding *g-def* **by** (*intro landau-o.big-mult-1 5 3 9*)
 hence (λx. *ln* (*real* (*s1-of x*) + *1*)) $\in O[?F](g)$
  using *10* **by** (*intro landau-ln-3 sum-in-bigo*, *auto*)
 hence *11*: (λx. *log 2* (*real* (*s1-of x*) + *1*)) $\in O[?F](g)$
  by (*simp add*:*log-def*)

 have *12*: (λx. *ln* (*real* (*s2-of x*) + *1*)) $\in O[?F]$(λx. *ln* (*1* / *real-of-rat* (ε-*of x*)))
  using *evt*[**where** ε=*2*] *6 3*
  by (*intro landau-ln-3 sum-in-bigo*, *auto*)

 have *13*: (λx. *log 2* (*real* (*s2-of x*) + *1*)) $\in O[?F](g)$
  unfolding *g-def*
  by (*rule landau-o.big-mult-1*, *rule landau-o.big-mult-1'*, *auto simp add*: *8 5 12*
log-def)

 have (λx. *real* (*k-of x*)) $\in O[?F](g1)$
  unfolding *g1-def* **using** *7 4*
  by (*intro landau-o.big-mult-1*, *simp-all*)
 hence (λx. *log 2* (*real* (*k-of x*) + *1*)) $\in O[?F](g1)$
  by (*simp add*:*log-def*) (*intro landau-ln-3 sum-in-bigo 8*, *auto*)
 hence *14*: (λx. *log 2* (*real* (*k-of x*) + *1*)) $\in O[?F](g)$
  unfolding *g-def* **by** (*intro landau-o.big-mult-1 3 5*)

 have *15*: (λx. *log 2* (*real* (*m-of x*) + *1*)) $\in O[?F](g)$
  unfolding *g-def* **using** *2 8 3*
  by (*intro landau-o.big-mult-1'*, *simp-all*)

 have (λx. *ln* (*real* (*n-of x*) + *1*)) $\in O[?F]$(λx. *ln* (*real* (*n-of x*)))
  by (*intro landau-ln-2*[**where** *a*=*2*] *eventually-mono*[*OF evt*[**where** *n*=*2*]] *sum-in-bigo*
1, *auto*)
 hence (λx. *log 2* (*real* (*n-of x*) + *1*)) $\in O[?F]$(λx. *ln* (*real* (*n-of x*)) + *ln* (*real*
(*m-of x*)))
  by (*intro landau-sum-1 evt*[**where** *n*=*1* **and** *m*=*1*])
   (*auto simp add*:*log-def*)
 hence *16*: (λx. *real* (*s1-of x*) $*$ *real* (*s2-of x*) $*$
  (*2* + *2* $*$ *log 2* (*real* (*n-of x*) + *1*) + *2* $*$ *log 2* (*real* (*m-of x*) + *1*))) $\in O[?F](g)$
  unfolding *g-def* **using** *9 6 5 2*
  by (*intro landau-o.mult sum-in-bigo*, *auto*)

have *fk-space-usage* = (λx. *fk-space-usage* (*k-of* x, *n-of* x, *m-of* x, ε*-of* x, δ*-of* x))
by (*simp add*:*case-prod-beta′ k-of-def n-of-def* ε*-of-def* δ*-of-def m-of-def*)
also have ... $\in$ *O*[*?F*](*g*)
using *10 11 13 14 15 16*
by (*simp add*:*fun-cong*[*OF s1-of-def*[*symmetric*]] *fun-cong*[*OF s2-of-def*[*symmetric*]] *Let-def*)
(*intro sum-in-bigo*, *auto*)
also have ... = *O*[*?F*](*?rhs*)
by (*simp add*:*case-prod-beta′ g1-def g-def n-of-def* ε*-of-def* δ*-of-def m-of-def k-of-def*)
finally show *?thesis* **by** *simp*
qed

end

9 Tutorial on the use of Pseudorandom-Objects

theory *Tutorial-Pseudorandom-Objects*
imports
Universal-Hash-Families.Pseudorandom-Objects-Hash-Families
Expander-Graphs.Pseudorandom-Objects-Expander-Walks
Equivalence-Relation-Enumeration.Equivalence-Relation-Enumeration
Median-Method.Median
Concentration-Inequalities.Bienaymes-Identity
Frequency-Moments.Frequency-Moments
begin

This section is a tutorial for the use of pseudorandom objects. Starting from the approximation algorithm for the second frequency moment by Alon et al. [1], we will improve the solution until we achieve a space complexity of $\mathcal{O}(\ln n+\varepsilon^{-2}\ln(\delta^{-1})\ln m)$, where n denotes the range of the stream elements, m denotes the length of the stream, ε denotes the desired accuracy and δ denotes the desired failure probability.

The construction relies on a combination of pseudorandom object, in particular an expander walk and two chained hash families.

hide-const (**open**) *topological-space-class.discrete*
hide-const (**open**) *Abstract-Rewriting.restrict*
hide-fact (**open**) *Abstract-Rewriting.restrict-def*
hide-fact (**open**) *Henstock-Kurzweil-Integration.integral-cong*
hide-fact (**open**) *Henstock-Kurzweil-Integration.integral-mult-right*
hide-fact (**open**) *Henstock-Kurzweil-Integration.integral-diff*

The following lemmas show a one-side and two-sided Chernoff-bound for $\{0,1\}$-valued independent identically distributed random variables. This to show the similarity with expander walks, for which similar bounds can be established: *expander-chernoff-bound-one-sided* and *expander-chernoff-bound*.

lemma *classic-chernoff-bound-one-sided*:

```
  fixes l :: nat
  assumes AE x in measure-pmf p. f x ∈ {0,1::real}
  assumes (∫ x. f x ∂p) ≤ μ l > 0 γ ≥ 0
  shows measure (prod-pmf {0..<l} (λ-. p)) {w. (∑ i<l. f (w i))/l−μ≥γ} ≤ exp
(− 2 ∗ real l ∗ γˆ2)
    (is ?L ≤ ?R)
proof −
  define ν where ν = real l∗(∫ x. f x ∂p)
  let ?p = prod-pmf {0..<l} (λ-. p)

  have 1: prob-space.indep-vars (measure-pmf ?p) (λ-. borel) (λi x. f (x i)) {0..<l}
   by (intro prob-space.indep-vars-compose2[OF - indep-vars-Pi-pmf] prob-space-measure-pmf)
auto

  have f (y i) ∈ {0..1} if y ∈ {0..<l} →E set-pmf p i ∈ {0..<l} for y i
  proof −
    have y i ∈ set-pmf p using that by auto
    thus ?thesis using assms(1) unfolding AE-measure-pmf-iff by auto
  qed
  hence 2: AE x in measure-pmf ?p. f (x i) ∈ {0..1}
    if i ∈ {0..<l} for i
    using that by (intro AE-pmfI) (auto simp: set-prod-pmf)

  have (∑ i=0..<l. (∫ x. f (x i) ∂?p)) = (∑ i<l. (∫ x. f x ∂map-pmf (λx. x i)
?p))
    by (auto simp:atLeast0LessThan)
  also have ... = (∑ i<l. (∫ x. f x ∂p)) by (subst Pi-pmf-component) auto
  also have ... = ν unfolding ν-def by simp
  finally have 3: (∑ i=0..<l. (∫ x. f (x i) ∂prod-pmf {0..<l} (λ-. p))) = ν by
simp

  have 4: ν ≤ real l ∗ μ unfolding ν-def using assms(2) by (simp add: mult-le-cancel-left)

  interpret Hoeffding-ineq measure-pmf ?p {0..<l} λi x. f (x i) (λ-. 0) (λ-. 1) ν
    using 1 2 unfolding 3 by unfold-locales auto

  have ?L ≤ measure ?p {x. (∑ i=0..<l. f (x i)) ≥ real l∗μ + real l∗γ}
   using assms(3) by (intro pmf-mono) (auto simp:field-simps atLeast0LessThan)
  also have ... ≤ measure ?p {x ∈ space ?p. (∑ i=0..<l. f (x i)) ≥ ν + real l∗γ}
    using 4 by (intro pmf-mono) auto
  also have ... ≤ exp (− 2 ∗ (real l ∗ γ)ˆ2 / (∑ i=0..<l. (1 − 0)²))
    using assms(3,4) by (intro Hoeffding-ineq-ge) auto
  also have ... = ?R using assms(3) by (simp add:power2-eq-square)
  finally show ?thesis by simp
qed

lemma classic-chernoff-bound:
  assumes AE x in measure-pmf p. f x ∈ {0,1::real} l > (0::nat) γ ≥ 0
  defines μ ≡ (∫ x. f x ∂p)
```

shows *measure* (*prod-pmf* {*0*..<*l*} (λ-. *p*)) {*w*. $|(\sum i<l.\ f\ (w\ i))/l-\mu|\geq\gamma$} ≤ *2*∗*exp* (−*2*∗*real l*∗γ^*2*)
(**is** *?L* ≤ *?R*)
proof −
have [*simp*]:*integrable p f* **using** *assms*(*1*) **unfolding** *AE-measure-pmf-iff*
by (*intro integrable-bounded-pmf boundedI*[**where** *B*=*1*]) *auto*
let *?w* = *prod-pmf* {*0*..<*l*} (λ-. *p*)
have *?L* ≤ *measure ?w* {*w*. $(\sum i<l.\ f\ (w\ i))/l-\mu\geq\gamma$} + *measure ?w* {*w*. $(\sum i<l.\ f\ (w\ i))/l-\mu\leq-(\gamma)$}
by (*intro pmf-add*) *auto*
also have ... ≤ *exp* (−*2*∗*real l*∗γ^*2*) + *measure ?w* {*w*. $-((\sum i<l.\ f\ (w\ i))/l-\mu)\geq\gamma$}
using *assms* **by** (*intro add-mono classic-chernoff-bound-one-sided*) (*auto simp*:*algebra-simps*)
also have ... ≤ *exp* (−*2*∗*real l*∗γ^*2*) + *measure ?w* {*w*. $((\sum i<l.\ 1-f\ (w\ i))/l-(1-\mu))\geq\gamma$}
using *assms*(*2*) **by** (*auto simp*: *sum-subtractf field-simps*)
also have ... ≤ *exp* (−*2*∗*real l*∗γ^*2*) + *exp* (−*2*∗*real l*∗γ^*2*)
using *assms* **by** (*intro add-mono classic-chernoff-bound-one-sided*) *auto*
also have ... = *?R* **by** *simp*
finally show *?thesis* **by** *simp*
qed

Definition of the second frequency moment of a stream.

definition *F2* :: *'a list* ⇒ *real* **where**
F2 xs = ($\sum$ *x* ∈ *set xs*. (*of-nat* (*count-list xs x*)^*2*))

lemma *prime-power-ls*: *is-prime-power* (*pro-size* (ℒ [− *1*, *1*]))
proof −
have *is-prime-power* ((*2*::*nat*)^*1*) **by** (*intro is-prime-powerI*) *auto*
thus *is-prime-power* (*pro-size* (ℒ [− *1*, *1*])) **by** (*auto simp*:*list-pro-size numeral-eq-Suc*)
qed

lemma *prime-power-h2*: *is-prime-power* (*pro-size* (ℋ *4 n* (ℒ [− *1*, *1*::*real*])))
by (*intro hash-pro-size-prime-power prime-power-ls*) *auto*

abbreviation Ψ **where** Ψ ≡ *pmf-of-set* {−*1*,*1*::*real*}

lemma *f2-exp*:
assumes *finite* (*set-pmf p*)
assumes ⋀*I*. *I* ⊆ {*0*..<*n*} ⟹ *card I* ≤ *4* ⟹ *map-pmf* (λ*x*. (λ*i*∈*I*. *x i*)) *p* = *prod-pmf I* (λ-. Ψ)
assumes *set xs* ⊆ {*0*..<*n*::*nat*}
shows (∫ *h*. ($\sum$ *x* ← *xs*. *h x*)^*2* ∂*p*) = *F2 xs* (**is** *?L* = *?R*)
proof −
let *?c* = (λ*x*. *real* (*count-list xs x*))

have [*simp*]: *integrable* (*measure-pmf p*) *f* **for** *f* :: - ⇒ *real*
by (*intro integrable-measure-pmf-finite assms*)

```
have 0:(∫ h. h x ∗ h y ∂p) = of-bool (x = y)
  (is ?L1 = ?R1) if x ∈ set xs y ∈ set xs for x y
proof −
  have xy-lt-n: x < n y < n  using assms that by auto
  have card-xy: card {x,y} ≤ 4 by (cases x = y) auto

  have ?L1 = (∫ h. (h x ∗ h y) ∂map-pmf (λf. restrict f {x,y}) p)
    by simp
  also have ... = (∫ h. (h x ∗ h y) ∂prod-pmf {x,y} (λ-. Ψ))
   using xy-lt-n card-xy by (intro integral-cong assms(2) arg-cong[where f=measure-pmf])
auto
  also have ... = of-bool (x = y) (is ?L2 = ?R2)
  proof (cases x = y)
    case True
    hence ?L2 = (∫ h. (h x ^2) ∂prod-pmf {x} (λ-. pmf-of-set {−1,1}))
      unfolding power2-eq-square by simp
    also have ... = (∫ x. x^2 ∂pmf-of-set {−1,1})
      unfolding Pi-pmf-singleton by simp
    also have ... = 1 by (subst integral-pmf-of-set) auto
    also have ... = ?R2 using True by simp
    finally show ?thesis by simp
  next
    case False
   hence ?L2 = (∫ h. (∏ i∈{x,y}. h i) ∂prod-pmf {x,y} (λ-. pmf-of-set {−1,1}))
by simp
    also have ... = (∏ i∈{x,y}. (∫ x. x ∂pmf-of-set {−1,1}))
      by (intro expectation-prod-Pi-pmf integrable-measure-pmf-finite) auto
    also have ... = 0 using False by (subst integral-pmf-of-set) auto
    also have ... = ?R2 using False by simp
    finally show ?thesis by simp
  qed
  finally show ?thesis by simp
qed

have ?L = (∫ h. (∑ x ∈ set xs. real (count-list xs x) ∗ h x)^2 ∂p)
  unfolding sum-list-eval by simp
also have ... = (∫ h. (∑ x ∈ set xs. (∑ y ∈ set xs. (?c x ∗ ?c y) ∗ h x ∗ h y))
∂p)
  unfolding power2-eq-square sum-distrib-left sum-distrib-right by (simp add:ac-simps)
also have ... = (∑ x ∈ set xs. (∑ y ∈ set xs. (∫ h. (?c x ∗ ?c y) ∗ h x ∗ h y
∂p))) by simp
also have ... = (∑ x ∈ set xs. (∑ y ∈ set xs. ?c x ∗ ?c y ∗ (∫ h. h x ∗ h y ∂p)))
  by (subst integral-mult-right[symmetric]) (simp-all add:ac-simps)
also have ... = (∑ x ∈ set xs. (∑ y ∈ set xs. ?c x ∗ ?c y ∗ of-bool (x = y)))
  by (intro sum.cong refl) (simp add: 0)
also have ... = (∑ x ∈ set xs. ?c x^2)
  unfolding of-bool-def by (simp add:if-distrib if-distribR sum.If-cases power2-eq-square)
also have ... = F2 xs unfolding F2-def by simp
finally show ?thesis by simp
```

qed

lemma *f2-exp-sq*:
  assumes *finite* (*set-pmf p*)
  assumes $\bigwedge I.\ I \subseteq \{0..<n\} \Longrightarrow card\ I \leq 4 \Longrightarrow map\text{-}pmf\ (\lambda x.\ (\lambda i{\in}I.\ x\ i))\ p =$ $prod\text{-}pmf\ I\ (\lambda\text{-}.\ \Psi)$
  assumes *set xs* $\subseteq \{0..<n::nat\}$
  shows $(\int h.\ ((\sum x{\leftarrow}xs.\ h\ x)\hat{\ }2)\hat{\ }2\ \partial p) \leq 3 * F2\ xs\hat{\ }2$ (**is** $?L \leq ?R$)
proof −
  let $?c = (\lambda x.\ real\ (count\text{-}list\ xs\ x))$

  have [*simp*]: *integrable* (*measure-pmf p*) *f* **for** *f* :: - ⇒ *real*
    by (*intro integrable-measure-pmf-finite assms*)

  define *S* **where** *S* = *set xs*

  have *a*: *finite S* **unfolding** *S-def* **by** *simp*

  define *Q* :: *nat* ⇒ *nat* ⇒ *nat* ⇒ *nat* ⇒ *real*
    where *Q a b c d* =
      $of\text{-}bool(a{=}b{\wedge}c{=}d{\wedge}a{\neq}c) + of\text{-}bool(a{=}c{\wedge}b{=}d{\wedge}a{\neq}b) +$
      $of\text{-}bool(a{=}d{\wedge}b{=}c{\wedge}a{\neq}b) + of\text{-}bool(a{=}b{\wedge}b{=}c{\wedge}c{=}d)$ **for** *a b c d*

  have *cases*: $(\int h.\ h\ a{*}h\ b{*}h\ c{*}h\ d\ \partial p) = Q\ a\ b\ c\ d$ (**is** $?L1 = ?R1$)
    if $a \in S\ b{\in}S\ c \in S\ d \in S$ **for** *a b c d*
  proof −
    have *card* {*a*,*b*,*c*,*d*} = *card* (*set* [*a*,*b*,*c*,*d*]) **by** (*intro arg-cong*[**where** *f*=*card*]) *auto*
    also have ... ≤ *length* [*a*,*b*,*c*,*d*] **by** (*intro card-length*)
    finally have *card*: *card* {*a*, *b*, *c*, *d*} ≤ *4* **by** *simp*

    have $?L1 = (\int h.\ h\ a{*}h\ b{*}h\ c{*}h\ d\ \partial map\text{-}pmf\ (\lambda f.\ restrict\ f\ \{a,b,c,d\})\ p)$ **by** *simp*
    also have ... $= (\int h.\ h\ a{*}h\ b{*}h\ c{*}h\ d\ \partial prod\text{-}pmf\ \{a,b,c,d\}\ (\lambda\text{-}.\ \Psi))$ **using** *that assms*(*3*)
      by (*intro integral-cong arg-cong*[**where** *f*=*measure-pmf*] *assms*(*2*) *card*) (*auto simp*:*S-def*)
    also have ... $= (\int h.\ (\prod i \leftarrow [a,b,c,d].\ h\ i)\ \partial prod\text{-}pmf\ \{a,b,c,d\}\ (\lambda\text{-}.\ \Psi))$ **by** (*simp add*:*ac-simps*)
    also have ... $= (\int h.\ (\prod i \in \{a,b,c,d\}.\ h\ i\hat{\ }count\text{-}list\ [a,b,c,d]\ i)\ \partial prod\text{-}pmf\ \{a,b,c,d\}\ (\lambda\text{-}.\ \Psi))$
      by (*subst prod-list-eval*) *auto*
    also have ... $= (\prod i \in \{a,b,c,d\}.\ (\int x.\ x\hat{\ }count\text{-}list\ [a,b,c,d]\ i\ \partial\Psi))$
      by (*intro expectation-prod-Pi-pmf integrable-measure-pmf-finite*) *auto*
    also have ... $= (\prod i \in \{a,b,c,d\}.\ of\text{-}bool\ (even\ (count\text{-}list\ [a,b,c,d]\ i)))$
      by (*intro prod.cong refl*) (*auto simp*:*integral-pmf-of-set*)
    also have ... $= (\prod i \in set\ (remdups\ [a,b,c,d]).\ of\text{-}bool\ (even\ (count\text{-}list\ [a,b,c,d]\ i)))$
      by (*intro prod.cong refl*) *auto*

also have ... = ($\prod$ *i* ← *remdups* [*a*,*b*,*c*,*d*]. *of-bool* (*even* (*count-list* [*a*,*b*,*c*,*d*] *i*)))
by (*intro prod.distinct-set-conv-list*) *auto*
also have ... = *Q a b c d* **unfolding** *Q-def* **by** *simp*
finally show *?thesis* **by** *simp*
qed

have *?L* = ($\int$ *h*. ($\sum$ *x* ∈ *S*. *real* (*count-list xs x*) ∗ *h x*)^*4* ∂*p*)
unfolding *S-def sum-list-eval* **by** *simp*
also have ... = ($\int$ *h*. ($\sum$ *a*∈*S*. ($\sum$ *b*∈*S*.($\sum$ *c*∈*S*.($\sum$ *d*∈*S*.(*?c a*∗*?c b*∗*?c c*∗*?c d*)∗*h a*∗*h b*∗*h c*∗*h d*)))) ∂*p*)
unfolding *power4-eq-xxxx sum-distrib-left sum-distrib-right* **by** (*simp add*:*ac-simps*)
also have ... = ($\sum$ *a*∈*S*.($\sum$ *b*∈*S*.($\sum$ *c*∈*S*.($\sum$ *d*∈*S*.($\int$ *h*. (*?c a*∗*?c b*∗*?c c*∗*?c d*)∗*h a*∗*h b*∗*h c*∗*h d* ∂*p*)))))
by *simp*
also have ... = ($\sum$ *a*∈*S*.($\sum$ *b*∈*S*.($\sum$ *c*∈*S*.($\sum$ *d*∈*S*. (*?c a*∗*?c b*∗*?c c*∗*?c d*)∗($\int$ *h*. *h a*∗*h b*∗*h c*∗*h d* ∂*p*)))))
by (*subst integral-mult-right*[*symmetric*]) (*simp-all add*:*ac-simps*)
also have ... = ($\sum$ *a*∈*S*.($\sum$ *b*∈*S*.($\sum$ *c*∈*S*.($\sum$ *d*∈*S*. (*?c a*∗*?c b*∗*?c c*∗*?c d*)∗(*Q a b c d*)))))
by (*intro sum.cong refl*) (*simp add*:*cases*)
also have ... = *1*∗($\sum$ *a*∈*S*. *?c a*^*4*) + *3*∗($\sum$ *a*∈*S*. ($\sum$ *b*∈*S*. *?c a*^*2* ∗ *?c b*^*2* ∗ *of-bool*(*a*≠*b*)))
unfolding *Q-def*
by (*simp add*: *sum.distrib distrib-left sum-collapse*[*OF a*] *ac-simps sum-distrib-left*[*symmetric*] *power2-eq-square power4-eq-xxxx*)
also have ... ≤ *3*∗($\sum$ *a*∈*S*. *?c a*^*4*) + *3*∗($\sum$ *a*∈*S*. ($\sum$ *b*∈*S*. *?c a*^*2* ∗ *?c b*^*2* ∗ *of-bool*(*a*≠*b*)))
by (*intro add-mono mult-right-mono sum-nonneg*) *auto*
also have ... = *3*∗($\sum$ *a*∈*S*. ($\sum$ *b*∈*S*. *?c a*^*2* ∗ *?c b*^*2* ∗ (*of-bool* (*a*=*b*) + *of-bool*(*a*≠*b*))))
using *a* **by** (*simp add*: *sum.distrib distrib-left*)
also have ... = *3*∗($\sum$ *a*∈*S*. ($\sum$ *b*∈*S*. *?c a*^*2* ∗ *?c b*^*2* ∗ *1*))
by (*intro sum.cong arg-cong2*[**where** *f*=(∗)] *refl*) *auto*
also have ... = *3* ∗ *F2 xs*^*2* **unfolding** *F2-def power2-eq-square*
by (*simp add*: *S-def sum-distrib-left sum-distrib-right ac-simps*)
finally show *?L* ≤ *3* ∗ *F2 xs*^*2* **by** *simp*
qed

lemma *f2-var*:
assumes *finite* (*set-pmf p*)
assumes $\bigwedge$*I*. *I* ⊆ {*0*..<*n*} ⟹ *card I* ≤ *4* ⟹ *map-pmf* (λ*x*. (λ*i*∈*I*. *x i*)) *p* = *prod-pmf I* (λ-. Ψ)
assumes *set xs* ⊆ {*0*..<*n*::*nat*}
shows *measure-pmf.variance p* (λ*h*. ($\sum$ *x*←*xs*. *h x*)^*2*) ≤ *2*∗ *F2 xs*^*2*
(**is** *?L* ≤ *?R*)
proof −
have [*simp*]: *integrable* (*measure-pmf p*) *f* **for** *f* :: - ⇒ *real*
by (*intro integrable-measure-pmf-finite assms*)

have *?L* = (∫ *h*. ((∑ *x*←*xs*. *h x*)^*2*)^*2* ∂*p*) − *F2 xs*^*2*
by (*subst measure-pmf.variance-eq*) (*simp-all add*:*f2-exp*[*OF assms*(*1*−*3*)])
also have ... ≤ *3* ∗ *F2 xs*^*2* − *F2 xs*^*2*
by (*intro diff-mono f2-exp-sq*[*OF assms*]) *auto*
finally show *?thesis* **by** *simp*
qed

lemma
assumes *s* ∈ *set-pmf* ($\mathcal{H}_P$ *4 n* ($\mathcal{L}$ [−*1*,*1*]))
assumes *set xs* ⊆ {*0*..<*n*}
shows *f2-exp-hp*: (∫ *h*. (∑ *x* ← *xs*. *h x*)^*2* ∂*sample-pro s*) = *F2 xs* (**is** *?T1*)
and *f2-exp-sq-hp*: (∫ *h*. ((∑ *x* ← *xs*. *h x*)^*2*)^*2* ∂*sample-pro s*) ≤ *3*∗ *F2 xs*^*2*
(**is** *?T2*)
and *f2-var-hp*: *measure-pmf.variance s* (λ*h*. (∑ *x* ← *xs*. *h x*)^*2*) ≤ *2*∗ *F2 xs*^*2*
(**is** *?T3*)
proof −
have *0*:*map-pmf* (λ*x*. *restrict x I*) (*sample-pro s*) = *prod-pmf I* (λ-. Ψ) (**is** *?L*
= -)
if *I* ⊆ {*0*..<*n*} *card I* ≤ *4* **for** *I*
proof −
have *?L* = *prod-pmf I* (λ-. *sample-pro* ($\mathcal{L}$ [− *1*, *1*]))
using *that* **by** (*intro hash-pro-pmf-distr*[*OF - assms*(*1*)] *prime-power-ls*) *auto*
also have ... = *prod-pmf I* (λ-. Ψ) **by** (*subst list-pro-2*) *auto*
finally show *?thesis* **by** *simp*
qed

show *?T1* **by** (*intro f2-exp*[*OF - - assms*(*2*)] *finite-pro-set 0*) *simp*
show *?T2* **by** (*intro f2-exp-sq*[*OF - - assms*(*2*)] *finite-pro-set 0*) *simp*
show *?T3* **by** (*intro f2-var*[*OF - - assms*(*2*)] *finite-pro-set 0*) *simp*
qed

lemmas *f2-exp-h* = *f2-exp-hp*[*OF hash-pro-in-hash-pro-pmf*[*OF prime-power-ls*]]
lemmas *f2-var-h* = *f2-var-hp*[*OF hash-pro-in-hash-pro-pmf*[*OF prime-power-ls*]]

lemma *F2-definite*:
assumes *xs* ≠ []
shows *F2 xs* > *0*
proof −
have *0* < *real* (*card* (*set xs*)) **using** *assms* **by** (*simp add*: *card-gt-0-iff*)
also have ... = (∑ *x* ∈ *set xs*. *1*) **by** *simp*
also have ... ≤ *F2 xs* **using** *count-list-gr-1* **unfolding** *F2-def* **by** (*intro sum-mono*)
force
finally show *?thesis* **by** *simp*
qed

The following algorithm uses a completely random function, accordingly it requires a lot of space: $\mathcal{O}(n + \ln m)$.

fun *example-1* :: *nat* ⇒ *nat list* ⇒ *real pmf*
where *example-1 n xs* =

do {
h ← *prod-pmf* {*0*..<*n*} (λ-. *pmf-of-set* {−*1*,*1*::*real*});
return-pmf ((∑*x* ← *xs*. *h x*)^*2*)
}

lemma *example-1-correct*:
assumes *set xs* ⊆ {*0*..<*n*}
shows
measure-pmf.expectation (*example-1 n xs*) *id* = *F2 xs* (**is** *?L1* = *?R1*)
measure-pmf.variance (*example-1 n xs*) *id* ≤ *2* ∗ *F2 xs*^*2* (**is** *?L2* ≤ *?R2*)
proof −
have *?L1* = (∫ *h*. (∑*x* ← *xs*. *h x*)^*2* ∂*prod-pmf* {*0*..<*n*} (λ-. Ψ))
by (*simp add*:*map-pmf-def*[*symmetric*])
also have ... = *?R1* **using** *assms* **by** (*intro f2-exp*)
(*auto intro*: *Pi-pmf-subset*[*symmetric*] *simp add*:*restrict-def set-Pi-pmf*)
finally show *?L1* = *?R1* **by** *simp*

have *?L2* = *measure-pmf.variance* (*prod-pmf* {*0*..<*n*} (λ-. Ψ)) (λ*h*. (∑*x* ← *xs*. *h x*)^*2*)
by (*simp add*:*map-pmf-def*[*symmetric*] *atLeast0LessThan*)
also have ... ≤ *?R2*
using *assms* **by** (*intro f2-var*)
(*auto intro*: *Pi-pmf-subset*[*symmetric*] *simp add*:*restrict-def set-Pi-pmf*)
finally show *?L2* ≤ *?R2* **by** *simp*
qed

This version replaces a the use of completely random function with a pseudorandom object, it requires a lot less space: $\mathcal{O}(\ln n + \ln m)$.

fun *example-2* :: *nat* ⇒ *nat list* ⇒ *real pmf*
where *example-2 n xs* =
do {
h ← *sample-pro* (ℋ *4 n* (ℒ [−*1*,*1*]));
return-pmf ((∑*x* ← *xs*. *h x*)^*2*)
}

lemma *example-2-correct*:
assumes *set xs* ⊆ {*0*..<*n*}
shows
measure-pmf.expectation (*example-2 n xs*) *id* = *F2 xs* (**is** *?L1* = *?R1*)
measure-pmf.variance (*example-2 n xs*) *id* ≤ *2* ∗ *F2 xs*^*2* (**is** *?L2* ≤ *?R2*)
proof −
have *?L1* = (∫ *h*. (∑*x* ← *xs*. *h x*)^*2* ∂*sample-pro* (ℋ *4 n* (ℒ [−*1*,*1*])))
by (*simp add*:*map-pmf-def*[*symmetric*])
also have ... = *?R1*
using *assms* **by** (*intro f2-exp-h*) *auto*
finally show *?L1* = *?R1* **by** *simp*

have *?L2* = *measure-pmf.variance* (*sample-pro* (ℋ *4 n* (ℒ [−*1*,*1*]))) (λ*h*. (∑*x* ← *xs*. *h x*)^*2*)

by (*simp add*:*map-pmf-def*[*symmetric*])
also have ... $\leq$ *?R2*
using *assms* **by** (*intro f2-var-h*) *auto*
finally show *?L2* $\leq$ *?R2* **by** *simp*
qed

The following version replaces the deterministic construction of the pseudo-random object with a randomized one. This algorithm is much faster, but the correctness proof is more difficult.

fun *example-3* :: *nat* $\Rightarrow$ *nat list* $\Rightarrow$ *real pmf*
where *example-3 n xs* =
do {
$h \leftarrow$ *sample-pro* =<< $\mathcal{H}_P$ *4 n* ($\mathcal{L}$ [-1,*1*]);
return-pmf (($\sum x \leftarrow$ *xs*. *h x*)^*2*)
}

lemma
assumes *set xs* $\subseteq$ {*0*..<*n*}
shows
measure-pmf.expectation (*example-3 n xs*) *id* = *F2 xs* (**is** *?L1* = *?R1*)
measure-pmf.variance (*example-3 n xs*) *id* $\leq$ *2* $*$ *F2 xs*^*2* (**is** *?L2* $\leq$ *?R2*)
proof –
let *?p* = $\mathcal{H}_P$ *4 n* ($\mathcal{L}$ [-1,*1*::*real*])
let *?q* = *bind-pmf ?p sample-pro*

have |*h x*| $\leq$ *1* **if** *that1*: *M* $\in$ *set-pmf ?p h* $\in$ *pro-set M x* $\in$ *set xs* **for** *h M x*
proof –
obtain *i* **where** *1*:*h* = *pro-select M i*
using *that1*(*2*) **unfolding** *set-sample-pro*[*of M*] **by** *auto*
have *h x* $\in$ *pro-set* ($\mathcal{L}$ [-1,*1*::*real*])
unfolding *1* **using** *that*(*1*) **by** (*intro hash-pro-pmf-range*[*OF prime-power-ls*])
auto
thus *?thesis* **by** (*auto simp*: *list-pro-set*)
qed

hence *0*: *bounded* ((λ*xa*. *xa x*) ‘ *set-pmf ?q*) **if** *x* $\in$ *set xs* **for** *x*
using *that* **by** (*intro boundedI*[**where** *B*=*1*]) *auto*

have ($\int$ *h*. ($\sum x \leftarrow$ *xs*. *h x*)^*2* ∂*?q*) = ($\int$ *s*. ($\int$ *h*. ($\sum x \leftarrow$ *xs*. *h x*)^*2* ∂*sample-pro*
s) ∂*?p*)
by (*intro integral-bind-pmf bounded-pow bounded-sum-list 0*)
also have ... = ($\int$ *s*. *F2 xs* ∂*?p*)
by (*intro integral-cong-AE AE-pmfI f2-exp-hp*[*OF - assms*]) *simp-all*
also have ... = *?R1* **by** *simp*
finally have *a*:($\int$ *h*. ($\sum x \leftarrow$ *xs*. *h x*)^*2* ∂*?q*) = *?R1* **by** *simp*
thus *?L1* = *?R1* **by** (*simp add*:*map-pmf-def*[*symmetric*])

have *?L2* = *measure-pmf.variance ?q* (λ*h*. ($\sum x \leftarrow$ *xs*. *h x*)^*2*)
by (*simp add*:*map-pmf-def*[*symmetric*])

also have ... = $(\int h.\ ((\sum x \leftarrow xs.\ h\ x)\hat{}2)\hat{}2\ \partial ?q) - (\int h.\ (\sum x \leftarrow xs.\ h\ x)\hat{}2\ \partial ?q)\hat{}2$
by (*intro measure-pmf.variance-eq integrable-bounded-pmf bounded-pow bounded-sum-list 0*)
also have ... = $(\int s.\ (\int h.\ ((\sum x \leftarrow xs.\ h\ x)\hat{}2)\hat{}2\ \partial sample\text{-}pro\ s)\ \partial ?p) - (F2\ xs)^2$
unfolding *a*
by (*intro arg-cong2*[**where** *f*=(−)] *integral-bind-pmf refl bounded-pow bounded-sum-list 0*)
also have ... $\leq (\int s.\ 3{*}F2\ xs\hat{}2\ \ \partial ?p) - (F2\ xs)\hat{}2$
by (*intro diff-mono integral-mono-AE′ AE-pmfI f2-exp-sq-hp*[*OF - assms*]) *simp-all*
also have ... = *?R2* **by** *simp*
finally show *?L2* ≤ *?R2* **by** *simp*
qed

context
fixes $\varepsilon\ \delta$ :: *real*
assumes *ε-gt-0*: $\varepsilon > 0$
assumes *δ-range*: $\delta \in \{0{<}..{<}1\}$
begin

By using the mean of many independent parallel estimates the following algorithm achieves a relative accuracy of ε, with probability $\frac{3}{4}$. It requires $\mathcal{O}(\varepsilon^{-2}(\ln n + \ln m))$ bits of space.

fun *example-4* :: *nat* ⇒ *nat list* ⇒ *real pmf*
where *example-4 n xs* =
do {
let s = *nat* ⌈*8* / *ε*^*2*⌉;
h ← *prod-pmf* {*0*..<*s*} (λ-. *sample-pro* ($\mathcal{H}$ *4 n* ($\mathcal{L}$ [−*1*,*1*])));
return-pmf $((\sum j < s.\ (\sum x \leftarrow xs.\ h\ j\ x)\hat{}2)/s)$
}

lemma *example-4-correct-aux*:
assumes *set xs* ⊆ {*0*..<*n*}
defines *s* ≡ *nat* ⌈*8* / *ε*^*2*⌉
defines $R \equiv (\lambda h :: nat \Rightarrow nat \Rightarrow real.\ (\sum j{<}s.\ (\sum x{\leftarrow}xs.\ h\ j\ x)\hat{}2)/real\ s)$
assumes *fin*: *finite* (*set-pmf p*)
assumes *indep*: *prob-space.k-wise-indep-vars* (*measure-pmf p*) *2* (λ-. *discrete*) (λ*i x*. *x i*) {..<*s*}
assumes *comp*: $\bigwedge i.\ i < s \Longrightarrow$ *map-pmf* (λ*x*. *x i*) *p* = *sample-pro* ($\mathcal{H}$ *4 n* ($\mathcal{L}$ [−*1*,*1*]))
shows *measure p* {*h*. |*R h* − *F2 xs*| > *ε* ∗ *F2 xs*} ≤ *1*/*4* (**is** *?L* ≤ *?R*)
proof (*cases xs* = [])
case *True* **thus** *?thesis* **by** (*simp add*:*R-def F2-def*)
next
case *False*
note *f2-gt-0* = *F2-definite*[*OF False*]
let *?p* = *sample-pro* ($\mathcal{H}$ *4 n* ($\mathcal{L}$ [−*1*,*1*::*real*]))

have [*simp*]: *integrable* (*measure-pmf p*) *f* **for** *f* :: - ⇒ *real*
  by (*intro integrable-measure-pmf-finite fin*)

have $8\ /\ \varepsilon^2 > 0$ **using** *ε-gt-0* **by** (*intro divide-pos-pos*) *auto*
hence *0*:$\lceil 8\ /\ \varepsilon^2 \rceil > 0$ **by** *simp*
hence *1*: $s > 0$ **unfolding** *s-def* **by** *simp*

have (∫ *h*. *R h* ∂*p*) = (∑ *j*<*s*. (∫ *h*. (∑ *x*←*xs*. *h j x*)^*2* ∂*p*))/*real s* **unfolding** *R-def* **by** *simp*
also have ... = (∑ *j*<*s*. (∫ *h*. (∑ *x*←*xs*. *h x*)^*2* ∂(*map-pmf*(λ*h*. *h j*)*p*)))/*real s* **by** *simp*
also have ... = (∑ *j*<*s*. (∫ *h*. (∑ *x*←*xs*. *h x*)^*2* ∂*?p*))/*real s*
  by (*intro sum.cong arg-cong2*[**where** *f*=(/)] *refl*) (*simp add*: *comp*)
also have ... = *F2 xs* **using** *1* **unfolding** *f2-exp-h*[*OF assms*(*1*)] **by** *simp*
finally have *exp-R*: (∫ *h*. *R h* ∂*p*) = *F2 xs* **by** *simp*

have *measure-pmf.variance p R* = *measure-pmf.variance p* (λ*h*. (∑ *j*<*s*. (∑ *x*←*xs*. *h j x*)^*2*))/*s*^*2*
  unfolding *R-def* **by** (*subst measure-pmf.variance-divide*) *simp-all*
also have ... = (∑ *j*<*s*. *measure-pmf.variance p* (λ*h*. (∑ *x*←*xs*. *h j x*)^*2*))/*real s*^*2*
  by (*intro arg-cong2*[**where** *f*=(/)] *refl measure-pmf.bienaymes-identity-pairwise-indep-2*
    prob-space.indep-vars-compose2[*OF* - *prob-space.k-wise-indep-vars-subset*[*OF* - *indep*]]
      prob-space-measure-pmf) (*auto intro*:*finite-subset*)
also have ... = (∑ *j*<*s*. *measure-pmf.variance*(*map-pmf*(λ*h*. *h j*)*p*)(λ*h*. (∑ *x*←*xs*. *h x*)^*2*))/*real s*^*2*
  by *simp*
also have ... = (∑ *j*<*s*. *measure-pmf.variance ?p* (λ*h*. (∑ *x*←*xs*. *h x*)^*2*))/ *real s*^*2*
  by (*intro sum.cong arg-cong2*[**where** *f*=(/)] *refl*) (*simp add*: *comp*)
also have ... ≤ (∑ *j*<*s*. *2* ∗ *F2 xs*^*2*)/*real s*^*2*
  by (*intro divide-right-mono sum-mono f2-var-h*[*OF assms*(*1*)]) *simp*
also have ... = *2* ∗ *F2 xs*^*2*/*real s* **by** (*simp add*:*power2-eq-square divide-simps*)
also have ... = *2* ∗ *F2 xs*^*2*/ ⌈*8*/*ε*^*2*⌉
  using *less-imp-le*[*OF 0*] **unfolding** *s-def* **by** (*subst of-nat-nat*) *auto*
also have ... ≤ *2* ∗ *F2 xs*^*2* / (*8*/ *ε*^*2*)
  using *ε-gt-0* **by** (*intro divide-left-mono mult-pos-pos*) *simp-all*
also have ... = *ε*^*2* ∗ *F2 xs*^*2*/*4* **by** *simp*
finally have *var-R*: *measure-pmf.variance p R* ≤ *ε*^*2* ∗ *F2 xs*^*2*/*4* **by** *simp*

have (∫ *h*. *R h* ∂*p*) = (∑ *j*<*s*. (∫ *h*. (∑ *x*←*xs*. *h j x*)^*2* ∂*p*))/*real s* **unfolding** *R-def* **by** *simp*
also have ... = (∑ *j*<*s*. (∫ *h*. (∑ *x*←*xs*. *h x*)^*2* ∂(*map-pmf*(λ*h*. *h j*)*p*)))/*real s* **by** *simp*
also have ... = (∑ *j*<*s*. (∫ *h*. (∑ *x*←*xs*. *h x*)^*2* ∂*?p*))/*real s*
  by (*intro sum.cong arg-cong2*[**where** *f*=(/)] *refl*) (*simp add*:*comp*)
also have ... = *F2 xs* **using** *1* **unfolding** *f2-exp-h*[*OF assms*(*1*)] **by** *simp*

finally have *exp-R*: $(\int h.\ R\ h\ \partial p) = F2\ xs$ **by** *simp*

have $?L \leq measure\ p\ \{h.\ |R\ h - F2\ xs| \geq \varepsilon * F2\ xs\}$ **by** (*intro pmf-mono*) *auto*
also have ... $\leq \mathcal{P}(h\ in\ p.\ |R\ h - (\int h.\ R\ h\ \partial p)| \geq \varepsilon * F2\ xs)$ **unfolding** *exp-R* **by** *simp*
also have ... $\leq measure\text{-}pmf.variance\ p\ R\ /\ (\varepsilon * F2\ xs)\hat{}2$
using *f2-gt-0 ε-gt-0* **by** (*intro measure-pmf.Chebyshev-inequality*) *simp-all*
also have ... $\leq (\varepsilon\hat{}2 * F2\ xs\hat{}2/4)\ /\ (\varepsilon * F2\ xs)\hat{}2$
by (*intro divide-right-mono var-R*) *simp*
also have ... $= 1/4$ **using** *ε-gt-0 f2-gt-0* **by** (*simp add:divide-simps*)
finally show *?thesis* **by** *simp*
qed

lemma *example-4-correct*:
assumes $set\ xs \subseteq \{0..<n\}$
shows $\mathcal{P}(\omega\ in\ example\text{-}4\ n\ xs.\ |\omega - F2\ xs| > \varepsilon * F2\ xs) \leq 1/4$ (**is** $?L \leq ?R$)
proof −
define $s :: nat$ **where** $s = nat\ \lceil 8\ /\ \varepsilon\hat{}2\rceil$
define R **where** $R\ h = (\sum j{<}s.\ (\sum x{\leftarrow}xs.\ h\ j\ x)\hat{}2)/s$ **for** $h :: nat \Rightarrow nat \Rightarrow real$

let $?p = sample\text{-}pro\ (\mathcal{H}\ 4\ n\ (\mathcal{L}\ [-1,1{::}real]))$
let $?q = prod\text{-}pmf\ \{..<s\}\ (\lambda\text{-}.\ ?p)$

have $?L = (\int h.\ indicator\ \{h.\ |R\ h - F2\ xs| > \varepsilon * F2\ xs\}\ h\ \partial\ ?q)$
by (*simp add:Let-def measure-bind-pmf R-def s-def indicator-def atLeast0LessThan*)
also have ... $= measure\ ?q\ \{h.\ |R\ h - F2\ xs| > \varepsilon * F2\ xs\}$ **by** *simp*
also have ... $\leq ?R$ **unfolding** *R-def s-def*
by (*intro example-4-correct-aux[OF assms] prob-space.k-wise-indep-vars-triv prob-space-measure-pmf indep-vars-Pi-pmf*)
(*auto intro*: *finite-pro-set simp add:Pi-pmf-component set-Pi-pmf*)
finally show *?thesis* **by** *simp*
qed

Instead of independent samples, we can choose the seeds using a second pair-wise independent pseudorandom object. This algorithm requires only $\mathcal{O}(\ln n + \varepsilon^{-2} \ln m)$ bits of space.

fun $example\text{-}5 :: nat \Rightarrow nat\ list \Rightarrow real\ pmf$
where $example\text{-}5\ n\ xs =$
do {
let $s = nat\ \lceil 8\ /\ \varepsilon\hat{}2\rceil$;
$h \leftarrow sample\text{-}pro\ (\mathcal{H}\ 2\ s\ (\mathcal{H}\ 4\ n\ (\mathcal{L}\ [-1,1])))$;
$return\text{-}pmf\ ((\sum j < s.\ (\sum x \leftarrow xs.\ h\ j\ x)\hat{}2)/s)$
}

lemma *example-5-correct-aux*:
assumes $set\ xs \subseteq \{0..<n\}$
defines $s \equiv nat\ \lceil 8\ /\ \varepsilon\hat{}2\rceil$
defines $R \equiv (\lambda h :: nat \Rightarrow nat \Rightarrow real.\ (\sum j{<}s.\ (\sum x{\leftarrow}xs.\ h\ j\ x)\hat{}2)/real\ s)$

shows *measure* (*sample-pro* ($\mathcal{H}$ *2 s* ($\mathcal{H}$ *4 n* ($\mathcal{L}$ [−*1*,*1*])))) {*h*. |*R h* − *F2 xs*| > ε ∗ *F2 xs*} ≤ *1*/*4*
proof −
let *?p* = *sample-pro* ($\mathcal{H}$ *2 s* ($\mathcal{H}$ *4 n* ($\mathcal{L}$ [−*1*,*1*::*real*])))

have *prob-space.k-wise-indep-vars ?p 2* (λ-. *discrete*) (λ*i x*. *x i*) {..<*s*}
using *hash-pro-indep*[*OF prime-power-h2*]
by (*simp add*: *prob-space.k-wise-indep-vars-def*[*OF prob-space-measure-pmf*])

thus *?thesis* **unfolding** *R-def s-def*
by (*intro example-4-correct-aux*[*OF assms*(*1*)] *finite-pro-set*)
(*simp-all add*:*hash-pro-component*[*OF prime-power-h2*])
qed

lemma *example-5-correct*:
assumes *set xs* ⊆ {*0*..<*n*}
shows $\mathcal{P}$(ω *in example-5 n xs*. |ω − *F2 xs*| > ε ∗ *F2 xs*) ≤ *1*/*4* (**is** *?L* ≤ *?R*)
proof −
define *s* :: *nat* **where** *s* = *nat* ⌈*8* / ε^*2*⌉
define *R* **where** *R h* = ($\sum$*j*<*s*. ($\sum$*x*←*xs*. *h j x*)^*2*)/*s* **for** *h* :: *nat* ⇒ *nat* ⇒ *real*

let *?p* = *sample-pro* ($\mathcal{H}$ *2 s* ($\mathcal{H}$ *4 n* ($\mathcal{L}$ [−*1*,*1*::*real*])))

have *?L* = ($\int$ *h*. *indicator* {*h*. |*R h* − *F2 xs*| > ε ∗ *F2 xs*} *h* ∂ *?p*)
by (*simp add*:*Let-def measure-bind-pmf R-def s-def indicator-def*)
also have ... = *measure ?p* {*h*. |*R h* − *F2 xs*| > ε ∗ *F2 xs*} **by** *simp*
also have ... ≤ *?R* **unfolding** *R-def s-def* **by** (*intro example-5-correct-aux*[*OF assms*])
finally show *?thesis* **by** *simp*
qed

The following algorithm improves on the previous one, by achieving a success probability of δ. This works by taking the median of $\mathcal{O}(\ln(\delta^{-1}))$ parallel independent samples. It requires $\mathcal{O}(\ln(\delta^{-1})(\ln n + \varepsilon^{-2}\ln m))$ bits of space.

fun *example-6* :: *nat* ⇒ *nat list* ⇒ *real pmf*
where *example-6 n xs* =
do {
let s = *nat* ⌈*8* / ε^*2*⌉; *let t* = *nat* ⌈*8* ∗ *ln* (*1*/δ)⌉;
h ← *prod-pmf* {*0*..<*t*} (λ-. *sample-pro* ($\mathcal{H}$ *2 s* ($\mathcal{H}$ *4 n* ($\mathcal{L}$ [−*1*,*1*]))));
return-pmf (*median t* (λ*i*. (($\sum$*j* < *s*. ($\sum$*x* ← *xs*. *h i j x*)^*2*)/ *s*)))
}

lemma *example-6-correct*:
assumes *set xs* ⊆ {*0*..<*n*}
shows $\mathcal{P}$(ω *in example-6 n xs*. |ω − *F2 xs*| > ε ∗ *F2 xs*) ≤ δ (**is** *?L* ≤ *?R*)
proof −
define *s* **where** *s* = *nat* ⌈*8* / ε^*2*⌉
define *t* **where** *t* = *nat* ⌈*8* ∗ *ln*(*1*/δ)⌉

define R **where** $R\ h = (\sum j{<}s.\ (\sum x{\leftarrow}xs.\ h\ j\ x)\hat{\ }2)/s$ **for** $h :: nat \Rightarrow nat \Rightarrow real$
define I **where** $I = \{w.\ |w - F2\ xs| \le \varepsilon * F2\ xs\}$

have $8 * ln\ (1\ /\ \delta) > 0$ **using** *δ-range* **by** (*intro mult-pos-pos ln-gt-zero*) *auto*
hence *t-gt-0*: $t > 0$ **unfolding** *t-def* **by** *simp*
have *int-I*: *interval I* **unfolding** *interval-def I-def* **by** *auto*

let $?p = sample\text{-}pro\ (\mathcal{H}\ 2\ s\ (\mathcal{H}\ 4\ n\ (\mathcal{L}\ [-1,1::real])))$
let $?q = prod\text{-}pmf\ \{0..{<}t\}\ (\lambda\text{-}.\ ?p)$

have $(\int h.\ (of\text{-}bool\ (R\ h \notin I)::real)\ \partial ?p) = (\int h.\ indicator\ \{h.\ R\ h \notin I\}\ h\ \partial ?p)$
unfolding *of-bool-def indicator-def* **by** *simp*
also have ... $= measure\ ?p\ \{h.\ R\ h \notin I\}$ **by** *simp*
also have ... $\le 1/4$
using *example-5-correct-aux*[*OF assms*] **unfolding** *R-def s-def I-def* **by** (*simp add:not-le*)
finally have *0*: $(\int h.\ (of\text{-}bool\ (R\ h \notin I)::real)\ \partial ?p) \le 1/4$ **by** *simp*

have $?L = (\int h.\ indicator\ \{h.\ |median\ t\ (\lambda i.\ R\ (h\ i)) - F2\ xs| > \varepsilon * F2\ xs\}\ h\ \partial\ ?q)$
by (*simp add:Let-def measure-bind-pmf R-def s-def indicator-def t-def*)
also have ... $= measure\ ?q\ \{h.\ median\ t\ (\lambda i.\ R\ (h\ i)) \notin I\}$
unfolding *I-def* **by** (*simp add:not-le*)
also have ... $\le measure\ ?q\ \{h.\ t \le 2 * card\ \{k.\ k < t \wedge R\ (h\ k) \notin I\}\}$
using *median-est-rev*[*OF int-I*] **by** (*intro pmf-mono*) *auto*
also have ... $= measure\ ?q\ \{h.\ (\sum k{<}t.\ of\text{-}bool(R\ (h\ k) \notin I))/real\ t - 1/4 \ge (1/4)\}$
using *t-gt-0* **by** (*intro arg-cong2*[**where** *f=measure*]) (*auto simp:Int-def divide-simps*)
also have ... $\le exp\ (-\ 2 * real\ t * (1/4)\hat{\ }2)$
by (*intro classic-chernoff-bound-one-sided t-gt-0 AE-pmfI 0*) *auto*
also have ... $= exp\ (-\ (real\ t\ /\ 8))$ **using** *t-gt-0* **by** (*simp add:power2-eq-square*)
also have ... $\le exp\ (-\ of\text{-}int\ \lceil 8 * ln\ (1\ /\ \delta)\rceil\ /\ 8)$ **unfolding** *t-def*
by (*intro iffD2*[*OF exp-le-cancel-iff*] *divide-right-mono iffD2*[*OF neg-le-iff-le*]) *auto*
also have ... $\le exp\ (-\ (8 * ln\ (1\ /\ \delta))\ /\ 8)$
by (*intro iffD2*[*OF exp-le-cancel-iff*] *divide-right-mono iffD2*[*OF neg-le-iff-le*]) *auto*
also have ... $= exp\ (-\ ln\ (1\ /\ \delta))$ **by** *simp*
also have ... $= \delta$ **using** *δ-range* **by** (*subst ln-div*) *auto*
finally show *?thesis* **by** *simp*
qed

The following algorithm uses an expander random walk, instead of independent samples. It requires only $\mathcal{O}(\ln n + \ln(\delta^{-1})\varepsilon^{-2}\ln m)$ bits of space.

fun *example-7* :: *nat* ⇒ *nat list* ⇒ *real pmf*
where *example-7 n xs* =
do {

let s = nat $\lceil$*8 / ε^2*$\rceil$; *let t = nat* $\lceil$*32 ∗ ln (1/δ)*$\rceil$;
h ← sample-pro (ℰ t (1/8) (ℋ 2 s (ℋ 4 n (ℒ [−1,1]))));
return-pmf (median t (λi. ((∑j < s. (∑x ← xs. h i j x)^2)/ s)))
}

lemma *example-7-correct*:
assumes *set xs ⊆ {0..<n}*
shows *𝒫(ω in example-7 n xs. |ω − F2 xs| > ε ∗ F2 xs) ≤ δ* (**is** *?L ≤ ?R*)
proof −
define *s t* **where** *s-def*: *s = nat* $\lceil$*8 / ε^2*$\rceil$ **and** *t-def*: *t = nat* $\lceil$*32 ∗ ln(1/δ)*$\rceil$
define *R* **where** *R h = (∑j<s. (∑x←xs. h j x)^2)/s* **for** *h :: nat ⇒ nat ⇒ real*
define *I* **where** *I = {w. |w − F2 xs| ≤ ε∗F2 xs}*

have *8 ∗ ln (1 / δ) > 0* **using** *δ-range* **by** (*intro mult-pos-pos ln-gt-zero*) *auto*
hence *t-gt-0*: *t > 0* **unfolding** *t-def* **by** *simp*
have *int-I*: *interval I* **unfolding** *interval-def I-def* **by** *auto*

let *?p = sample-pro (ℋ 2 s (ℋ 4 n (ℒ [−1,1::real])))*
let *?q = sample-pro (ℰ t (1/8) (ℋ 2 s (ℋ 4 n (ℒ [−1,1]))))*

have *(∫h. (of-bool (R h ∉ I)::real) ∂?p) = (∫h. indicator {h. R h ∉ I} h ∂?p)*
by (*simp add:of-bool-def indicator-def*)
also have ... = *measure ?p {h. R h ∉ I}* **by** *simp*
also have ... ≤ *1/4*
using *example-5-correct-aux[OF assms]* **unfolding** *R-def s-def I-def* **by** (*simp add:not-le*)
finally have ∗: *(∫h. (of-bool (R h ∉ I)::real) ∂?p) ≤ 1/4* **by** *simp*

have *?L = (∫h. indicator {h. |median t (λi. R (h i)) − F2 xs| > ε ∗ F2 xs} h ∂ ?q)*
by (*simp add:Let-def measure-bind-pmf R-def s-def indicator-def t-def*)
also have ... = *measure ?q {h. median t (λi. R (h i)) ∉ I}*
unfolding *I-def* **by** (*simp add:not-le*)
also have ... ≤ *measure ?q {h. t ≤ 2 ∗ card {k. k < t ∧ R (h k) ∉ I}}*
using *median-est-rev[OF int-I]* **by** (*intro pmf-mono*) *auto*
also have ... = *measure ?q {h. 1/8 + 1/8 ≤ (∑k<t. of-bool(R (h k) ∉ I))/real t − 1/4}*
using *t-gt-0* **by** (*intro arg-cong2[***where** *f=measure] Collect-cong refl*)
(*auto simp add:of-bool-def sum.If-cases Int-def field-simps*)
also have ... ≤ *exp (− 2 ∗ real t ∗ (1/8)*^{2}*)*
by (*intro expander-chernoff-bound-one-sided t-gt-0* ∗) *auto*
also have ... = *exp (− (real t / 32))* **using** *t-gt-0* **by** (*simp add:power2-eq-square*)
also have ... ≤ *exp (− of-int* $\lceil$*32 ∗ ln (1 / δ)*$\rceil$ */ 32)* **unfolding** *t-def*
by (*intro iffD2[OF exp-le-cancel-iff] divide-right-mono iffD2[OF neg-le-iff-le]*) *auto*
also have ... ≤ *exp (− (32 ∗ ln (1 / δ)) / 32)*
by (*intro iffD2[OF exp-le-cancel-iff] divide-right-mono iffD2[OF neg-le-iff-le]*) *auto*

also have ... = *exp* (− *ln* (*1* / δ)) **by** *simp*
also have ... = δ **using** δ*-range* **by** (*subst ln-div*) *auto*
finally show *?thesis* **by** *simp*
qed

end

end

A Informal proof of correctness for the F_0 algorithm

This appendix contains a detailed informal proof for the new Rounding-KMV algorithm that approximates F_0 introduced in Section 6 for reference. It follows the same reasoning as the formalized proof.

Because of the amplification result about medians (see for example [1, §2.1]) it is enough to show that each of the estimates the median is taken from is within the desired interval with success probability $\frac{2}{3}$. To verify the latter, let $a_1, \ldots, a_m$ be the stream elements, where we assume that the elements are a subset of $\{0, \ldots, n-1\}$ and $0 < \delta < 1$ be the desired relative accuracy. Let p be the smallest prime such that $p \geq \max(n, 19)$ and let h be a random polynomial over $GF(p)$ with degree strictly less than 2. The algoritm also introduces the internal parameters t, r defined by:

$$t := \lceil 80\delta^{-2} \rceil \qquad\qquad r := 4 \log_2 \lceil \delta^{-1} \rceil + 23$$

The estimate the algorithm obtains is R, defined using:

$$H := \{\lfloor h(a) \rfloor_r | a \in A\} \qquad R := \begin{cases} tp\,(\min_t(H))^{-1} & \text{if } |H| \geq t \\ |H| & \text{othewise,} \end{cases}$$

where $A := \{a_1, \ldots, a_m\}$, $\min_t(H)$ denotes the t-th smallest element of H and $\lfloor x \rfloor_r$ denotes the largest binary floating point number smaller or equal to x with a mantissa that requires at most r bits to represent.[1] With these definitions, it is possible to state the main theorem as:

$$P(|R - F_0| \leq \delta |F_0|) \geq \frac{2}{3}.$$

which is shown separately in the following two subsections for the cases $F_0 \geq t$ and $F_0 < t$.

[1] This rounding operation is called *truncate-down* in Isabelle, it is defined in HOL-Library.Float.

A.1 Case $F_0 \geq t$

Let us introduce:

$$H^* := \{h(a) | a \in A\}^{\#} \qquad\qquad R^* := tp\left(\min_t^{\#}(H^*)\right)^{-1}$$

These definitions are modified versions of the definitions for H and R: The set H^* is a multiset, this means that each element also has a multiplicity, counting the number of *distinct* elements of A being mapped by h to the same value. Note that by definition: $|H^*| = |A|$. Similarly the operation $\min_t^{\#}$ obtains the t-th element of the multiset H (taking multiplicities into account). Note also that there is no rounding operation $\lfloor \cdot \rfloor_r$ in the definition of H^*. The key reason for the introduction of these alternative versions of H, R is that it is easier to show probabilistic bounds on the distances $|R^* - F_0|$ and $|R^* - R|$ as opposed to $|R - F_0|$ directly. In particular the plan is to show:

$$P\left(|R^* - F_0| > \delta' F_0\right) \leq \frac{2}{9}, \text{ and} \tag{1}$$

$$P\left(|R^* - F_0| \leq \delta' F_0 \wedge |R - R^*| > \frac{\delta}{4} F_0\right) \leq \frac{1}{9} \tag{2}$$

where $\delta' := \frac{3}{4}\delta$. I.e. the probability that R^* has not the relative accuracy of $\frac{3}{4}\delta$ is less that $\frac{2}{9}$ and the probability that assuming R^* has the relative accuracy of $\frac{3}{4}\delta$ but that R deviates by more that $\frac{1}{4}\delta F_0$ is at most $\frac{1}{9}$. Hence, the probability that neither of these events happen is at least $\frac{2}{3}$ but in that case:

$$|R - F_0| \leq |R - R^*| + |R^* - F_0| \leq \frac{\delta}{4} F_0 + \frac{3\delta}{4} F_0 = \delta F_0. \tag{3}$$

Thus we only need to show Equation 1 and 2. For the verification of Equation 1 let

$$Q(u) = |\{h(a) < u \mid a \in A\}|$$

and observe that $\min_t^{\#}(H^*) < u$ if $Q(u) \geq t$ and $\min_t^{\#}(H^*) \geq v$ if $Q(v) \leq t - 1$. To see why this is true note that, if at least t elements of A are mapped by h below a certain value, then the t-smallest element must also be within them, and thus also be below that value. And that the opposite direction of this conclusion is also true. Note that this relies on the fact that H^* is a multiset and that multiplicities are being taken into account, when computing the t-th smallest element. Alternatively, it is also possible to write $Q(u) = \sum_{a \in A} 1_{\{h(a) < u\}}$[2], i.e., Q is a sum of pairwise independent $\{0, 1\}$-valued random variables, with expectation $\frac{u}{p}$ and variance $\frac{u}{p} - \frac{u^2}{p^2}$.

[2] The notation 1_A is shorthand for the indicator function of A, i.e., $1_A(x) = 1$ if $x \in A$ and 0 otherwise.

[3] Using lineariy of expectation and Bienaymé's identity, it follows that $\mathrm{Var}\, Q(u) \leq \mathrm{E}\, Q(u) = |A|up^{-1} = F_0 up^{-1}$ for $u \in \{0, \ldots, p\}$.
For $v = \left\lfloor \frac{tp}{(1-\delta')F_0} \right\rfloor$ it is possible to conclude:

$$t - 1 \leq^{4} \frac{t}{(1-\delta')} - 3\sqrt{\frac{t}{(1-\delta')}} - 1 \leq \frac{F_0 v}{p} - 3\sqrt{\frac{F_0 v}{p}} \leq \mathrm{E}Q(v) - 3\sqrt{\mathrm{Var}Q(v)}$$

and thus using Tchebyshev's inequality:

$$\begin{aligned} P\left(R^* < \left(1-\delta'\right) F_0\right) &= P\left(\mathrm{rank}_t^{\#}(H^*) > \frac{tp}{(1-\delta')F_0}\right) \\ &\leq P(\mathrm{rank}_t^{\#}(H^*) \geq v) = P(Q(v) \leq t-1) \\ &\leq P\left(Q(v) \leq \mathrm{E}Q(v) - 3\sqrt{\mathrm{Var}Q(v)}\right) \leq \frac{1}{9}. \end{aligned} \tag{4}$$

Similarly for $u = \left\lceil \frac{tp}{(1+\delta')F_0} \right\rceil$ it is possible to conclude:

$$t \geq \frac{t}{(1+\delta')} + 3\sqrt{\frac{t}{(1+\delta')}} + 1 + 1 \geq \frac{F_0 u}{p} + 3\sqrt{\frac{F_0 u}{p}} \geq \mathrm{E}Q(u) + 3\sqrt{\mathrm{Var}Q(v)}$$

and thus using Tchebyshev's inequality:

$$\begin{aligned} P\left(R^* > \left(1+\delta'\right) F_0\right) &= P\left(\mathrm{rank}_t^{\#}(H^*) < \frac{tp}{(1+\delta')F_0}\right) \\ &\leq P(\mathrm{rank}_t^{\#}(H^*) < u) = P(Q(u) \geq t) \\ &\leq P\left(Q(u) \geq \mathrm{E}Q(u) + 3\sqrt{\mathrm{Var}Q(u)}\right) \leq \frac{1}{9}. \end{aligned} \tag{5}$$

Note that Equation 4 and 5 confirm Equation 1. To verfiy Equation 2, note that

$$\min_t(H) = \lfloor \min_t^{\#}(H^*) \rfloor_r \tag{6}$$

if there are no collisions, induced by the application of $\lfloor h(\cdot) \rfloor_r$ on the elements of A. Even more carefully, note that the equation would remain true, as long as there are no collision within the smallest t elements of H^*. Because Equation 2 needs to be shown only in the case where $R^* \geq (1-\delta')F_0$, i.e., when $\min_t^{\#}(H^*) \leq v$, it is enough to bound the probability of a collision in the range $[0; v]$. Moreover Equation 6 implies $|\min_t(H) - \min_t^{\#}(H^*)| \leq \max(\min_t^{\#}(H^*), \min_t(H))2^{-r}$ from which it is possible to derive $|R^* - R| \leq \frac{\delta}{4} F_0$. Another important fact is that h is injective with probability $1 - \frac{1}{p}$,

[3] A consequence of h being chosen uniformly from a 2-independent hash family.
[4] The verification of this inequality is a lengthy but straightforward calculcation using the definition of δ' and t.

this is because h is chosen uniformly from the polynomials of degree less than 2. If it is a degree 1 polynomial it is a linear function on $GF(p)$ and thus injective. Because $p \geq 18$ the probability that h is not injective can be bounded by $1/18$. With these in mind, we can conclude:

$$\begin{aligned}
& P\left(|R^* - F_0| \leq \delta' F_0 \wedge |R - R^*| > \frac{\delta}{4} F_0\right) \\
\leq\ & P\left(R^* \geq (1-\delta')F_0 \wedge \min_t^{\#}(H^*) \neq \min_t(H) \wedge h \text{ inj.}\right) + P(\neg h \text{ inj.}) \\
\leq\ & P\left(\exists a \neq b \in A. \lfloor h(a) \rfloor_r = \lfloor h(b) \rfloor_r \leq v \wedge h(a) \neq h(b)\right) + \frac{1}{18} \\
\leq\ & \frac{1}{18} + \sum_{a \neq b \in A} P\left(\lfloor h(a) \rfloor_r = \lfloor h(b) \rfloor_r \leq v \wedge h(a) \neq h(b)\right) \\
\leq\ & \frac{1}{18} + \sum_{a \neq b \in A} P\left(|h(a) - h(b)| \leq v 2^{-r} \wedge h(a) \leq v(1 + 2^{-r}) \wedge h(a) \neq h(b)\right) \\
\leq\ & \frac{1}{18} + \sum_{a \neq b \in A} \sum_{\substack{a', b' \in \{0, \ldots, p-1\} \wedge a' \neq b' \\ |a' - b'| \leq v 2^{-r} \wedge a' \leq v(1+2^{-r})}} P(h(a) = a') P(h(b) = b') \\
\leq\ & \frac{1}{18} + \frac{5 F_0^2 v^2}{2p^2} 2^{-r} \leq \frac{1}{9}.
\end{aligned}$$

which shows that Equation 2 is true.

A.2 Case $F_0 < t$

Note that in this case $|H| \leq F_0 < t$ and thus $R = |H|$, hence the goal is to show that: $P(|H| \neq F_0) \leq \frac{1}{3}$. The latter can only happen, if there is a collision induced by the application of $\lfloor h(\cdot) \rfloor_r$. As before h is not injective with probability at most $\frac{1}{18}$, hence:

$$
\begin{aligned}
& P\left(|R - F_0| > \delta F_0\right) \leq P\left(R \neq F_0\right) \\
\leq\ & \frac{1}{18} + P\left(R \neq F_0 \wedge h \text{ inj.}\right) \\
\leq\ & \frac{1}{18} + P\left(\exists a \neq b \in A. \lfloor h(a) \rfloor_r = \lfloor h(b) \rfloor_r \wedge h \text{ inj.}\right) \\
\leq\ & \frac{1}{18} + \sum_{a \neq b \in A} P\left(\lfloor h(a) \rfloor_r = \lfloor h(b) \rfloor_r \wedge h(a) \neq h(b)\right) \\
\leq\ & \frac{1}{18} + \sum_{a \neq b \in A} P\left(|h(a) - h(b)| \leq p 2^{-r} \wedge h(a) \neq h(b)\right) \\
\leq\ & \frac{1}{18} + \sum_{a \neq b \in A} \sum_{\substack{a', b' \in \{0, \ldots, p-1\} \\ a' \neq b' \wedge |a' - b'| \leq p 2^{-r}}} P(h(a) = a') P(h(b) = b') \\
\leq\ & \frac{1}{18} + F_0^2 2^{-r+1} \leq \frac{1}{18} + t^2 2^{-r+1} \leq \frac{1}{9}.
\end{aligned}
$$

Which concludes the proof. □

References

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