

Babai's Nearest Plane Algorithm

Eric Ren, Sage Binder, and Katherine Kosaian

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Abstract

γ -CVP is the problem of finding a vector in L that is within γ times the closest possible to t , where L is a lattice and t is a target vector. If the basis for L is LLL-reduced, Babai's Closest Hyperplane algorithm solves γ -CVP for $\gamma = 2^{n/2}$, where n is the dimension of the lattice L , in time polynomial in n . This session formalizes said algorithm, using the AFP formalization of LLL [2, 1] and adapting a proof of correctness from the lecture notes of Stephens-Davidowitz [4].

Contents

1	Introduction	1
2	Locale setup for Babai	4
3	Coordinates	5
4	Lattice Lemmas	8
5	Lemmas on closest distance	8
6	More linear algebra lemmas	9
7	Coord-Invariance	11
8	Main Theorem	12

1 Introduction

The (exact) *closest vector problem* (CVP) is the problem of finding the closest vector within a lattice L to a target vector t . This is equivalent to finding the shortest vector in the *lattice coset* $L - t := \{l - t : l \in L\}$. There is a corresponding family of weaker problems, γ -CVP (where γ is some real parameter), where one needs only find a vector in $L - t$ whose length is at most γ times the shortest possible. Through a reduction to the *shortest*

vector problem [4], solutions to these problems may be used to factor rational polynomials. This problem is therefore of cryptographic interest.

Although exact CVP (or 1-CVP) is NP-Complete [3], Babai's Nearest Plane Algorithm solves $2^{n/2}$ -CVP, where n is the dimension of L , in polynomial time, provided that L is presented using an LLL-reduced basis with parameter $\alpha = 4/3$. The proof in this document is mostly a straightforward algebraicization of the proof in Stephens-Davidowitz' lecture notes. It makes use of the coordinate systems defined by the original basis (denoted β) and the Gram-Schmidt orthogonalization of that basis (denoted $\tilde{\beta}$). Let $[u]_\beta$ denote the representation of a vector u under β , with coordinates $[u]_\beta^j$; $j = 1, \dots, n$ (likewise for $\tilde{\beta}$). Also, let s_i denote the output of the algorithm after step i and let d be the shortest lattice coset vector, as witnessed by the vector v . The proof works by analysing the coordinates of $[s_n]_{\tilde{\beta}}$, showing that all are at most $1/2$ and that some later coordinates are exactly those of $[v]_{\tilde{\beta}}$.

The algorithm modifies coordinate $n - i$ in both bases for the last time in step i (formalized in lemma `coord_invariance`), during which both coordinates are decreased below $1/2$ (formalized in lemma `small_coord`). Combined, these facts imply that the output s_n has $|[s_n]_{\tilde{\beta}}^j| \leq 1/2$ for all indices j .

Since $\tilde{\beta}$ is orthogonal, we have

$$\|s_n\|^2 = \sum_{i=1}^n \left([s_n]_{\tilde{\beta}}^i \|\tilde{\beta}_i\| \right)^2, \quad (1)$$

so the preceding coordinate bounds $\|s_n\|^2$ by $\frac{1}{4} \sum_{i=1}^n \|\tilde{\beta}_i\|^2$. If the $\tilde{\beta}_i$ are all short compared to d , this bound suffices. In fact, if there is any short vector $\tilde{\beta}_I$ in $\tilde{\beta}$ then because β is LLL-reduced, any vector preceding $\tilde{\beta}_I$ in $\tilde{\beta}$ will not be much longer. This bounds the first I terms in Equation 1. By selecting I maximal, we may assume that $\tilde{\beta}$ ends in a series of $n - I$ long vectors. In this case it can be shown $[v]_{\tilde{\beta}}^j$ and $[s_n]_{\tilde{\beta}}^j$ differ by an integral amount for $j = I + 1, \dots, n$. Therefore, if $[v]_{\tilde{\beta}}^j$ and $[s_n]_{\tilde{\beta}}^j$ differ at all, they differ by at least 1, which would mean $|[v]_{\tilde{\beta}}^j| \geq 1/2$, since $|[s_n]_{\tilde{\beta}}^j| \leq 1/2$. This would force v to be longer than d , a contradiction. So $[v]_{\tilde{\beta}}^j = [s_n]_{\tilde{\beta}}^j$ for $j = I + 1, \dots, n$, which gives a tighter bound on the last $n - I$ terms in equation 1.

Precisely, let I denote $\max\{i : \|\tilde{\beta}_i\| \leq 2d\}$, meaning for all indices $j > I$, $\|\tilde{\beta}_j\| > 2d$. Now, for all $j > I$, $d^2 = \|v\|^2 \geq ([v]_{\tilde{\beta}}^j)^2 \|\tilde{\beta}_j\|^2 > ([v]_{\tilde{\beta}}^j)^2 \cdot 4d^2$, meaning $1/4 > (\tilde{\beta}^j)^2$, or $1/2 > |[v]_{\tilde{\beta}}^j|$. Since $|[s_j]_{\tilde{\beta}}^j| \leq 1/2$ from the previous section, $|[v]_{\tilde{\beta}}^j - [s_j]_{\tilde{\beta}}^j| < 1$. Using properties of the change-of-basis between $\beta, \tilde{\beta}$ formalized in the LLL AFP session, we show that $[v]_{\tilde{\beta}}^j - [s_j]_{\tilde{\beta}}^j =$

$[v]_\beta^j - [s_j]_\beta^j = [v - s_j]_\beta^j$, so that $\left| [v - s_j]_\beta^j \right| < 1$. But since $v - s_j$ lies in the lattice, $[v - s_j]_\beta^j$ is integral, so $\left| [v - s_j]_\beta^j \right| = 0$, meaning $[v]_\beta^j = [s_j]_\beta^j$. Lemma **coord_invariance** gives that $[v]_\beta^j = [s_j]_\beta^j = [s_n]_\beta^j$. This is formalized by lemma **correct_coord**.

Now $\|s_n\|^2 = \sum_{i=1}^n ([s_n]_{\tilde{\beta}}^i \|\tilde{\beta}_i\|)^2$, since $\tilde{\beta}$ is orthogonal. Splitting the sum around I equates this to $\sum_{i=1}^I ([s_n]_{\tilde{\beta}}^i)^2 + \sum_{i=I+1}^n ([s_n]_{\tilde{\beta}}^i)^2$. Lemma **small_coord** bounds the terms in the first sum by $\|\tilde{\beta}_i\|^2/4$, while lemma **correct_coord** bounds the terms in the second sum by d^2 , giving $\|s_n\|^2 \leq (n - I)d^2 + \sum_{i=1}^I \|\tilde{\beta}_i\|^2/4$. If β is LLL-reduced with parameter α , $\|\tilde{\beta}_i\|^2 \leq \alpha^I \|\tilde{\beta}_I\|^2$ for all $i \leq I$, which, by the definition of I , is at most $4d^2$. So $\|s_n\|^2 \leq ((n - I) + I\alpha^I)d^2 \leq n\alpha^n d^2$. The standard choice of $\alpha = 4/3$ gives $\|s_n\|^2 \leq 2^n d^2$. All of this is formalized in the final section, which culminates in the main theorem.

To avoid having to prove that a shortest vector exists, we use the definition $\inf\{\|u - t\| : u \in L\}$ for d instead of $\min\{\|u - t\| : u \in L\}$ and rephrase the arguments above to allow $\|v\|$ to exceed d by a small constant factor ϵ . This workaround and its details are contained in the section on the closest distance and negligibly change the rest of the proof.

theory *Babai-Algorithm*

imports *LLL-Basis-Reduction.LLL*

HOL.Archimedean-Field

HOL-Analysis.Inner-Product

begin

fun *calculate-c*:: *rat vec* $\Rightarrow$ *rat vec list* $\Rightarrow$ *nat* $\Rightarrow$ *int* **where**

calculate-c s L1 n = *round* ((*s* $\cdot$ (*L1*! ((*dim-vec s*) - *n*)))) / (*sq-norm-vec* (*L1*! ((*dim-vec s*) - *n*))))

fun *update-s*:: *rat vec* $\Rightarrow$ *rat vec list* $\Rightarrow$ *rat vec list* $\Rightarrow$ *nat* $\Rightarrow$ *rat vec* **where**

update-s sn M Mt n = ((*rat-of-int* (*calculate-c sn Mt n*)) $\cdot_v$ *M*) / ((*dim-vec sn*) - *n*))

fun *Babai-Help*:: *rat vec* $\Rightarrow$ *rat vec list* $\Rightarrow$ *rat vec list* $\Rightarrow$ *nat* $\Rightarrow$ *rat vec* **where**

Babai-Help s M Mt 0 = *s* |

Babai-Help s M Mt (Suc n) = (*let* *B* = (*Babai-Help s M Mt n*) *in* *B* - (*update-s B M Mt (Suc n)*))

definition *Babai*:: *rat vec* $\Rightarrow$ *rat vec list* $\Rightarrow$ *rat vec* **where**

Babai s M = *Babai-Help s M* (*gram-schmidt* (*dim-vec s*) *M*) (*dim-vec s*)

```

end
theory Babai
  imports Babai-Algorithm

```

```

begin

```

This theory contains the proof of correctness of the algorithm. The main theorem is "theorem Babai-Correct", under the locale "Babai-with-assms". To use the theorem, one needs to show that lattice, the vectors in the lattice basis, and the target vector all have the same dimension, that the lattice basis vectors are linearly independent and form an invertible matrix, and that the lattice basis is LLL-weakly-reduced.

2 Locale setup for Babai

```

locale Babai =
  fixes M :: int vec list
  fixes t :: rat vec
  assumes length-M: length M = dim-vec t
begin

```

```

abbreviation n where n ≡ length M
definition α where (α::rat) = 4/3
sublocale LLL n n M α⟨proof⟩

```

```

abbreviation coset::rat vec set where coset≡{(map-vec rat-of-int x)-t|x. x∈L}
abbreviation Mt where Mt ≡ gram-schmidt n (RAT M)

```

```

definition s :: nat ⇒ rat vec where
  s i = Babai-Help (uminus t) (RAT M) Mt i

```

```

definition closest-distance-sq:: real where
  closest-distance-sq = Inf {real-of-rat (sq-norm x::rat) |x. x ∈ coset}
end

```

Locale setup with additional assumptions required for main theorem

```

locale Babai-with-assms = Babai+
  fixes mat-M mat-M-inv:: rat mat
  assumes basis: lin-indep M
  defines mat-M ≡ mat-of-cols n (RAT M)
  defines mat-M-inv ≡

```

(if (invertible-mat mat-M) then SOME B. (inverts-mat B mat-M) $\wedge$ (inverts-mat
 mat-M B) else (0_m n n))
 assumes inv:invertible-mat mat-M
 assumes reduced:weakly-reduced M n
 assumes non-trivial:0<n
 begin

lemma dim-vecs-in-M:
 shows $\forall v \in \text{set } M. \text{dim-vec } v = \text{length } M$
 <proof>

lemma inv1:mat-M * mat-M-inv = 1_m n
 <proof>

lemma inv2:mat-M-inv * mat-M = 1_m n
 <proof>

sublocale rats: vec-module TYPE(rat) n<proof>

lemma M-dim: dim-row mat-M = n dim-col mat-M = n
 <proof>

lemma M-inv-dim: dim-row mat-M-inv = n dim-col mat-M-inv = n
 <proof>

lemma Babai-to-Help:
 shows $s \cdot n = \text{Babai-Algorithm.Babai } (\text{uminus } t) \text{ (RAT } M)$
 <proof>

3 Coordinates

This section sets up the use of the lattice basis and its GS orthogonalization
 as coordinate systems and some properties of that coordinate system. The
 important lemma here is coord-invariance, which shows that after step i of
 the algorithm, all coordinates (in both systems) after n-i are invariant.

definition lattice-coord :: rat vec $\Rightarrow$ rat vec
 where lattice-coord a = mat-M-inv *_v a

lemma dim-preserve-lattice-coord:
 fixes v::rat vec
 assumes dim-vec v=n

shows $\text{dim-vec } (\text{lattice-coord } v) = n$ $\langle \text{proof} \rangle$
lemma *vec-to-col*:
assumes $i < n$
shows $(\text{RAT } M)!i = \text{col mat-}M\ i$
 $\langle \text{proof} \rangle$

lemma *unit*:
assumes $i < n$
shows $\text{lattice-coord } ((\text{RAT } M)!i) = \text{unit-vec } n\ i$
 $\langle \text{proof} \rangle$

lemma *linear*:
fixes $i::\text{nat}$
fixes $v1::\text{rat vec}$
and $v2::\text{rat vec}$
and $q::\text{rat}$
assumes $\text{dim-vec } v1 = n$
assumes $\text{dim-}2:\text{dim-vec } v2 = n$
assumes $0 \leq i$
assumes $\text{dim-}i:i < n$
shows $(\text{lattice-coord } (v1 + (q \cdot_v v2)))\$i = (\text{lattice-coord } v1)\$i + q * ((\text{lattice-coord } v2)\$i)$
 $\langle \text{proof} \rangle$

lemma *sub-s*:
fixes $i::\text{nat}$
assumes $0 \leq i$
assumes $i < n$
shows $s\ (\text{Suc } i) = (s\ i) -$
 $(\text{rat-of-int } (\text{calculate-c } (s\ i)\ Mt\ (\text{Suc } i)) \cdot_v (\text{RAT } M)!(\text{dim-vec } (s\ i)) - (\text{Suc } i)))$
 $\langle \text{proof} \rangle$

lemma *M-locale-1*:
shows $\text{gram-schmidt-fs-Rn } n\ (\text{RAT } M)$
 $\langle \text{proof} \rangle$

lemma *M-locale-2*:
shows $\text{gram-schmidt-fs-lin-indpt } n\ (\text{RAT } M)$
 $\langle \text{proof} \rangle$

lemma *more-dim*: $\text{length } (\text{RAT } M) = n$
 $\langle \text{proof} \rangle$

lemma *Mt-gso-connect*:
fixes $j::\text{nat}$
assumes $j < n$
shows $Mt!j = \text{gs.gso } j$

$\langle proof \rangle$

lemma *access-index-M-dim*:

assumes $0 \leq i$

assumes $i < n$

shows $dim-vec \ (map \ of-int-hom.vec-hom \ M \ ! \ i) = n$

$\langle proof \rangle$

lemma *s-dim*:

fixes $i::nat$

assumes $i \leq n$

shows $dim-vec \ (s \ i) = n \wedge (s \ i) \in carrier-vec \ n$

$\langle proof \rangle$

lemma *dim-vecs-in-Mt*:

fixes $i::nat$

assumes $i < n$

shows $dim-vec \ (Mt!i) = n$

$\langle proof \rangle$

lemma *upper-tri*:

fixes $i::nat$

and $j::nat$

assumes $j > i$

assumes $j < n$

shows $((RAT \ M)!i) \cdot (Mt!j) = 0$

$\langle proof \rangle$

lemma *one-diag*:

fixes $i::nat$

assumes $0 \leq i$

assumes $i < n$

shows $((RAT \ M)!i) \cdot (Mt!i) = sq-norm \ (Mt!i)$

$\langle proof \rangle$

lemma *coord-invariance*:

fixes $j::nat$

fixes $k::nat$

fixes $i::nat$

assumes $k \leq j$

assumes $j+i \leq n$

assumes $k > 0$

shows $(lattice-coord \ (s \ (j+i)))\$ (n-k) = (lattice-coord \ (s \ j))\$ (n-k) \\ \wedge (s \ (j+i)) \cdot Mt!(n-k) = (s \ j) \cdot Mt!(n-k)$

$\langle proof \rangle$

lemma *small-orth-coord*:

fixes $i::nat$

assumes $1 \leq i$

assumes $i \leq n$

shows $abs \ (\ (s \ i) \cdot Mt!(n-i)) \leq \ (sq\text{-}norm \ (Mt!(n-i)))*(1/2)$
 $\langle proof \rangle$
lemma *lattice-carrier*: $L \subseteq carrier\text{-}vec \ n$
 $\langle proof \rangle$

4 Lattice Lemmas

lemma *lattice-sum-close*:
fixes $u::int \ vec$ **and** $v::int \ vec$
assumes $u \in L \ v \in L$
shows $u+v \in L$
 $\langle proof \rangle$

lemma *lattice-smult-close*:
fixes $u::int \ vec$ **and** $q::int$
assumes $u \in L$
shows $q \cdot_v u \in L$
 $\langle proof \rangle$

lemma *smult-vec-zero*:
fixes $v :: 'a::ring \ vec$
shows $0 \cdot_v v = 0_v \ (dim\text{-}vec \ v)$
 $\langle proof \rangle$

lemma *coset-s*:
fixes $i::nat$
assumes $i \leq n$
shows $s \ i \in coset$
 $\langle proof \rangle$

lemma *subtract-coset-into-lattice*:
fixes $v::rat \ vec$
fixes $w::rat \ vec$
assumes $v \in coset$
assumes $w \in coset$
shows $(v-w) \in of\text{-}int\text{-}hom.vec\text{-}hom' \ L$
 $\langle proof \rangle$

lemma *t-in-coset*:
shows $uminus \ t \in coset$
 $\langle proof \rangle$

5 Lemmas on closest distance

lemma *closest-distance-sq-pos*: $closest\text{-}distance\text{-}sq \geq 0$
 $\langle proof \rangle$

definition *witness*::rat vec $\Rightarrow$ rat $\Rightarrow$ bool
where *witness* *v* *eps-closest* = (*sq-norm* *v* $\leq$ *eps-closest* $\wedge$ *v* $\in$ *coset* $\wedge$ *dim-vec* *v* = *n*)

definition *epsilon*::real **where** *epsilon* = 11/10

definition *close-condition*::rat $\Rightarrow$ bool
where *close-condition* *eps-closest* $\equiv$
 (if *closest-distance-sq* = 0 then 0 $\leq$ *real-of-rat* *eps-closest*
 else *real-of-rat* (*eps-closest*) > *closest-distance-sq*)
 $\wedge$ (*real-of-rat* (*eps-closest*) $\leq$ *epsilon* * *closest-distance-sq*)

lemma *close-rat*:
obtains *eps-closest*::rat
where *close-condition* *eps-closest*
 <proof>

definition *eps-closest*::rat
where *eps-closest* = (if $\exists r$. *close-condition* *r* then SOME *r*. *close-condition* *r*
 else 0)

lemma *eps-closest-lemma*: *close-condition* *eps-closest*
 <proof>

lemma *rational-tri-ineq*:
fixes *v*::rat vec
fixes *w*::rat vec
assumes *dim-vec* *v* = *dim-vec* *w*
shows (*sq-norm* (*v*+*w*)) $\leq$ 4*(*Max* {(*sq-norm* *v*), (*sq-norm* *w*)})
 <proof>

lemma *witness-exists*:
shows $\exists v$. *witness* *v* *eps-closest*
 <proof>

6 More linear algebra lemmas

lemma *carrier-Ms*:
shows *mat-M* $\in$ *carrier-mat* *n* *n* *mat-M-inv* $\in$ *carrier-mat* *n* *n*
 <proof>

lemma *carrier-L*:
fixes *v*::rat vec
assumes *dim-vec* *v* = *n*
shows *lattice-coord* *v* $\in$ *carrier-vec* *n*
 <proof>

lemma *sumlist-index-commute*:
fixes *Lst*::rat vec list

```

fixes  $i::nat$ 
assumes  $set\ Lst \subseteq carrier\_vec\ n$ 
assumes  $i < n$ 
shows  $(gs.sumlist\ Lst)\$i = sum\_list\ (map\ (\lambda j. (Lst!\$j)\$i)\ [0..<(length\ Lst)])$ 
 $\langle proof \rangle$ 

```

lemma *mat-mul-to-sum-list*:

```

fixes  $A::rat\ mat$ 
fixes  $v::rat\ vec$ 
assumes  $dim\_vec\ v = dim\_col\ A$ 
assumes  $dim\_row\ A = n$ 
shows  $A *_v v = gs.sumlist\ (map\ (\lambda j. v\$j \cdot_v (col\ A\ j))\ [0..<dim\_col\ A])$ 
 $\langle proof \rangle$ 

```

lemma *recover-from-lattice-coord*:

```

fixes  $v::rat\ vec$ 
assumes  $dim\_vec\ v = n$ 
shows  $v = gs.sumlist\ (map\ (\lambda i. (lattice\_coord\ v)\$i \cdot_v (RAT\ M)!\$i)\ [0..<n])$ 
 $\langle proof \rangle$ 

```

lemma *sumlist-linear-coord*:

```

fixes  $Lst::int\ vec\ list$ 
assumes  $\bigwedge i. i < length\ Lst \implies dim\_vec\ (Lst!\$i) = n$ 
shows  $lattice\_coord\ (map\_vec\ rat\_of\_int\ (sumlist\ Lst)) = gs.sumlist\ (map\ lattice\_coord\ (RAT\ Lst))$ 
 $\langle proof \rangle$ 

```

lemma *integral-sum*:

```

fixes  $l::nat$ 
assumes  $\bigwedge j1. j1 < l \implies$ 
 $map\ f\ [0..<l]!\$j1 \in \mathbb{Z}$ 
shows  $sum\_list$ 
 $(map\ f\ [0..<l]) \in \mathbb{Z}$ 
 $\langle proof \rangle$ 

```

lemma *int-coord*:

```

fixes  $i::nat$ 
assumes  $0 \leq i$ 
assumes  $i < n$ 
fixes  $v::int\ vec$ 
assumes  $v \in L$ 
assumes  $dim\_vec\ v = n$ 
shows  $(lattice\_coord\ (map\_vec\ rat\_of\_int\ v))\$i \in \mathbb{Z}$ 
 $\langle proof \rangle$ 

```

lemma *int-coord-for-rat*:

```

fixes  $i::nat$ 
assumes  $0 \leq i$ 
assumes  $i < n$ 
fixes  $v::rat\ vec$ 
assumes  $v \in of-int-hom.vec-hom\ L$ 
assumes  $dim-vec\ v = n$ 
shows  $(lattice-coord\ v)\$i \in \mathbb{Z}$ 
 $\langle proof \rangle$ 

```

7 Coord-Invariance

This section shows that the algorithm output matches true closest (or near-closest) vector in some trailing coordinates.

definition I **where**

```

 $I = (if\ (\{i \in \{0..<n\}. ((sq-norm\ (Mt!i)::rat)) \leq 4*eps-closest\}::nat\ set) \neq \{\}$ 
 $\quad then\ Max\ (\{i \in \{0..<n\}. ((sq-norm\ (Mt!i)::rat)) \leq 4*eps-closest\}::nat\ set)\ else$ 
 $-1)$ 

```

lemma $I-geq$:

shows $I \geq -1$

$\langle proof \rangle$

lemma $I-leq$:

shows $I < n$

$\langle proof \rangle$

lemma $index-geq-I-big$:

fixes $i::nat$

assumes $i > I$

assumes $i < n$

shows $((sq-norm\ (Mt!i)::rat)) > 4*eps-closest$

$\langle proof \rangle$

lemma $scalar-prod-gs-from-lattice-coord$:

fixes $i::nat$

fixes $v::rat\ vec$

assumes $dim-vec\ v = n$

assumes $i < n$

shows $v \cdot Mt!i = sum-list\ (map\ (\lambda k. (lattice-coord\ v)\$k * (((RAT\ M)!k) \cdot Mt!i))$
 $[i..<n])$

$\langle proof \rangle$

lemma $correct-coord-help$:

fixes $i::nat$

assumes $i < (int\ n) - I$

assumes $witness\ v\ (eps-closest)$

assumes $0 < i$

shows $(lattice-coord\ (s\ i))\$(n-i) = (lattice-coord\ v)\$(n-i)$

$\wedge ((s\ i) \cdot Mt!(n-i) = v \cdot Mt!(n-i))$
 $\langle proof \rangle$

lemma *correct-coord*:
fixes $v::rat\ vec$
fixes $k::nat$
assumes $witness\ v\ eps-closest$
assumes $I < k$
assumes $k < n$
shows $(s\ n) \cdot Mt!(k) = v \cdot Mt!(k)$
 $\langle proof \rangle$

8 Main Theorem

This section culminates in the main theorem.

lemma *sq-norm-from-Mt*:
fixes $v::rat\ vec$
assumes $v-carr:v \in carrier-vec\ n$
shows $sq-norm\ v = sum-list\ (map\ (\lambda i. (v \cdot Mt!i) \hat{\wedge} 2 / (sq-norm\ (Mt!i)))\ [0..<n])$
 $\langle proof \rangle$

lemma *bound-help*:
fixes $N::nat$
shows $real-of-rat\ ((rat-of-int\ N) * \alpha \hat{\wedge} N) * epsilon \leq 2 \hat{\wedge} N$
 $\langle proof \rangle$

lemma *present-bound-nicely*:
fixes $N::nat$
shows $real-of-rat\ ((rat-of-int\ N) * \alpha \hat{\wedge} N * eps-closest) \leq 2 \hat{\wedge} N * closest-distance-sq$
 $\langle proof \rangle$

lemma *basis-decay*:
fixes $i::nat$
fixes $j::nat$
assumes $i < n$
assumes $i + j < n$
shows $sq-norm\ (Mt!i) \leq \alpha \hat{\wedge} j * sq-norm(Mt!(i+j))$
 $\langle proof \rangle$

lemma *basis-decay-cor*:
fixes $i::nat$
fixes $j::nat$
assumes $i < n$
assumes $j < n$
assumes $i \leq j$
shows $sq-norm\ (Mt!i) \leq \alpha \hat{\wedge} n * sq-norm(Mt!j)$
 $\langle proof \rangle$

```

theorem Babai-Correct:
  shows real-of-rat ((sq-norm (s n))::rat)  $\leq 2^{\wedge}n * \text{closest-distance-sq} \wedge s\ n \in \text{coset}$ 
  <proof>

end
end

```

References

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